Instruction Sets Ch 10-11

Characteristics
Operands
Operations
Addressing
Instruction Formats

Instruction Set

(käskykanta)

- Collection of instructions that CPU understands
- Only interface to CPU from outside
- CPU executes a program ⇔ CPU executes given instructions “one at a time”
  - fetch-execute cycle

Fig. 10.1 (Fig. 9.1 [Stal99])

Machine Instruction

Fig. 10.1 (Fig. 9.1 [Stal99])

- Opcode
  - What should I do? Math? Move? Jump?
- Source operand references
  - Where is the data to work on? Reg? Memory?
- Result operand reference
  - Where should I put the result? Reg? Memory?
- Next instruction reference
  - Where is the next instruction? Default? Jump?

Instruction Representation

- Bit presentation:
  - binary program
- Assembly language
  - symbolic program
- Symbolic assembly language

0x2465A080
Opcode, operands
LOAD R1,=0x6678
Symbolic opcode
Virtual or physical address?
LOAD R1,TotalSum
Symbolic value?

Fig. 10.11 (Fig. 9.11 [Stal99])

Instruction Set Design (5)

- Operation types (operaatiotyyppi)
  - How many? What type? Simple? Complex?
- Data types (tietotyyppi)
  - Just a few? Many?
- Instruction format (käskyn muoto)
  - fixed length? Varying length? Nr of operands?
- Number of addressable registers
  - too many ⇒ too long instructions?
  - too few ⇒ too hard to optimise code?
- Addressing (tiedon osoitus)
  - What modes to use to address data and when?

Good Instruction Set (2)

- Good target for compiler
  - Easy to compile?
  - Possible to compile code that runs fast?
  - Easy to compile code that runs fast?
- Allows fast execution of programs
  - How many meaningless instructions per second? MIPS? GFLOPS?
  - How fast does my program run?
    - Solve linear system of 1000 variables?
    - Set of data base queries?
    - Connect a phone call in reasonable time?

Good Instruction Set (contd) (5)

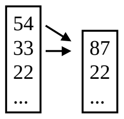
- Beautiful & Aesthetic
  - Orthogonal (ortogonaalinen)
    - Simple, no special registers, no special cases, any data type or addressing mode can be used with any instruction
  - Complete (täydellinen)
    - Lots of operations, good for all applications
  - Regular (säännöllinen)
    - Specific instruction field has always same meaning
  - Streamlined (virtaviivainen)
    - Easy to define what resources are used

Good Instruction Set (contd) (2)

- Easy to implement
  - 18 months vs. 36 months?
  - Who will be 1st in market? Who will get development monies back and who will not?
- Scalability (skaalautuva)
  - Speed up clock speed 10X, does it work?
  - Double address length, does design extend?
    - E.g., 32 bits ⇒ 64 bits ⇒ 128 bits?

Number of Operands? (4)

- 3? ADD A,B,C Mem(A) ← mem(B) + mem(C)
  - Normal case now ADD R1, R2, R3 r1 ← r2+r3
- 2? ADD R1, R2 r1 ← r1+r2
  - 1 operand and result the same
- 1? ADD A acc ← acc+mem(A)
  - 1 operand and result in implicit <u>accumulator</u> (register)
- 0? ADD
  - All operands and result in implicit <u>stack</u>



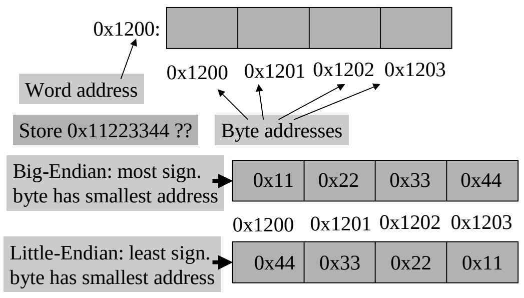
Instruction Set Architecture (ISA) Basic Classes

- Accumulator architecture
- Stack architecture
- General Purpose Register (GPR) architecture
  - only one type of registers, good for all
  - 2 or 3 operands
- Load/Store architecture
  - only load/store instructions access memory
  - 3 operand ALU instructions

```
LOAD R3, C
LOAD R2,B
ADD   R1,R2,R3
STORE R1,A
```

Big vs. Little Endian (3)

- How are multi-byte values stored


0x1200:
Word address
0x1200 0x1201 0x1202 0x1203
Store 0x11223344 ??
Byte addresses
Big-Endian: most sign. byte has smallest address
0x11 0x22 0x33 0x44
0x1200 0x1201 0x1202 0x1203
Little-Endian: least sign. byte has smallest address
0x44 0x33 0x22 0x11

