

Signaling Using the NSIS Framework; Messaging and End-host Mobility Applications

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Scenario: Battery monitoring

- Network of battery powered (wireless) routers ($R_1..R_3$) and three hosts ($H_1..H_3$)

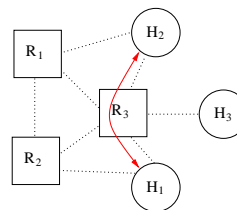
- Hosts H_1 and H_2 communicate very actively



- Router R_3 runs low on battery

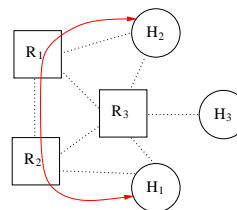
- What about connectivity of H_3 ?

- We can keep H_3 connected, if we have:

- Mechanism to alter routing
- **Battery status information**

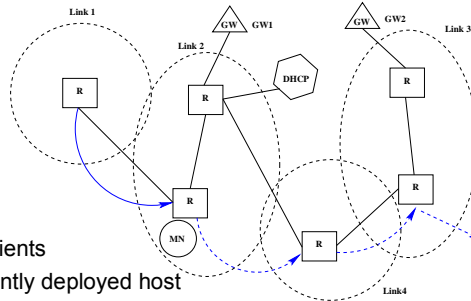


 Route, active communication
 Direct route exists





Scenario: Mobility



- Mesh access network
- Mobile Node (MN)
- Support for unmodified clients
 - Compatible with currently deployed host software
- Equal nodes in a decentralized topology
 - Co-operating, scalable, robust, autonomous



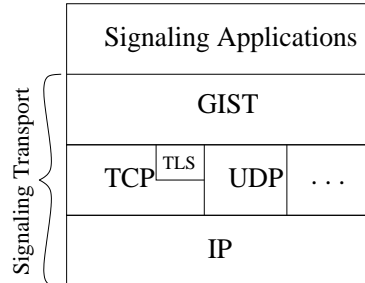
NSIS Framework

- Next Steps In Signaling WG @ IETF
- Signaling framework
- Signaling transport separated from application logic
 - GIST (General Internet Signaling Transport)
 - Uses existing protocols and methods (TCP, UDP, TLS...)
- NSLPs (NSIS Signaling Layer Protocol) make up the application layer (QoS, NAT/Firewall)
- Uses mainly TLV (Type – Length – Value) objects



GIST services

- GIST does all the 'dirty work'
- Next (NSLP level) node discovery
- Path-coupled transport
 - Reliable
 - Unreliable
- Channel security

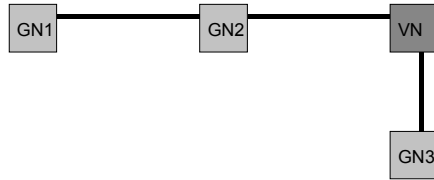


GIST NSLP node discovery

- When NSLP wants to send a message towards some host
- GIST finds the next GIST node on the path running the same NSLP
 - There may be multiple GIST (and non-GIST) nodes between sender and recipient



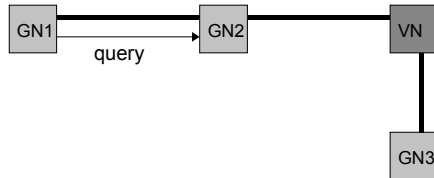
GIST NSLP node discovery



———— Data flow
==== Signaling flow



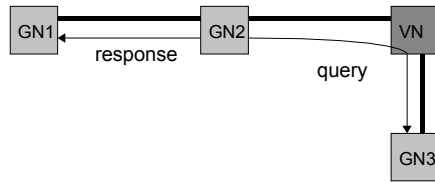
GIST NSLP node discovery



———— Data flow
==== Signaling flow



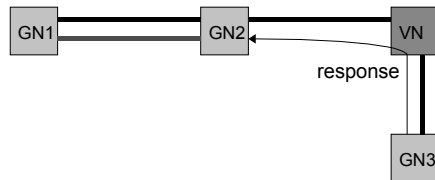
GIST NSLP node discovery



———— Data flow
——— Signaling flow



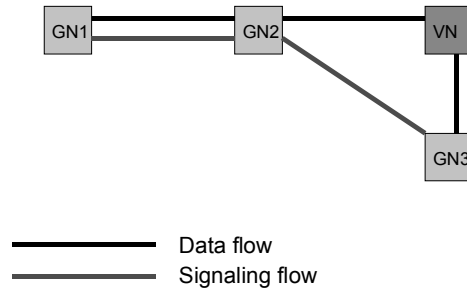
GIST NSLP node discovery



———— Data flow
——— Signaling flow

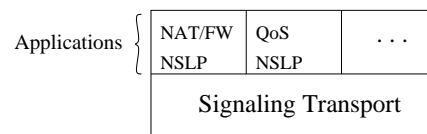


GIST NSLP node discovery



Signaling Applications (NSLPs)

- Use services provided by GIST
- May use services provided by other NSLPs
 - NAT/FW traversal
- Forwards its own messages (identified by NSLP ID) along the path towards signaling end host





NSIS Framework and GIST

- Anything to ask about NSIS or GIST?



