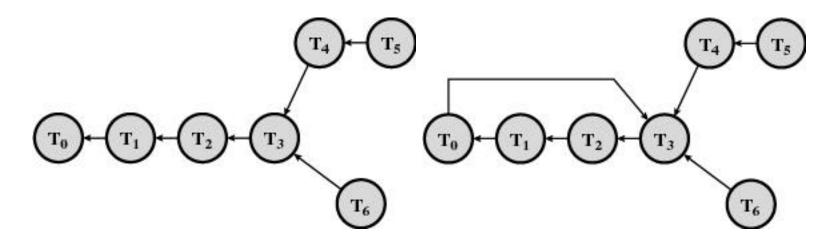
```
/* Data object D<sub>i</sub> receiving a lock_request(T<sub>i</sub>) */
                                                                                              /* Transaction T<sub>i</sub> receiving an update message */
     if (Locked_by(D_i) == null)
                                                                                                   if (Wait_for(T_i) != Wait_for(T_i))
          send(granted);
     else
                                                                                                        Wait\_for(T_i) = Wait\_for(T_i);
                                                                                                   if (intersect(Wait_for(T_i), Request_Q(T_i)) = null)
          send not granted to T<sub>i</sub>;
                                                                                                        update(Wait_for(Ti), Request_Q(Tj);
          send Locked_by(D<sub>i</sub>) to T<sub>i</sub>
                                                                                                    else
                                                                                                              DECLARE DEADLOCK;
                                                                                                              /* initiate deadlock resolution as follows */
/* Transaction T<sub>i</sub> makes a lock request for data object D<sub>i</sub> */
                                                                                                              /* T<sub>i</sub> is chosen as the transaction to be aborted */
     send lock_request(T_i) to D_i;
                                                                                                              /* T<sub>i</sub> releases all the data objects it holds */
     wait for granted/not granted;
                                                                                                              send_clear(T<sub>i</sub>, Held_by(T<sub>i</sub>));
     if (granted)
                                                                                                              allocate each data object Di held by Ti to the first
                                                                                                                   requester T_k in Request_Q(T_i);
          Locked_by(D_i) = T_i;
                                                                                                              for (every transaction T<sub>n</sub> in Request_Q(Tj) requesting
          Held_by(T_i) = ;
                                                                                                                   data object Di held by Ti)
     else /* suppose D<sub>i</sub> is being used by transaction T<sub>i</sub> */
                                                                                                                   Enqueue(T_n, Request\_Q(T_k));
          Held_by(T_i) = T_i;
          Enqueue(T_i, Request_Q(T_i));
          if (Wait_for(Tj) == null)
               Wait_for(T_i) = T_i;
                                                                                              /* Transaction T_k receiving a clear (T_i, T_k) message */
          else
                                                                                                    purge the tuple having Tj as the requesting transaction from
               Wait_for(T<sub>j</sub>) = Wait_for(T<sub>j</sub>);
                                                                                                   Request_Q(Tk);
          update(Wait_for(T<sub>i</sub>), Request_Q(T<sub>i</sub>));
```

Figure 14.14 A Distributed Deadlock Detection Algorithm



Transaction	Wait_for	Held_by	Request_Q	Transaction	Wait_for	Held_by	Request_Q
$T_0$	nil	nil	$T_1$	$T_0$	$T_0$	$T_3$	$T_1$
$T_1$	$T_0$	$T_0$	$T_2$	$T_1$	$T_0$	$T_0$	$T_2$
$T_2$	$T_0$	$T_1$	$T_3$	$T_2$	$T_0$	$T_1$	$T_3$
$T_3$	$T_0$	$T_2$	$T_4, T_6$	$T_3$	$T_0$	$T_2$	$T_4, T_6, T_0$
$T_4$	$T_0$	$T_3$	$T_5$	$T_4$	$T_0$	$T_3$	$T_5$
$T_5$	$T_0$	$T_4$	nil	$T_5$	$T_0$	$T_4$	nil
$T_6$	$T_0$	$T_3$	nil	$T_6$	$T_0$	$T_3$	nil

(a) State of system before request

(b) State of system after  $T_0$  makes a request to  $T_3$ 

Figure 14.15 Example of Distributed Deadlock Detection Algorithm of Figure 14.14