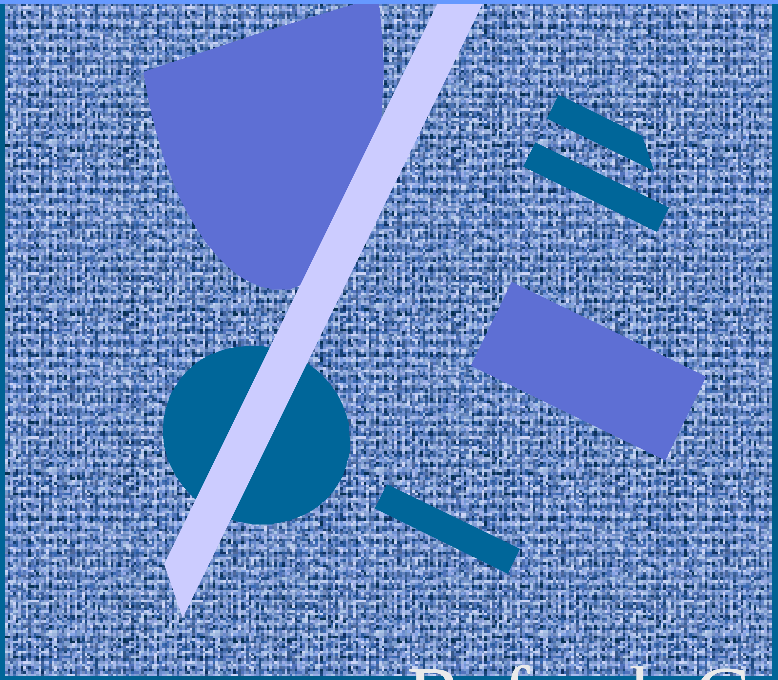
Computer System Overall Structure

Ch 1-7



Review

Overall Picture

Refresh Computer Organization I (TiTo)

Computer System

- Data movement, storage, and processing

Figs 1.3, 1.4

- Control

Figs 1.5, 1.6

Figs 3.2, 3.3, 3.9

- System and I/O Buses
- Internal and external memories
- Input/Output systems
- Operating Systems support

System & I/O Buses

- Bus configurations Fig 3.18
- Local (internal, memory) bus (sisäinen väylä)
  - – inside CPU chip
  - – connects CPU to cache
- System bus (systeemiväylä)
  - – connects CPU to memory
- I/O bus (I/O väylä)
  - – connects CPU & memory to I/O devices
- Implementation details later on

Internal and External Memories

- Memory hierarchy (muistihierarkia)

