

On the Interpolation of Product-Based Message Passing Heuristics for SAT

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Outline

- 1 Goals
- 2 Message Passing
 - Message Passing on a conceptual level
 - Product-based MP heuristics
- 3 Interpolation and ISI
 - Interpolation
 - Indirect Structural Interpolation (ISI)
 - The product-based MP Hierarchy
- 4 $\rho\sigma\text{PMP}^i$
- 5 Conclusions

Goals

- 1 Provide better access to MP for the SAT community.
 - Provide a consistent notational frame to explain all currently available MP heuristics.
 - Explain the functioning of all these heuristics.
 - Explain their respective strengths and weaknesses.
 - Explain where they differ.
- 2 Extend our knowledge about MP.
 - Provide more general/flexible MP heuristics.
 - Integrate MP into a CDCL solver (used to initialize VSIDS and phase-saving).

Message Passing on a conceptual level (1)

- Message Passing (MP) is a class of algorithms
- $H \in \text{MP}$ can be understood as variable and value ordering heuristics in the context of SAT
- The main goal of H is to provide *biases* for all variables of a CNF F
- $\forall v \in \mathcal{V} : \beta_H(v) \in [-1.0, 1.0]$
- The biases can be used to guide search (CDCL or SLS)

Message Passing on a conceptual level (1)

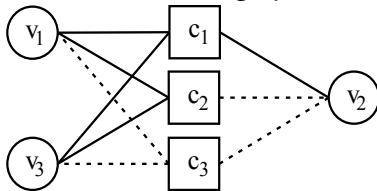
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- $\forall v \in \mathcal{V} : \beta_H(v) \in [-1.0, 1.0]$
- The biases can be used to guide search (CDCL or SLS)
- Given the formula F , what does H do to compute the biases?

Message Passing on a conceptual level (2)

Example

$$F = (v_1 \vee v_2 \vee v_3) \wedge (v_1 \vee \bar{v}_2 \vee v_3) \wedge (\bar{v}_1 \vee \bar{v}_2 \vee \bar{v}_3)$$

It is helpful to understand F as a *factor graph*.

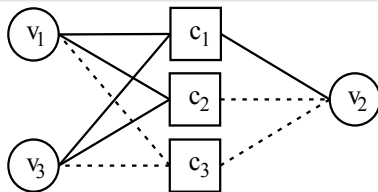


- Undirected, bipartite graph
- Two types of nodes (variable nodes (circles), clause nodes (squares))
- Two types of edges (positive edges (solid), negative edges (dashed))
- Edges constitute literal occurrences

Message Passing on a conceptual level (3)

Example

$$F = (v_1 \vee v_2 \vee v_3) \wedge (v_1 \vee \bar{v}_2 \vee v_3) \wedge (\bar{v}_1 \vee \bar{v}_2 \vee \bar{v}_3)$$



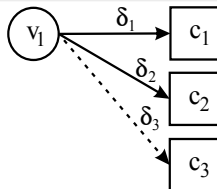
- H sends around messages along the edges.
- Assume variable v is contained in clause c as literal l

Two types of messages.

Message Passing on a conceptual level (4)

Example

$$F = (v_1 \vee v_2 \vee v_3) \wedge (v_1 \vee \bar{v}_2 \vee v_3) \wedge (\bar{v}_1 \vee \bar{v}_2 \vee \bar{v}_3)$$



1. Disrespect Messages (from variable nodes towards clause nodes):

- $\delta_H(l, c) \in [0.0, 1.0]$
- The chance that l will *not* satisfy c

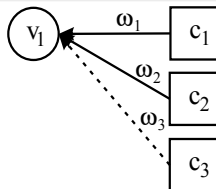
Intuitive meaning of $\delta_H(l, c) \approx 1.0$:

Variable v tells clause c that it cannot satisfy it.

Message Passing on a conceptual level (5)

Example

$$F = (v_1 \vee v_2 \vee v_3) \wedge (v_1 \vee \bar{v}_2 \vee v_3) \wedge (\bar{v}_1 \vee \bar{v}_2 \vee \bar{v}_3)$$



2. Warning Messages (from clause nodes towards variable nodes):

- $\omega_H(c, v) \in [0.0, 1.0]$
- The chance that *no other* literal in c can satisfy c

Intuitive meaning of $\omega_H(c, v) \approx 1.0$:

Clause c is telling variable v , that it needs it to be satisfied.

