Lesson 2

Concurrency at Programming Language Level

Ch 2 [BenA 06]

Abstraction
Pseudo-language
BACI
Ada, Java, etc.

Levels of Abstraction

- Granularity of operations
  - Invoke a library module
  - Statement in high level programming language
  - Instruction in machine language
- Atomic statement
  - Anything that we can guarantee to be atomic
    - Executed completely “at once”
    - Always the same correct atomic result
    - Result does not depend on anybody else
  - Can be at any granularity
  - Can *trust* on that atomicity

Atomic Statement

- Atomicity guaranteed somehow
  - Machine instruction: HW
    - Memory bus transaction
  - Programming language statement, set of statements, or set of machine instructions
    - SW
      - Manually coded
      - Disable interrupts
      - OS synchronization primitives
  - Library module
    - SW
      - Manually coded inside
      - <u>Provided automatically</u> to the user by programming environment

Load R1, Y

Read mem(0x35FA8300)

?

-- start atomic
Load R1, Y
Sub R1, =1
Jpos R1, Here
-- end atomic

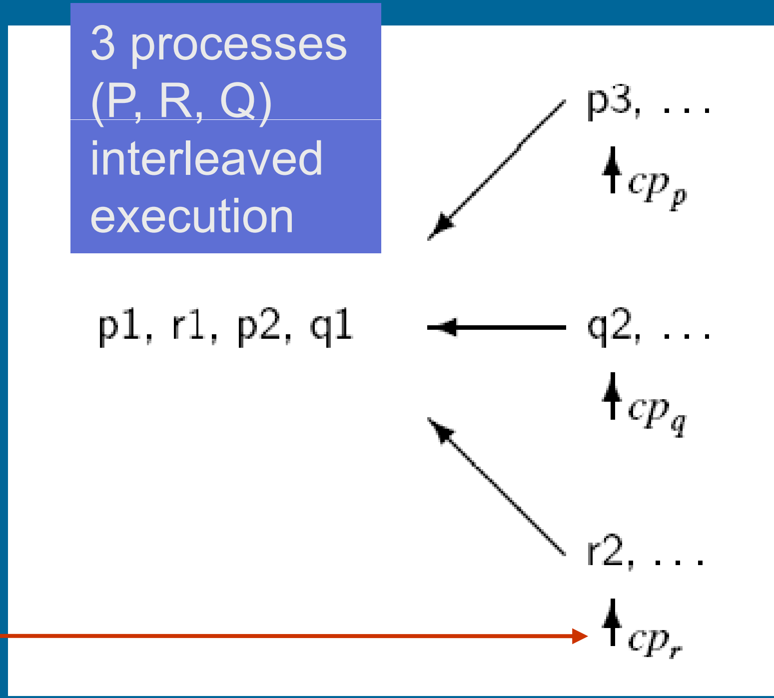
Monitors
Ch 7 [BenA 06]

Concurrent Program

- Sequential process
  - Successive atomic statements

P: p1 → p2 → p3 → p4 …

  - Control pointer (= program counter)



- Concurrent program
  - Finite set of sequential processes working for same goal
  - Arbitrary interleaving of atomic statements in different processes

P: p1 → p2

Q: q1 → q2

p1→q1→p2→q2,
p1→q1→q2→p2,
p1→p2→q1→q2,
q1→p1→q2→p2,
q1→p1→p2→q2,
q1→q2→p1→p2.

?