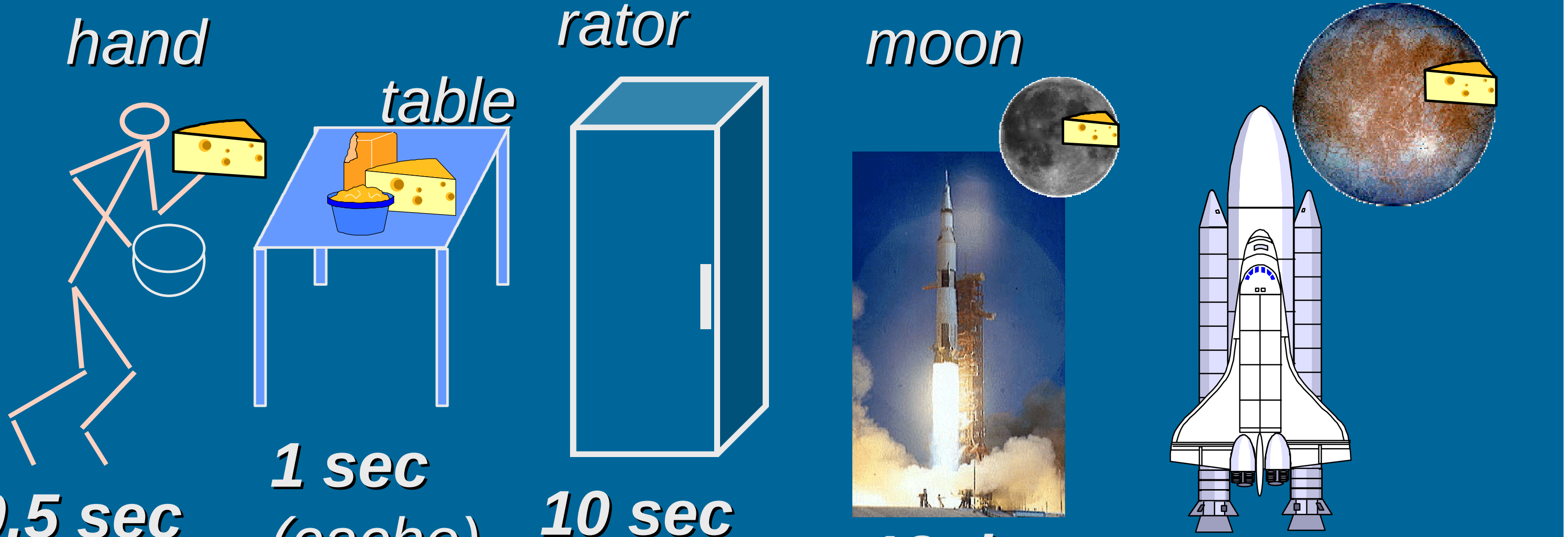
  Fig 4.1

  - Registers, L1 Cache, L2 Cache
  - Main memory, Disk cache
  - Disk, Optical, Tape
  - File server (local, via LAN)
  - Remote server (via WWW?)
- Storage capacity vs. access time (saantiaika)

  Fig 4.3 [Stal96]

Teemu’s Cheesecake (5)

Register, on-chip cache, memory, disk, and tape speeds relative to times locating cheese for the cheese cake you are baking...



hand

0.5 sec
(register)

table

1 sec
(cache)

refridge-
rator

10 sec
(memory)

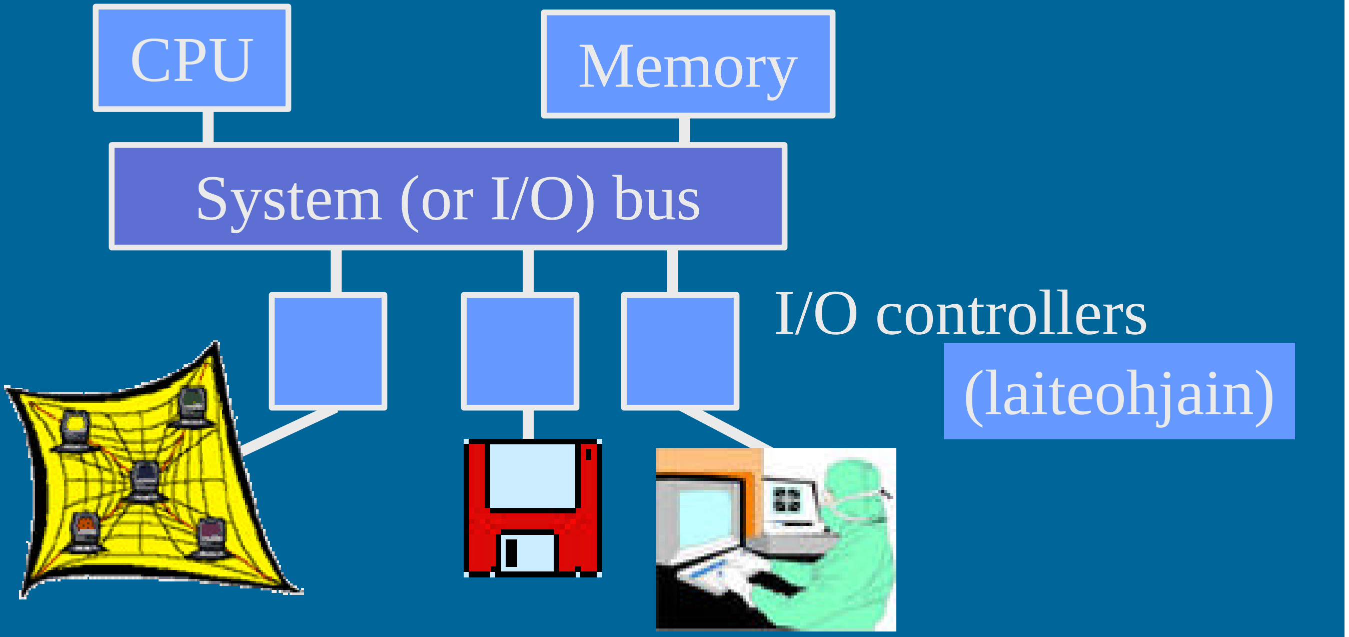
moon

12 days
(disk)