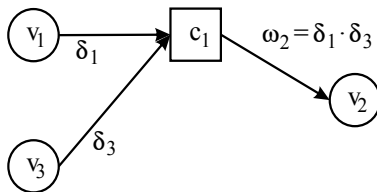
Message Passing on a conceptual level (6)

Example

$$F = (v_1 \vee v_2 \vee v_3) \wedge (v_1 \vee \bar{v}_2 \vee v_3) \wedge (\bar{v}_1 \vee \bar{v}_2 \vee \bar{v}_3)$$

For all product-based MP heuristics, the waring message is computed by

$$\omega_H(c, v) = \prod_{l \in c \setminus \{v, \bar{v}\}} \delta_H(l, c)$$



Message Passing on a conceptual level (7)

For all product-based MP heuristics, the *cavity freedom values* are computed by

$$[0.0, 1.0] \ni S_{\mathbf{H}}(l, c) = \begin{cases} \prod_{d \in C_v^-} [1 - \omega_{\mathbf{H}}(d, v)], l = v \\ \prod_{d \in C_v^+} [1 - \omega_{\mathbf{H}}(d, v)], l = \bar{v} \end{cases}$$

Intuitive meaning:

How happy are the other clauses if l satisfies c ?

$$[0.0, 1.0] \ni U_{\mathbf{H}}(l, c) = \begin{cases} \prod_{d \in C_v^+ \setminus \{c\}} [1 - \omega_{\mathbf{H}}(d, v)], l = v \\ \prod_{d \in C_v^- \setminus \{c\}} [1 - \omega_{\mathbf{H}}(d, v)], l = \bar{v} \end{cases}$$

Intuitive meaning:

How happy are the other clauses if l does not satisfy c ?

Message Passing on a conceptual level (8)

In summary:

- Computed δ_H values allow us to compute the ω_H values
- Computed ω_H values allow us to compute the S_H, U_H values

However:

- H will not send around messages arbitrarily
- H performs *clause updates* $\forall c \in F$
- The ordering of the clauses in which they receive updates is determined by a *random clause permutation* $\pi \in \mathcal{S}_m$

Message Passing on a conceptual level (9)

Following $\pi \in \mathcal{S}_m$, each clause is updated exactly once.

Basically, a clause update for clause c consists of three steps.

- ① Compute $\forall l \in c : \delta_H(l, c)$
- ② Using the δ , compute $\forall v \in c : \omega_H(c, v)$
- ③ Using the ω , compute $\forall l \in c : S_H(l, c), U_H(l, c)$

Where do the δ values come from in order to compute a clause update?

Message Passing on a conceptual level (10)

We need the terms of *iteration* and *cycle* to explain that.

- Doing the clause updates for all clauses exactly once is called an *iteration*.
- A *cycle* is a finite tuple of iterations.
- Iterations and cycles capture the notion of *passing time* while H performs its computations.
- An iteration is a single point in time, a cycle is a time-frame.

We denote the specific values computed in iteration z of cycle y with

- ${}^y_z\delta_H(l, c)$
- ${}^y_z\omega_H(c, v)$
- ${}^y_zS_H(l, c)$
- ${}^y_zU_H(l, c)$

Message Passing on a conceptual level (11)

Again, in order to compute the clause update for iteration z in cycle y

- ① Compute $\forall l \in c : {}^y_z\delta_H(l, c)$
- ② Using the ${}^y_z\delta_H(l, c)$, compute $\forall v \in c : {}^y_z\omega_H(c, v)$
- ③ Using the ${}^y_z\omega_H(c, v)$, compute $\forall l \in c : {}^y_zS_H(l, c), {}^y_zU_H(l, c)$

Again, where do the ${}^y_z\delta_H(l, c)$ values come from in order to compute a clause update?

Message Passing on a conceptual level (12)

The initialization for cycle y happens in iteration $z = 0$.