~~p1 → q2 → p2 → q1~~

Program State, Pseudo-language

pseudo-kieli

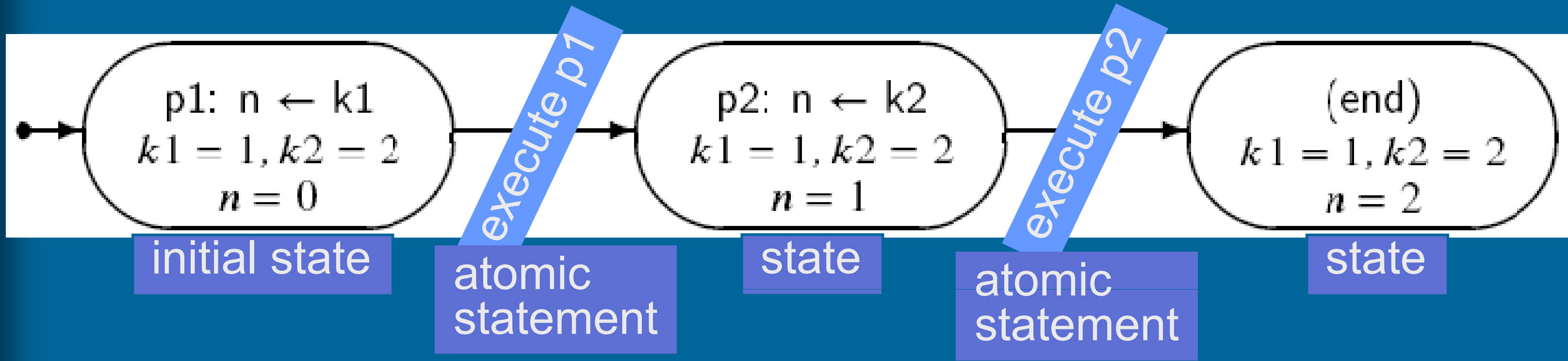
- Sequential program

Algorithm 2.2: Trivial sequential program
integer n ← 0
integer k1 ← 1
integer k2 ← 2
p1: n ← k1
p2: n ← k2

cp

state

- State
  - next statement to execute (cp, i.e., PC)
  - variable values

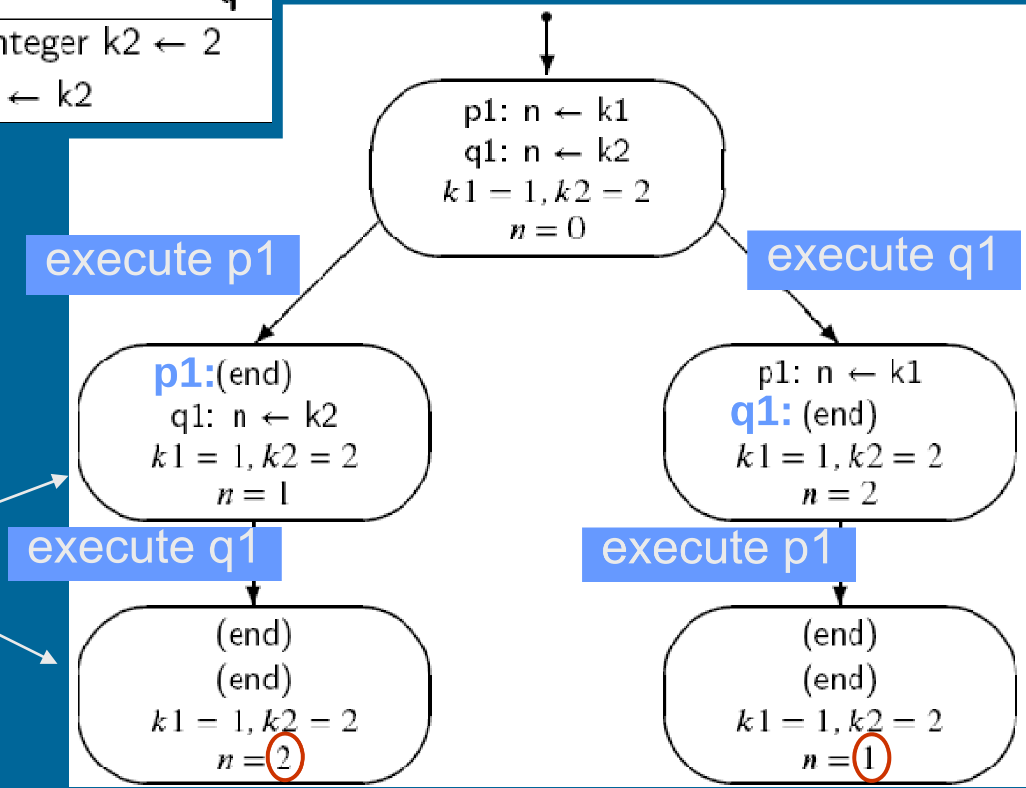


(Global) Program State

- Concurrent program

Algorithm 2.1: Trivial concurrent program	
integer n ← 0	
p	q
integer k1 ← 1	integer k2 ← 2
p1: n ← k1	q1: n ← k2