Europa
(Jupiter)

4 years
(tape)

Input/Output Systems (3)



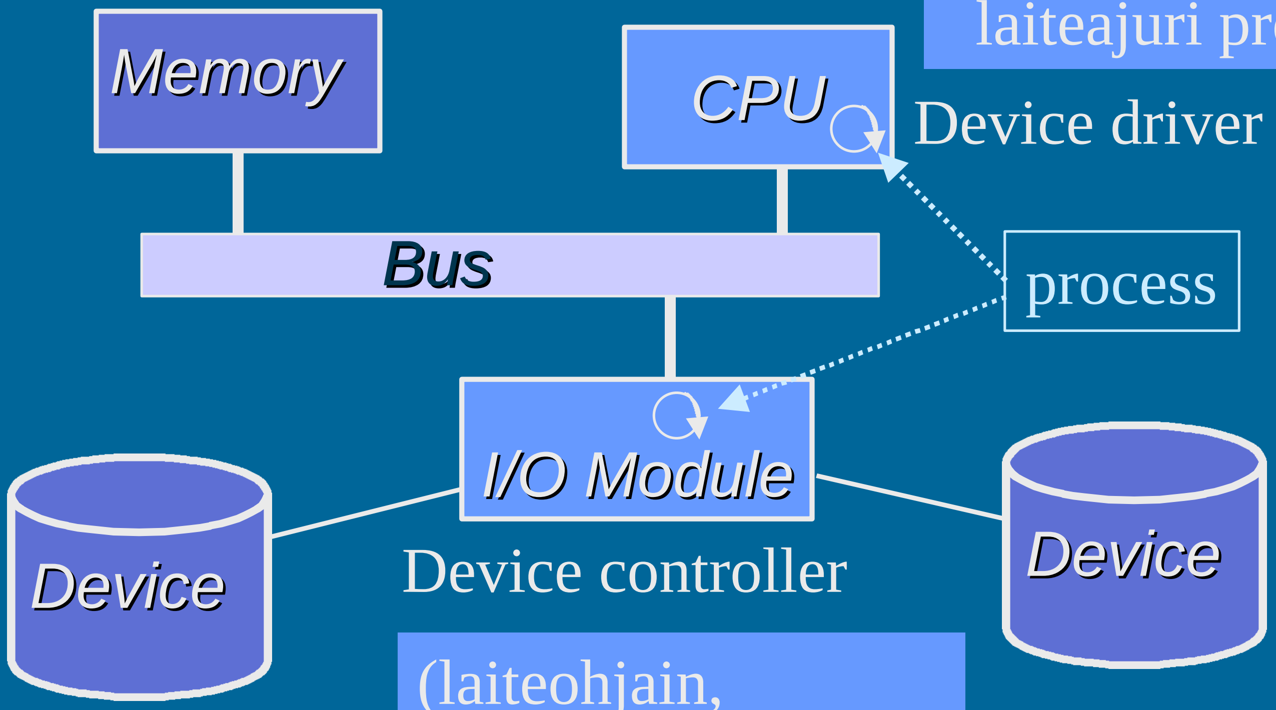
- Three categories
  - I/O with people — Video display, joy-stick, ...
  - I/O with machines — CD, disk, ...
  - Communication — Ethernet, token ring, ...