- $\forall c \in F : \forall l \in c$: initialize randomly with ${}^y_0\delta_H(l, c) \in_R (0.0, 1.0)$
- The values for ${}^y_0\omega_H(c, v), {}^y_0S_H(l, c), {}^y_0U_H(l, c)$ then directly follow with the definitions.

The clause updates for cycle y and iteration $z > 0$ are defined recursive.

- Rely on ${}_{z-1}^yS_H(l, c), {}_{z-1}^yU_H(l, c)$ in order to compute ${}^y_z\delta_H(l, c)$.

How exactly is ${}^y_z\delta_H(l, c)$ computed using these values?

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How exactly is ${}^y_z\delta_H(l, c)$ computed using these values?

- *This must be defined by H!*

Message Passing on a conceptual level (13)

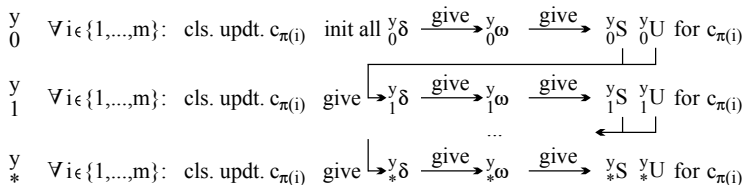
For Belief Propagation (BP) this is defined as

$$\bullet \quad {}_z^y\delta_{\text{BP}}(l, c) = \frac{{}_{z-1}^yU_{\text{BP}}(l, c)}{{}_{z-1}^yU_{\text{BP}}(l, c) + {}_{z-1}^yS_{\text{BP}}(l, c)} \left(= \frac{U}{U + S} \right)$$

Message Passing on a conceptual level (14)

We now know

- ... how cycles start.
- ... how the iterations are done.

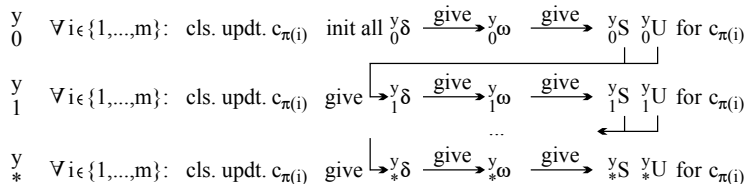


We do not know

Message Passing on a conceptual level (14)

We now know

- ... how cycles start.
- ... how the iterations are done.



We do not know

- ... how a cycle terminates.
- What we need is an *abort condition*.

Message Passing on a conceptual level (15)

The abort conditions for a product-based MP heuristics is defined as

- $\forall c \in F : \forall v \in c : | \overset{y}{z} \omega_{\text{H}}(c, v) - \overset{y}{z-1} \omega_{\text{H}}(c, v) | < \omega_{\text{max}}$
- In practice $\omega_{\text{max}} = 0.01$

The iteration of cycle y in which the abort condition holds is denoted $*$.

The messages

- $\overset{y}{*} \delta_{\text{H}}(l, c)$
- $\overset{y}{*} \omega_{\text{H}}(c, v)$

are called *equilibrium messages*.

The $\overset{y}{*} \omega_{\text{H}}(c, v)$ are used to compute the biases for cycle y .

Message Passing on a conceptual level (16)

Computing biases is done in three steps using the ${}^y\omega_{\mathbf{H}}(c, v)$.

- 1 Compute the *variable freedom* to be assigned to true (\mathcal{T}) or false (\mathcal{F})

$${}^y\mathcal{T}_{\mathbf{H}}(v) = \prod_{c \in C_v^-} [1 - {}^y\omega_{\mathbf{H}}(c, v)] \quad {}^y\mathcal{F}_{\mathbf{H}}(v) = \prod_{c \in C_v^+} [1 - {}^y\omega_{\mathbf{H}}(c, v)]$$

- 2 Compute *magnetization values* using \mathcal{T} and \mathcal{F}

$${}^y\mu_{\mathbf{H}}^+(v), {}^y\mu_{\mathbf{H}}^-(v), {}^y\mu_{\mathbf{H}}^{\pm}(v) \in [0.0, 1.0]$$

These give ${}^y\mu_{\mathbf{H}}(v) = {}^y\mu_{\mathbf{H}}^+(v) + {}^y\mu_{\mathbf{H}}^-(v) + {}^y\mu_{\mathbf{H}}^{\pm}(v)$