Messaging NSLP: Outline

- Motivation
- Architecture
- Protocol Messages
- API
- Messaging Sessions
- Status and Next Steps



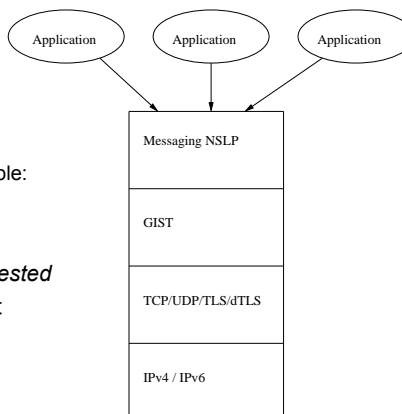
Messaging NSLP: Motivation

- A new kind of generic messaging protocol could have a lot of use cases
 - Router-to-Router communications
 - Host-to-Router communications
 - Router-to-Host communications
- New NSLP to provide generic messaging service to Messaging applications



Messaging NSLP: Layers

- Layered architecture
 - Messaging NSLP provides transport to messaging applications
 - Application of the first example: Monitoring application
- Messaging NSLP is *not interested* in content it carries – it is just transport





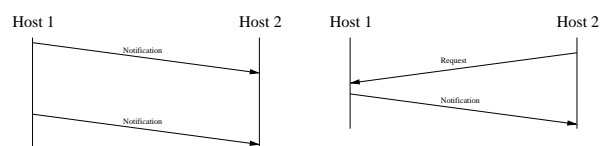
Messaging NSLP

- While GIST provides quite nice tools to signaling application developers, it still has some limitations
 - NSLP ID reservation bureaucracy
 - Quite complex API
- Messaging NSLP needs only one NSLP ID
 - Multiple messaging applications share the same GIST level identifier (NSLP ID)



Messaging NSLP: Messages

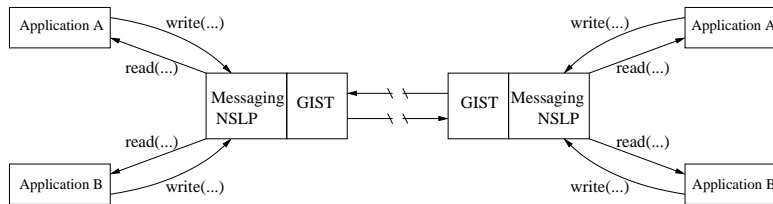
- Just three types of messages
 - Request
 - Notification
 - May be sent with or without request
 - Error





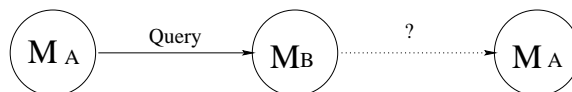
Messaging NSLP: API

- Simple API for Signaling Application development
 - Open(..) to open new messaging session
 - Write(..) to send data to remote application
 - Read(..) to read data sent by remote application
 - Close(..) to end messaging session



Messaging NSLP: Messaging Sessions

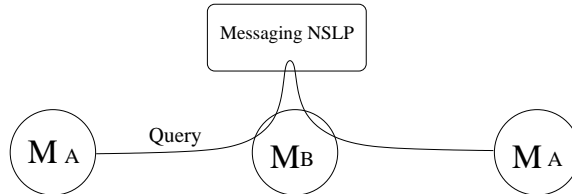
- Session opened by calling *Open(..)*
 - New Session identifier allocated
 - Triggers eventually GIST Query
 - The GIST peering process is invisible to application
- What if nodes running Messaging NSLP do not have the same Messaging applications?
- GIST decision whether to peer or not is based on NSLP IDs





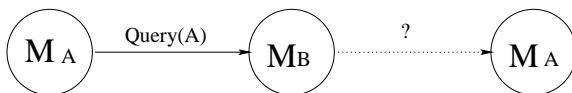
Messaging NSLP: Conditional peering

- Vanilla GIST can NOT always carry data within query messages
- Each Messaging NSLP node must act as relays
 - Works, but is very inefficient



Messaging NSLP: Conditional peering

- To get conditional peering really working, we need to modify GIST
 - Additional GIST object
 - We can use solution shown at previous slide as a backup





Messaging NSLP: Status and next steps

- First version of specification is ready
 - Works on paper...
 - ...and implementation has begun
- A couple of example applications are needed
- Tests and benchmarks
- Papers (including my MSc thesis...)
- Internet Draft?