I/O Module (Device Controller) (1)

- I/O Module

Fig 6.4

(keskusyksikkö, laiteajuri prosessi)



Device driver

process

Device controller

(laiteohjain, laiteohjain prosessi)

Direct vs. Interrupt-driven I/O (2)

- Direct, I.e., programmed I/O (suora I/O)
  - CPU controls I/O directly
  - CPU spins (waits) while I/O device works
  - I/O device transfers one word at a time
- Interrupt-driven I/O (keskeyttävä I/O)
  - CPU gives one I/O command, does a process switch, and continues with some other work
  - when I/O is done, I/O controller interrupts the CPU, and original process is made ready to run again

Direct vs. Interrupt-driven I/O (contd) (2)

- Direct Memory Access (DMA)
  - – I/O controller can directly access memory
    - o/w access only to “data registers”
  - – interrupt CPU only after (a big) block transfer
- I/O channels and I/O processors
  - – I/O controller is smart
  - – I/O controller manages complete I/O jobs
    - each with many DMA transfers?
    - many I/O jobs in queue at a time?

Memory-Mapped I/O (3)

(muistiinkuvattu I/O)

- Each device controlled via device registers
  - data, status, control (laiterekisterit)
- Device registers are addressed similarly as memory
  - with normal read/write instructions (vs. specific machine instructions for I/O)
  - device controller acts also as a memory card
- Device registers are physically located in the device controller which recognises certain memory addresses belonging to it

SCSI - Small Computer System Interconnect (3)

- Parallel data interface
  - 8,16, or 32 parallel data lines (wires)
  - 9 control lines
- Max 7 devices
- Arbitration
  - select who can use
  - the one with the highest priority wins
  - priority = SCSI id selected for the device

Operating Systems Support

- User/computer interface

(käyttöliittymä)

Fig 7.1

- Resource manager

Fig 7.2

(resurssien hallinta)

- Process manager

(prosessien hallinta)

Fig 7.8

(prosessin tilat)

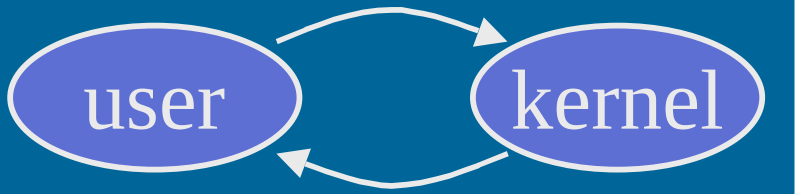
- Process Control Block (PCB)

(prosessin kontrollilohko)

Fig 7.9

Processor States

(suorittimen tilat)

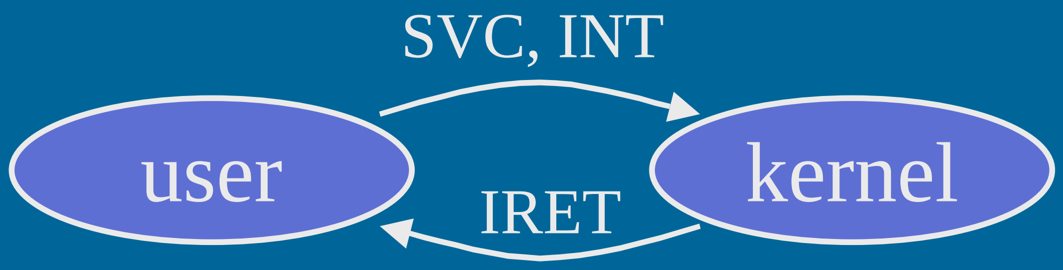


user

kernel

- User mode (normal mode) (käyttäjätila)
  - can use only non-privileged instructions
  - can access only memory in user-space
- Kernel mode (privileged mode) (etuoikeutettu tila)
  - can use all machine instructions, including privileged instructions (etuoikeutetut konekäskyt)
  - can access all memory, including kernel memory (KJ:n ytimen omat muistialueet)

Changing Processor Mode



SVC, INT

user

kernel

IRET

- User mode → kernel mode
  - – interrupt or explicit SVC instruction
  - – interrupt handler checks for rights to change mode

(keskeytyskäsittelijä)

- Kernel mode → user mode
  - – privileged machine instruction
  - – return from interrupt (e.g., IRET)
  - – returns control & restores previous mode

-- End of Chapter 1-7: Intro --