- 3 Compute the biases

$${}^y\beta_{\mathbf{H}}^+(v) = \frac{{}^y\mu_{\mathbf{H}}^+(v)}{{}^y\mu_{\mathbf{H}}(v)} \quad {}^y\beta_{\mathbf{H}}^-(v) = \frac{{}^y\mu_{\mathbf{H}}^-(v)}{{}^y\mu_{\mathbf{H}}(v)} \quad {}^y\beta_{\mathbf{H}}(v) = {}^y\beta_{\mathbf{H}}^+(v) - {}^y\beta_{\mathbf{H}}^-(v)$$

Message Passing on a conceptual level (17)

Where do the ${}^y\mu_{\mathcal{H}}^+(v)$, ${}^y\mu_{\mathcal{H}}^-(v)$, ${}^y\mu_{\mathcal{H}}^{\pm}(v) \in [0.0, 1.0]$ come from?
Again, this must be defined by $H!$

For Belief Propagation (BP), this is defined as

- ${}^y\mu_{\text{BP}}^+(v) = {}^y\mathcal{T}_{\text{BP}}(v)$
- ${}^y\mu_{\text{BP}}^-(v) = {}^y\mathcal{F}_{\text{BP}}(v)$
- ${}^y\mu_{\text{BP}}^{\pm}(v) = 0$

Therefore, ${}^y\mu_{\text{BP}}(v) = {}^y\mathcal{T}_{\text{BP}}(v) + {}^y\mathcal{F}_{\text{BP}}(v)$.

Finally, for BP, it is ${}^y\beta_{\text{BP}}(v) = \frac{{}^y\mathcal{T}_{\text{BP}}(v) - {}^y\mathcal{F}_{\text{BP}}(v)}{{}^y\mathcal{T}_{\text{BP}}(v) + {}^y\mathcal{F}_{\text{BP}}(v)}$

Product-based MP heuristics (1)

Well known product-based MP heuristics.

EMBPG

EMSPG

Level 0

BP

SP

EMBPG

EMSPG

BP

SP

- All the basic MP heuristics have different strengths and weaknesses.
- Introducing MP into a solver to guide its search is problematic.
- The necessity to choose basically means: However you choose, you choose wrong!

Product-based MP heuristics (2)

Increase the flexibility of MP heuristics in order to overcome the "robustness problem".

How to create a more flexible MP heuristic?

Product-based MP heuristics (2)

Increase the flexibility of MP heuristics in order to overcome the "robustness problem".

How to create a more flexible MP heuristic?

- Interpolation!

Interpolation (1)

What is it, that needs to be achieved in order to create an interpolation?
 Given two product-based MP heuristics H_1 and H_2 , we want an interpolation ρH^i , s.t.

- interpolation parameter $\rho \in [0.0, 1.0]$
- Setting $\rho = 0$ will make ρH^i mimic H_1 , i.e. $\beta_{H_1}(v) = \beta_{\rho H^i}(v, 0)$
- Setting $\rho = 1$ will make ρH^i mimic H_2 , i.e. $\beta_{H_2}(v) = \beta_{\rho H^i}(v, 1)$
- Setting $\rho \in (0.0, 1.0)$ results in a gradual adaption between H_1, H_2
 - gradually adapt the convergence behavior
 - gradually adapt the carefulness to present biases

Interpolation (2)

Equations used in all product-based MP heuristics.

During Iterations

- Disrespect message ${}^y\delta_{\mathbf{H}}(l, c)$
- Warning message ${}^y\omega_{\mathbf{H}}(l, c)$
- Literal cavity freedom values ${}^yS_{\mathbf{H}}(l, c), {}^yU_{\mathbf{H}}(l, c)$

After convergence, provided ${}^y\omega_{\mathbf{H}}(l, c)$

- Variable freedom ${}^y\mathcal{T}_{\mathbf{H}}(v), {}^y\mathcal{F}_{\mathbf{H}}(v)$
- Variable magnetization ${}^y\mu_{\mathbf{H}}^+(v), {}^y\mu_{\mathbf{H}}^-(v), {}^y\mu_{\mathbf{H}}^{\pm}(v), {}^y\mu_{\mathbf{H}}(v)$
- Variable bias ${}^y\beta_{\mathbf{H}}^+(v), {}^y\beta_{\mathbf{H}}^-(v), {}^y\beta_{\mathbf{H}}(v)$

Interpolation (3)

Equations that must be **defined by H** itself.