- Local state for each *process*:
  - cp
  - Variable values
    - Local & global
- Global state for *program*
  - *All* cp's
  - *All* local variables
  - *All* global variables



Possible Program States

- List of processes in program
  - List of values for each process
    - cp
    - value of each local/global/shared variable

p1: n ← k1
q1: n ← k2
$k1 = 1, k2 = 2$
$n = 0$

state: { { p1: n ← k1 **– process p**
k1 = 1 }
{ q1: n ← k2 **– process q**
k2 = 2 }
n = 0 **– shared variable**
}

- Nr of possible states can be (very) large
  - Not all states are reachable states!

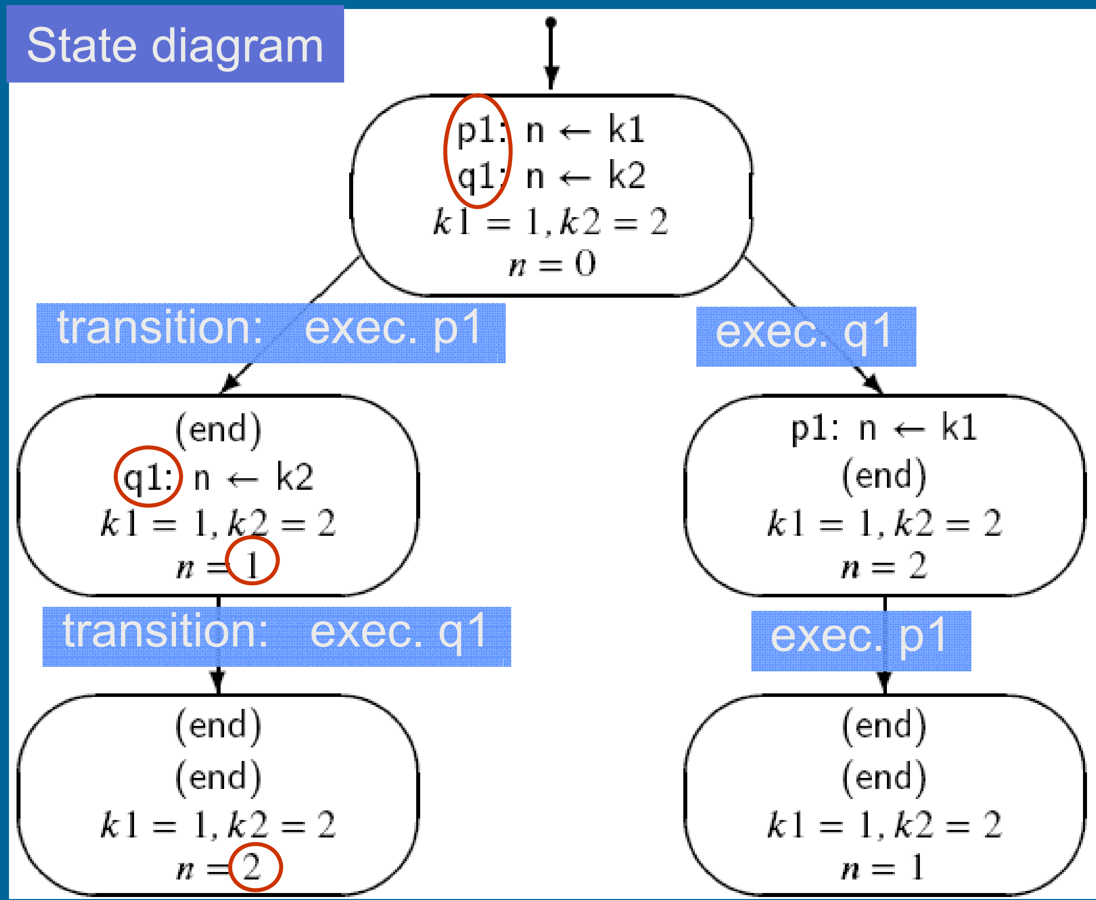
(saavutettavissa, saavutettava tila)

unreachable state: { { p1: n ← k1
k1 = *2* }
{ q1: n ← k2
k2 = *1* }
n = *3*
}

State Diagram and Scenarios

Process p	Process q	n	k1	k2
p1: n←k1	q1: n←k2	0	1	2
(end)	**q1: n←k2**	1	1	2
(end)	(end)	2	1	2

Scenario 1 (left side)



- Transitions from one possible state to another
  - Executed statement must be one of those in the 1st state
- State diagram for concurrent program
  - Contains all reachable states and transitions
  - All possible executions are included, they are all correct!