Big vs. Little Endian

- Address of multi-byte data items is the same in both representations
- Only internal byte order varies
- Must decide one way or the other
  - Math circuits must know which presentation used
    - Little-Endian may be faster ….
  - Must consider when moving data via network
- Power-PC: bi-endian - both modes at use
  - can change it per process basis
  - kernel mode selected separately

Data (Operands, Result) Location

- Register
  - close, fast
  - limited number of them
  - need to load/store values from/to memory sometimes (often)
    - big problem! 50% of compiler time to decide
    - register allocation problem
- Memory
  - far away
  - only possibility for large data sets
    - vectors, arrays, sets, tables, objects, ...

acc r2, r8

register stack f4, f15

memory stack (hw regs have mem addresses)

memory 0x345670

cache?

Aligned Data (4)

2 byte (16-bit) half-word has byte address:	0010…10010
4 byte (32-bit) word has byte address:	0010…10100
8 byte (64-bit) doubleword has byte address:	0010…11000

- Aligned data
  - faster memory access
    - 32-bit data loaded as one memory load
- Non-aligned data
  - saves mem, more bus traffic!
    - 32-bit non-aligned data requires 2 memory loads (each 4 bytes) and combining data into one 32-bit data item

11	22	33	44

		11	22
33	44		

Data Types (8)

• Address	16b, 32b, 64b, 128b?
• Integer	16b, 32b, 64b?
• Floating point	32b, 64b, 82b?
• Decimal	18 digits (9 bytes) packed decimal?
• Character	1 byte = 8b IRA = ASCII, EBCDIC?
• String	finite, arbitrary length? Length denotation?
• Logical data	1 bit (Boolean value, bit field)?
• Vector, array, record, ….	

Size of Operand

• 1 word, 32 bits	int, float, addr
• 2 words, 64 bits	double float, addr
• 4 words, 128 bits	addr
• 1 byte (8 bits)	char
• 2 bytes	short int
• 1 bit	logical values

Example: Pentium II Data Types

- General data types
  - 8-bit byte
  - 16-bit word
  - 32-bit doubleword
  - 64-bit quadword
- Not aligned
- Little Endian
- Specific data types — Table 10.2 (Tbl. 9.2 [Stal99])
- Numerical data types — Figure 10.4 (Fig. 9.4 [Stal99])

(for Pentium II)

Operation Types

Table 10.3 (Tbl 9.3 [Stal99])

- Data transfer
  - CPU ↔ memory
- ALU operations
  - INT, FLOAT, BOOLEAN, SHIFT, CONVERSION
- I/O
  - read from device, start I/O operation
- Transfer of control
  - jump, branch, call, return, IRET, NOP
- System control
  - HALT, SYSENTER, SYSEXIT, …
  - CPUID returns current HW configuration
    - size of L1 & L2 caches, etc

Table 10.4 (Tbl 9.4 [Stal99])

Data References (2)

- Where is data?
  - in memory
  - in registers
  - in instruction itself
- How to refer to data?
  - various addressing modes
  - multi-phase data access
    - how is data location determined (addressing mode)
    - compute data address (register? effective address?)
    - access data

Addressing Modes (Ch 10)

Fig. 11.1 (Fig 10.1 [Stal99])

Table. 11.1 (Tbl 10.1 [Stal99])

- Immediate Data in instruction
- Direct Memory address of data in instruction
- Indirect Address of memory address of data in instruction (pointer)
- Register Data in register (best case?)
- Register Indirect Register has memory address (pointer)
- Displacement Addr = reg value + constant
- Stack Data in stack pointed by some register

Displacement Address

- Effective address = (R1) + A

Contents of R1 Constant from instruction

- Constant is often small (8 bits, 16 bits?)
- Many uses

– PC relative	JUMP <u>-40(PC)</u>
– Base register address	CALL <u>Summation(BX)</u>
– Array index	ADDF F2, F2, <u>Table(R5)</u>
– Record field	MUL F4, F6, <u>Salary(R8)</u>
– Stack references	STORE F2, <u>-4(FP)</u>

More Addressing Modes

size of operand

- Autoincrement $EA = (R), R \leftarrow (R) + S$

  E.g., CurrIndex = i++;
- Autodecrement $R \leftarrow (R) - S, EA = (R)$

  E.g., CurrIndex = --i;
- Autoincrement deferred $EA = Mem(R), R \leftarrow (R) + S$