During Iterations

- Disrespect message ${}_z^y\delta_{\mathbf{H}}(l, c)$
- Warning message ${}_z^y\omega_{\mathbf{H}}(l, c)$
- Literal cavity freedom values ${}_z^yS_{\mathbf{H}}(l, c), {}_z^yU_{\mathbf{H}}(l, c)$

After convergence, provided ${}_z^y\omega_{\mathbf{H}}(l, c)$

- Variable freedom ${}_z^y\mathcal{T}_{\mathbf{H}}(v), {}_z^y\mathcal{F}_{\mathbf{H}}(v)$
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- Variable bias ${}_z^y\beta_{\mathbf{H}}^+(v), {}_z^y\beta_{\mathbf{H}}^-(v), {}_z^y\beta_{\mathbf{H}}(v)$

Must be defined for the **interpolation**.

ISI (1)

ISI is a technique to derive ρH^i given H_1 and H_2 .

It uses an interpolation parameter $\rho \in [0.0, 1.0]$.

It derives

- $y_z \delta_{\rho H}^i(l, c, \rho), y \mu_{\rho H}^{i+}(v, \rho), y \mu_{\rho H}^{i-}(v, \rho), y \mu_{\rho H}^{i\pm}(v, \rho)$

given

- $y_z \delta_{H_1}(l, c), y \mu_{H_1}^+(v), y \mu_{H_1}^-(v), y \mu_{H_1}^{\pm}(v)$
- $y_z \delta_{H_2}(l, c), y \mu_{H_2}^+(v), y \mu_{H_2}^-(v), y \mu_{H_2}^{\pm}(v)$

How exactly does it work? Exemplary explanation.

Assume we want to

- interpolate BP and SP
- using interpolation parameter $\rho \in [0.0, 1.0]$
- in order to derive the interpolation ρSP^i

ISI (2)

Step 1. derives ${}_z^y\delta_{\rho\text{SP}}^i(l, c, \rho)$ using

- ${}_z^y\delta_{\text{BP}}(l, c) = \frac{U}{U+S}$
- ${}_z^y\delta_{\text{SP}}(l, c) = \frac{U(1-S)}{U(1-S)+S}$

Linearly interpolate!

Numerator:

$$(1 - \rho)\{U\} + \rho\{U(1 - S)\} = \dots = U(1 - \rho S)$$

Denominator:

$$(1 - \rho)\{U + S\} + \rho\{U(1 - S) + S\} = \dots = U(1 - \rho S) + S$$

Combine:

$${}_z^y\delta_{\rho\text{SP}}^i(l, c, \rho) = \frac{U(1 - \rho S)}{U(1 - \rho S) + S}$$

ISI (3)

Step 2. derives ${}^y\mu_{\rho\text{SP}}^{i+}(v, \rho)$ using

- ${}^y\mu_{\text{BP}}^+(v) = {}^y\mathcal{T}_{\text{BP}}(v)$ ($= \mathcal{T}$)
- ${}^y\mu_{\text{SP}}^+(v) = {}^y\mathcal{T}_{\text{SP}}(v)(1 - {}^y\mathcal{F}_{\text{SP}}(v))$ ($= \mathcal{T}(1 - \mathcal{F})$)

Linearly interpolate!

$$(1 - \rho)\{\mathcal{T}\} + \rho\{\mathcal{T}(1 - \mathcal{F})\} = \dots = \mathcal{T}(1 - \rho\mathcal{F}) = {}^y\mu_{\rho\text{SP}}^{i+}(v, \rho)$$

Step 3. derives ${}^y\mu_{\rho\text{SP}}^{i-}(v, \rho)$ in a similar way.

Step 4. derives ${}^y\mu_{\rho\text{SP}}^{i\pm}(v, \rho)$ in a similar way.

In the end, all four defining functions for ρSP^i have been derived.

