

<p>Process 1</p> <p>Outgoing channels</p> <p>2 sent 1,2,3,4,5,6</p> <p>3 sent 1,2,3,4,5,6</p> <p>Incoming channels</p>	<p>Process 3</p> <p>Outgoing channels</p> <p>2 sent 1,2,3,4,5,6,7,8</p> <p>Incoming channels</p> <p>1 received 1,2,3 stored 4,5,6</p> <p>2 received 1,2,3 stored 4</p> <p>4 received 1,2,3</p>
<p>Process 2</p> <p>Outgoing channels</p> <p>3 sent 1,2,3,4</p> <p>4 sent 1,2,3,4</p> <p>Incoming channels</p> <p>1 received 1,2,3,4 stored 5,6</p> <p>3 received 1,2,3,4,5,6,7,8</p>	<p>Process 4</p> <p>Outgoing channels</p> <p>3 sent 1,2,3</p> <p>Incoming channels</p> <p>2 received 1,2 stored 3,4</p>

Figure 14.6 An Example of a Snapshot

```

if (!token_present)
{
    clock++;
    broadcast (Request, clock, i);
    wait (access, token);
    token_present = true;
}

token_held = true;
<critical section>;

token[i] = clock;
token_held = false;
for (int j = i + 1; j < n; j++)
{
    if (request(j) > token[j] && token_present)
    {
        token_present = false;
        send (access, token[j]);
    }
}
for (j = 1; j <= i-1; j++)
{
    if (request(j) > token[j] && token_present)
    {
        token_present = false;
        send(access, token[j]);
    }
}

```

(a) First Part

```

if (received (Request, k, j))
{
    request (j) = max(request(j), k);
    if (token_present && !token_held)
        <text of postlude>;
}

```

(b) Second Part

Notation

send (j, access, token)	end message of type access, with token, by process j
broadcast (request, clock, i)	send message from process i of type request, with time-stamp clock, to all other processes
received (request, t, j)	receive message from process j of type request, with time-stamp t

Figure 14.11 Token-Passing Algorithm (for process P_i)

```
if (e(T2) < e(T1))  
    halt_T2 ('wait');  
else  
    kill_T2 ('die');
```

(a) Wait-die method

```
if (e(T2) < e(T1))  
    kill_T1 ('wound');  
else  
    halt_T2 ('wait');
```

(b) Wound-wait method

Figure 14.13 Deadlock Prevention Methods