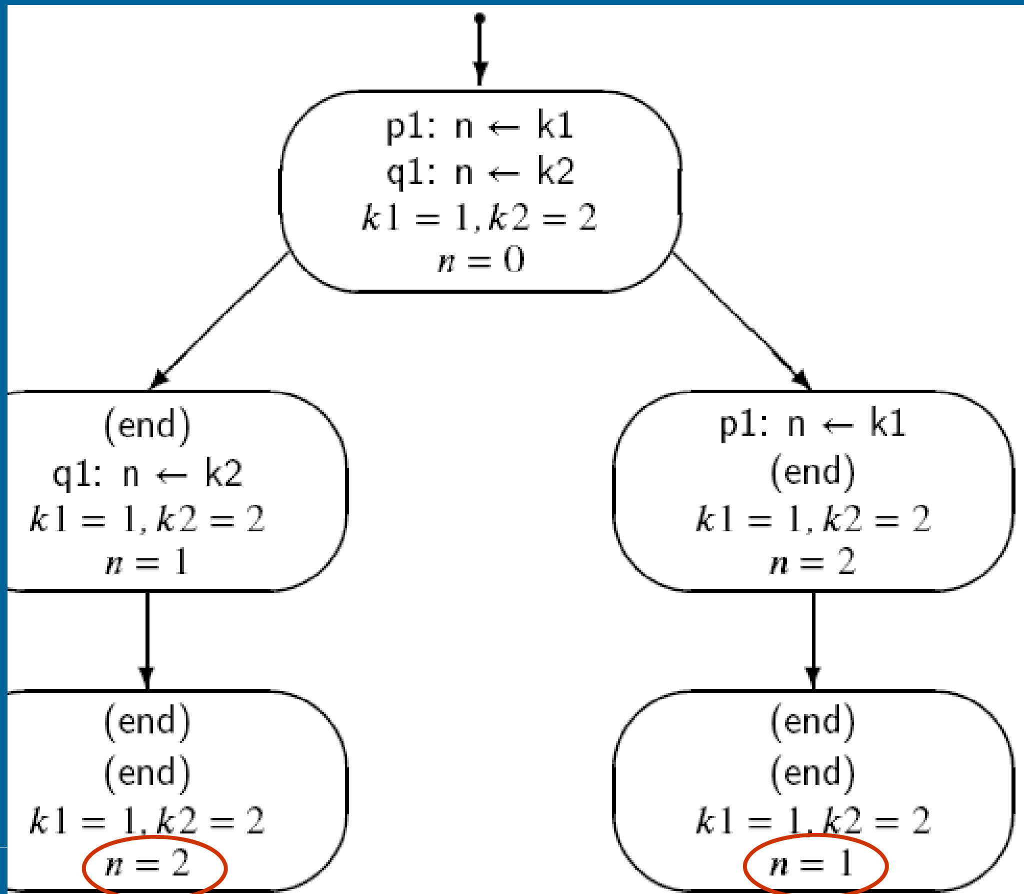
Atomic Statements

Algorithm 2.1: Trivial concurrent program	
integer n ← 0	
p	q
integer k1 ← 1	integer k2 ← 2
p1: n ← k1	q1: n ← k2

- Two scenar
  - Both corre
  - Different r

NO need to have the same result! Statements do the same, but overall result may be different. (see p. 19 [BenA 06])

- Atomic?
  - Assignme
  - Boolean e
  - Increment



Algorithm 2.3: Atomic assignment statements	
integer $n \leftarrow 0$	
p	**q**
p1: $n \leftarrow n + 1$	q1: $n \leftarrow n + 1$