The product-based MP Hierarchy (1)

The basic product-based MP heuristics.

EMBPG

EMSPG

BP

SP

Level 0

BP

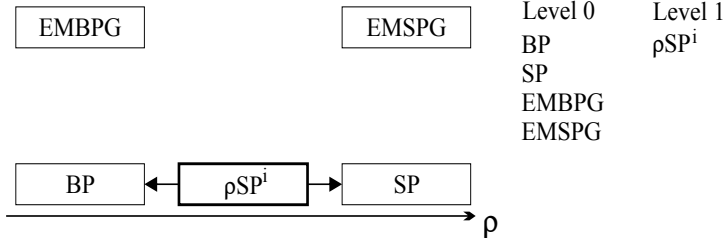
SP

EMBPG

EMSPG

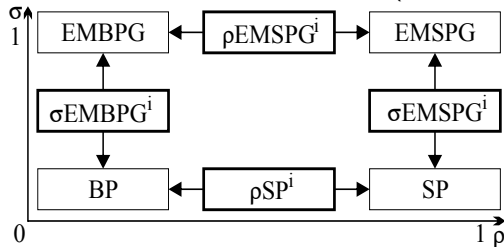
The product-based MP Hierarchy (2)

The first level of interpolations (applying ISI once).



The product-based MP Hierarchy (3)

The first level of interpolations (applying ISI once).



Level 0

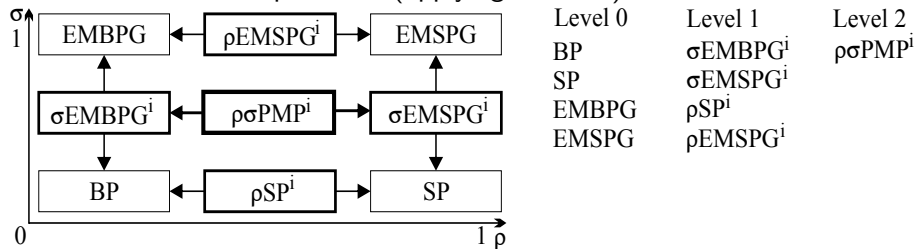
BP
SP
EMBPG
EMSPG

Level 1

σEMBPG^i
 σEMSPG^i
 ρSP^i
 ρEMSPG^i

The product-based MP Hierarchy (4)

The first level of interpolations (applying ISI once).



$\rho\sigma\text{PMP}^i$ (1)

Why is $\rho\sigma\text{PMP}^i$ so special?

- It is the most general product-based MP heuristic.
- It can mimic the behavior of all others.
- It can provide MP behavior that cannot be achieved by any other heuristic.

Each point in the parameter plane $(\rho, \sigma) \in [0.0, 1.0]^2$ characterizes a specific MP behavior.

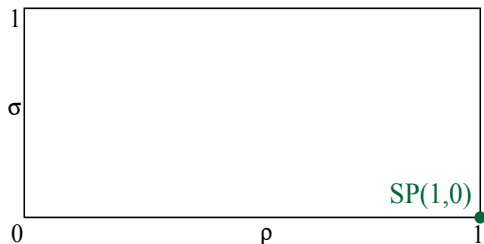


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Best behavior given?

Can use:

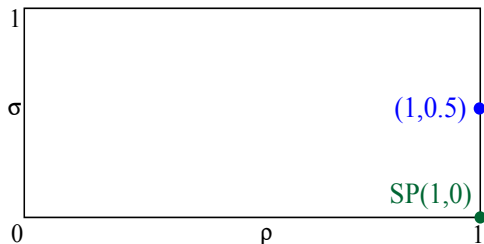
SP , ρSP^i , σEMSPG^i , $\rho\sigma\text{PMP}^i$

$\rho\sigma\text{PMP}^i$ (3)

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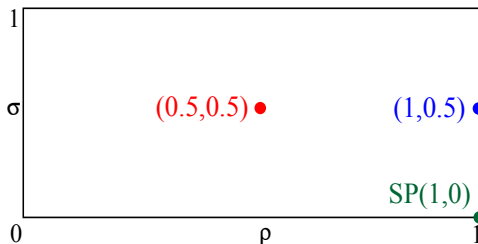
σEMSPG^i , $\rho\sigma\text{PMP}^i$

$\rho\sigma\text{PMP}^i$ (4)

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SP , ρSP^i , σEMSPG^i , $\rho\sigma\text{PMP}^i$

Can use:

σEMSPG^i , $\rho\sigma\text{PMP}^i$

Can use:

$\rho\sigma\text{PMP}^i$

$\rho\sigma\text{PMP}^i$ (5)

Why is that good in order to introduce MP into a solver?