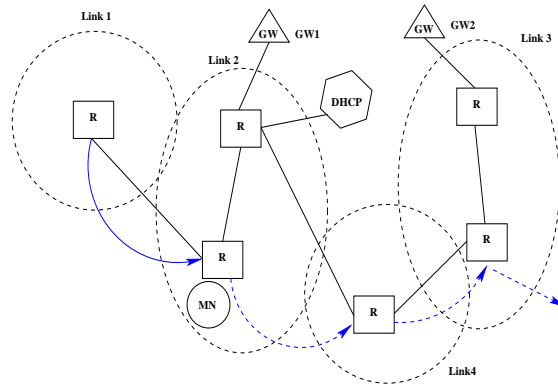


Messaging NSLP

- Thank You for listening
- Questions or comments?



NSLP for End-host mobility



Outline of presentation

- Overview
- Design criteria
- Related work
- Protocol messages
- Example sequence
- Phases of development and future work
- Questions



Overview

- Initial version functionality
 - Provide “best effort” mobility for unmodified clients
 - Edge access routers act on behalf of their clients
 - Host and source host routes
 - Proactive information exchange between mesh routers
- Clients
 - Arrive in network: DHCP
 - Move: send data or DHCP
- Mesh routers running NSIS
 - Detect clients
 - Set routes, monitor routes
 - Proactive information propagation



Design criteria

- Any routing protocol between ARs
 - MANET, ad hoc, regular IGP
 - GW known or a route 'upstream' available
 - Routing handles topology changes between ARs
- Address management through DHCP
 - Centralized or distributed
- Multicast available (for multicast GIST)
- IPv4 or IPv6; all or part of the nodes



Related work

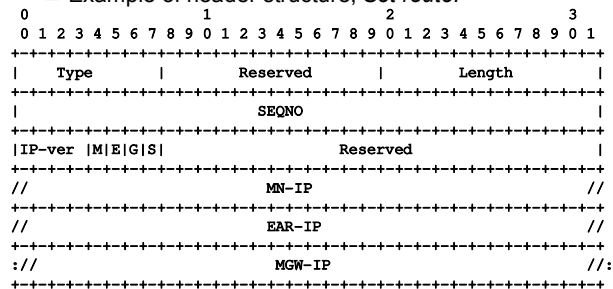
- MIP (Perkins RFC3344, Johnson et. al RFC3775)
- BCMP (Boukis et. al, IEEE globecom 2003)
- CellularIP Valko, ACM CCR 1999)
- HMIP (Soliman et. al RFC4140)
- NetLMM WG (ProxyMIP?)



Protocol messages

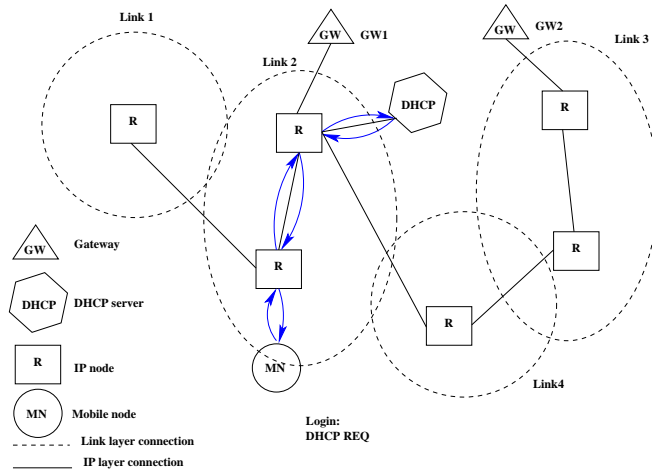
- Info
- Info refresh
- Set route
- Route reply
- Route refresh

■ Example of header structure, *Set route*:

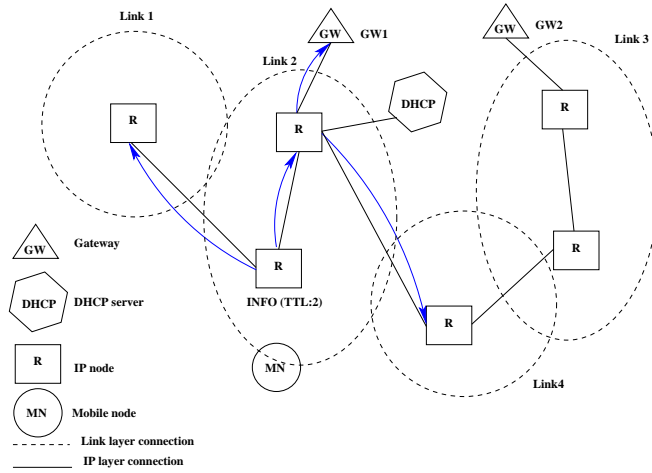




Example sequence: Login

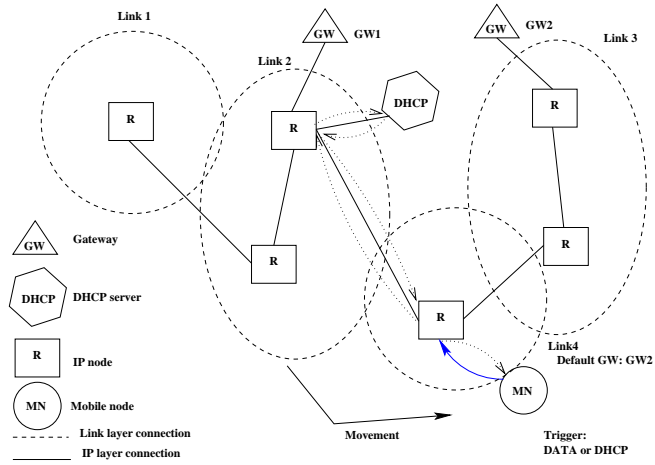


Example sequence: Info

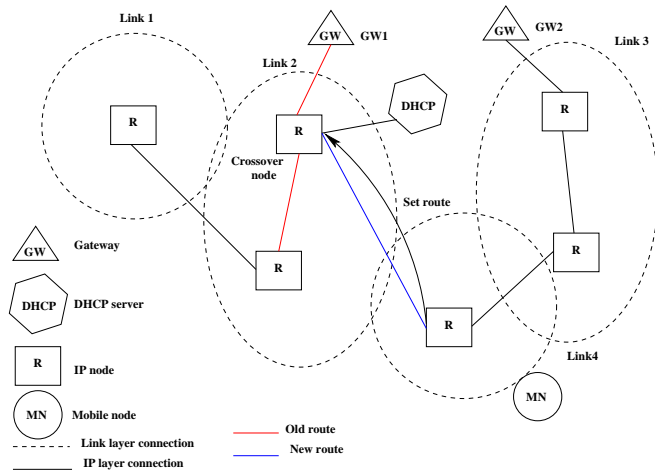




Example sequence: Handover

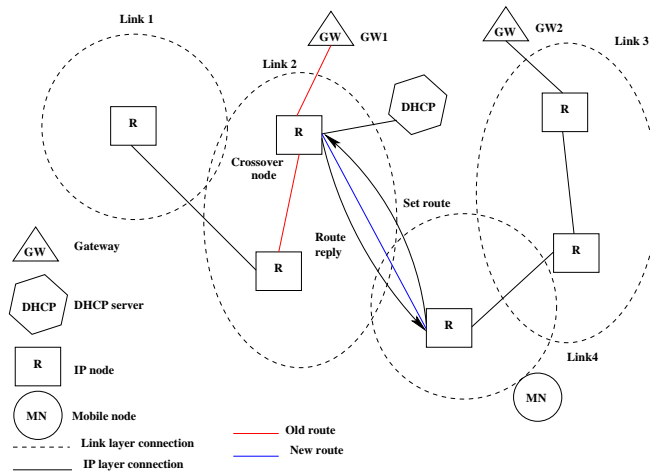


Example sequence: Set route





Example sequence: Route reply



Example sequence

- Terminating at the crossover node optimizes:
 - Delay in setting route
 - Space on routers, only routers that need them have host routes set
- CRN compares *Set route* to Mobile Node Info and detects different peer



Phases of development and future work 1/2

- Three phases: basic, basic plus and extended
- Basic
 - Like described above, only a baseline
 - Handovers not very smooth
 - No route optimization
 - Lacks persistence if information has not propagated
- Basic plus
 - Forcefully change IP of client to e.g. optimize routes
 - Reactive exchange of client information



Phases of development and future work 2/2

- Extended
 - The client participates in the signaling, NSLP on host
 - Co-operative handovers, planned change of IP
 - Session authentication
 - CARD + CTP functionality
 - Utilize other metrics for handover, available bandwidth and battery etc.
- Currently working on
 - Implementation
 - My MSc thesis on the basic phase
 - More detailed presentations for future workshops



Questions

- Thank You for listening
- Any questions or comments?