## 58093 String Processing Algorithms (Autumn 2011)

Exercises 3 (15 November)

- 1. Describe how to modify the LSD radix sort algorithm to handle strings of varying lengths. The time complexity should be the one given in Theorem 1.14.
- Ω(dp(R)) is a lower bound for string sorting for any algorithm if characters can be accessed only one at a time. However, for a small alphabet, it is possible to pack several characters into one machine word. Then multiple characters can be accessed simultaneously and treated as if they were a single *super-character*. For example, the string abbaba over the alphabet Σ = {a, b} can be thought of as the string (ab, ba, ab) over the alphabet Σ<sup>2</sup>. Algorithms taking advantage of this are called *super-alphabet* algorithms.

Develop a super-alphabet version of MSD radix sort. What is the time complexity?

- 3. Use the lcp comparison technique to modify the standard insertion sort algorithm so that it sorts strings in  $\mathcal{O}(dp(\mathcal{R}) + n^2)$  time.
- 4. Let  $\mathcal{R} = \{$ manne, manu, minna, salla, saul, sauli, vihtori $\}$ .
  - (a) Give the compact trie of  $\mathcal{R}$ .
  - (b) Give the balanced compact ternary trie of  $\mathcal{R}$ .
- 5. Show that the number of nodes in a trie  $trie(\mathcal{R})$  is exactly  $||\mathcal{R}|| lcp(\mathcal{R}) + 1$ , where  $||\mathcal{R}||$  is the total length of the strings in  $\mathcal{R}$  and  $lcp(\mathcal{R})$  is as defined in Exercise 2.5. *Hint:* Consider the construction of  $trie(\mathcal{R})$  using Algorithm 2.2.
- 6. Give an example showing that the worst case time complexity of string binary search without precomputed lcp information is  $\Omega(m \log n)$ .
- 7. Define

$$MLCP[mid] = \max\{LLCP[mid], RLCP[mid]\}$$
$$D[mid] = \begin{cases} 0 & \text{if } MLCP[mid] = LLCP[mid] \\ 1 & \text{otherwise} \end{cases}$$

Show that, if we store the arrays MLCP and D instead of LLCP and RLCP, we can compute LLCP[mid] and RLCP[mid] when needed during the string binary search.