- This circumvents the need to choose from all the available MP heuristics.
- The interpolation parameters ρ, σ can be tuned automatically for each class of formulas.

In the context of a CDCL search:

- 1 Use $\rho\sigma\text{PMP}^i$ to compute biases.
- 2 Use a specifically tuned MP behavior for the formula class.
- 3 Use the biases to initialize VSIDS and phase-saving.

Empirical results from parameter tuning

Benchmark	S/U	Solver Performance					
		DimetheusJW		DimetheusMP			
		%	PAR10	%	PAR10	ρ	σ
battleship	S	47.4	10627.2	89.5	2130.1	0.5002	0.0025
battleship	U	55.6	8919.7	55.6	8890.4	0.4463	1.0000
em-all	S	75.0	5263.7	100.0	75.4	0.8606	0.1295
em-compact	S	0.0	20000.0	37.5	12728.5	0.9229	0.7946
em-explicit	S	75.0	5473.3	100.0	157.1	0.2932	0.2698
em-fbcolors	S	12.5	17723.3	37.5	12662.9	0.0000	0.1731
grid-pebbling	S	100.0	16.5	100.0	8.0	0.9931	0.3890
grid-pebbling	U	88.9	2226.9	100.0	4.7	0.5884	0.0035
sgen1	S	16.7	16677.7	27.8	14460.9	0.0937	0.6563
k3-r4.200	S	0.0	20000.0	100.0	22.7	0.9929	0.0004
k3-r4.237	S	0.0	20000.0	75.0	5026.8	0.9961	0.0000
k4-r9.000	S	0.0	20000.0	100.0	10.0	0.8592	0.0000
k4-r9.526	S	0.0	20000.0	100.0	5.2	0.9530	0.0000

Conclusions

- 1 Provided better access to MP for the SAT community.
 - We provided a unified and consistent notational frame to explain all currently available MP heuristics.
 - We explained the functioning of all these heuristics.
 - We explained their respective strengths and weaknesses.
 - We explained where they differ.
- 2 Extend our knowledge about MP.
 - We provided a hierarchy of generality regarding product-based MP heuristics.
 - We clarified what an interpolation is and how they are derived.
 - Integrated MP into a CDCL solver (used to initialize VSIDS and phase-saving) to get more empirical insight.

Thanks you for your attention!

You can send disrespect messages and questions to
oliver@gableske.net

Thank you for your attention.

Check the paper

O. Gableske

*On the Interpolation between Product-Based
Message Passing Heuristics for SAT*

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The difference between BP and SP

With $\rho, S, U, \mathcal{T}, \mathcal{F} \in [0.0, 1.0]$

Disrespect messages:

- ${}^y_z\delta_{\text{BP}}(l, c) = \frac{U}{U + S}$ ${}^y_z\delta_{\text{SP}}(l, c) = \frac{U(1 - S)}{U(1 - S) + S}$
- ${}^y_z\delta_{\rho\text{SP}}^i(l, c, \rho) = \frac{U(1 - \rho S)}{U(1 - \rho S) + S}$

Bias computations:

- ${}^y\beta_{\text{BP}}(v) = \frac{\mathcal{T} - \mathcal{F}}{\mathcal{T} + \mathcal{F}}$ ${}^y\beta_{\text{SP}}(v) = \frac{\mathcal{T} - \mathcal{F}}{\mathcal{T} + \mathcal{F} - \mathcal{T}\mathcal{F}}$
- ${}^y\beta_{\rho\text{SP}}^i(v, \rho) = \frac{\mathcal{T} - \mathcal{F}}{\mathcal{T} + \mathcal{F} - \rho\mathcal{T}\mathcal{F}}$