  E.g., Sum = Sum + (*ptrX++);
- Scaled $EA = A + (R_j) + (R_i) * S$

  E.g., double X;
  X = Tbl[i][j];

Pentium II Addressing Modes

- Immediate
  - 1, 2, 4 bytes
- Register operand
  - 1, 2, 4, 8 byte registers
  - not all registers with every instruction
- Operands in Memory Fig. 11.2 (Fig 10.2 [Stal99])
  - compute effective address and combine with segment register to get linear address (virtual address)

Table 11.2 (Tbl 10.2 [Stal99])

Instruction Format (4)

- How to represent instructions in memory?
- How long instruction
  - – Descriptive or dense? Code size?
- Fast to load?
  - – In many parts?
  - – One operand description at a time?
- Fast to parse (I.e., split into logical components)?
  - – All instruction same size & same format?
  - – Very few formats?

Instruction Format (contd) (3)

- How many addressing modes?
  - – Fewer is better, but harder to compile to
- How many operands?
  - – 3 gives you more flexibility, but takes more space
- How many registers?
  - – 16 regs $\rightarrow$ need 4 bits to name it
  - – 256 regs $\rightarrow$ need 8 bits to name it
  - – need at least 16-32 for easy register allocation
  - – How many registers, that can be referenced in <u>one instruction</u> vs. referenced <u>overall</u>?

Instruction Format (contd) (3)

- How many register sets?
  - – A way to use more registers without forcing long instructions for naming them
  - – One register set for each subroutine call?
  - – One for indexing, one for data?
- Address range, number of bits in displacement
  - – more is better, but it takes space
- Address granularity
  - – byte is better, but word address is shorter

Pentium Instruction Set (5)

- CISC - Complex Instruction Set Computer
- At most one memory address
- "Everything" is optional
- "Nothing" is fixed
- Difficult to parse
  - – all latter fields and their interpretation depend on earlier fields

Fig. 11.8 (Fig 10.8 [Stal99])

Pentium Instruction Prefix Bytes (4)

Fig. 11.8 (Fig 10.8 (a)[Stal99])

- Instruction prefix (optional)
  - – LOCK - exclusive use of shared memory
  - – REP - repeat instruction for string characters
- Segment override (optional)
  - – override default segment register
  - – default is implicit, no need to store it every instruction
- Address size (optional)
  - – use the other (16 or 32 bit) address size
- Operand size (optional)
  - – use the other (16 or 32 bit) operand size

Pentium Instruction Fields (3)

Fig. 11.8 (Fig 10.8 (a)[Stal99])

- Opcode
  - – specific bit for byte size data
- Mod r/m (optional)
  - – data in reg (8) or in mem?
  - – which addressing mode of 24?
  - – can also specify opcode further for some opcodes
- SIB (optional) – Scale/Index/Base
  - – extra field needed for some addressing modes
  - – scale for scaled indexing
  - – index register
  - – base register

Pentium Instruction Fields (contd) (2)

Fig. 11.8

(Fig 10.8 (a)[Stal99])

- Displacement (optional)
  - for certain addressing modes
  - 1, 2, or 4 bytes
- Immediate (optional)
  - for certain addressing modes
  - 1, 2, or 4 bytes

PowerPC Instruction Format (7)

Fig. 11.9

(Fig 10.9 [Stal99])

- RISC - Reduced Instruction Set Computer
- Fixed length, just a few formats
- Only 2 addressing modes for data
- Only load/store instructions access memory
- 32 general purpose registers can be used everywhere
- Fixed data size
  - no string ops
- Simple branches
  - CR-field determines which compare result to use
  - L-bit determines whether a subroutine call
  - A-bit determines if branch is absolute or PC-relative

-- End of Chapters 10-11: Instruction Sets --

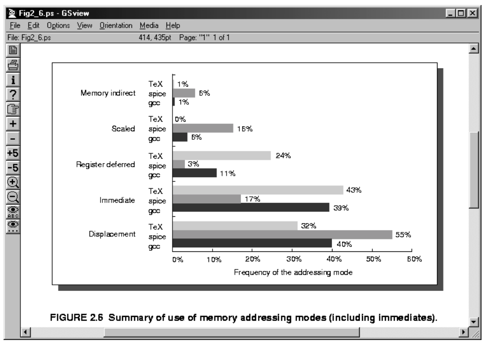


FIGURE 2.6 Summary of use of memory addressing modes (including immediates).

(Hennnessy-Patterson, Computer Architecture, 2nd Ed, 1996)