- Two scenarios for execution
  - Both correct
  - Both have the same result

P first, and then Q

Process p	Process q	n
p1: $n \leftarrow n+1$	q1: $n \leftarrow n+1$	0
(end)	**q1: $n \leftarrow n+1$**	1
(end)	(end)	2

Q first, and then P

Process p	Process q	n
p1: $n \leftarrow n+1$	**q1: $n \leftarrow n+1$**	0
p1: $n \leftarrow n+1$	(end)	1
(end)	(end)	2

Algorithm 2.3: Atomic assignment statements	
integer n ← 0	
p	q
p1: n ← n + 1	q1: n ← n + 1

Same statements with smaller atomic granularity:

Algorithm 2.4: Assignment statements with one global reference	
integer n ← 0	
p	q
integer temp	integer temp
p1: temp ← n	q1: temp ← n
p2: n ← temp + 1	q2: n ← temp + 1

Too Small Atomic Granularity

Algorithm 2.4: Assignment statements with one global reference	
integer n ← 0	
p	**q**
integer temp	integer temp
p1: temp ← n	q1: temp ← n
p2: n ← temp + 1	q2: n ← temp + 1

- Scenario 1
  - OK

Process p	Process q	n	p.temp	q.temp
p1: temp←n	q1: temp←n	0	?	?
p2: n←temp+1	q1: temp←n	0	0	?
(end)	**q1: temp←n**	1	0	?
(end)	**q2: n←temp+1**	1	0	1
(end)	(end)	2	0	1

- Scenario 2
  - Bad result

Process p	Process q	n	p.temp	q.temp
p1: temp←n	q1: temp←n	0	?	?
p2: n←temp+1	**q1: temp←n**	0	0	?
p2: n←temp+1	q2: n←temp+1	0	0	0
(end)	**q2: n←temp+1**	1	0	0
(end)	(end)	1	0	0

- From now on
  - Assignments and Boolean evaluations are atomic!

Correctness

- What is the correct answer?
- Usually clear for sequential programs
- Can be fuzzy for concurrent programs
  - Many correct answers?
  - What is <u>intended semantics</u> of the program?
  - Run programs 100 times, each time get different answer?
    - Each answer is correct, if program is correct!
    - Does not make debugging easier!
    - Usually can not test all possible scenarios (too many!)
  - How to define correctness for concurrent programs?
    - <u>Safety properties</u> = properties that are <u>always true</u>
    - <u>Liveness properties</u> = properties that <u>eventually become true</u>

"turvallisuus"

"elävyys"

Safety and Liveness

- Safety property
  - property must be true <u>all the time</u>
    - "Identity"
      - memFree + memAllocated = memTotal
    - Mouse cursor is displayed
    - System responds to new commands
- Liveness property
  - Property must <u>eventually become true</u>
    - Variable n value = 2
    - System prompt for next command is shown
    - Control will resume to calling program
    - Philosopher will get his turn to eat
    - Eventually the mouse cursor is not displayed
    - Program will terminate
- Duality of safety and liveness properties
  - { P_i will get his turn to eat } ≡ not { P_i will never get his turn to eat }
  - { n value will become 2 } ≡ not { n value is always ≠ 2 }

safety-ominaisuus, turvallisuus

identiteetti, invariantti

elävyys, liveness-ominaisuus

Linear Temporal Logic (LTL)

(lineaarinen) temporaalilogiikka

- Define safety and liveness properties for certain state in some (arbitrary) scenario
  - Example of Modal Temporal Logic (MDL), logic on concepts like possibility, impossibility, and necessity
- Alternative: Branching Temporal Logic (BTL)
  - Properties true in some or all states starting from the given state
    - More complex, because all future states must be covered
  - Common Temporal Logic (CTL)
    - Can be checked automatically
      - Every time computation reaches given state
    - SMV model checker
    - NuSMV model checker

Fairness

reiluus

- (Weakly) fair scenario
  - Wanted condition eventually occurs
    - Nobody is locked out forever
    - Will a philosopher ever get his turn to eat?
    - Will an algorithm eventually stop?

Algorithm 2.5: Stop the loop A	
integer n ← 0 boolean flag ← false	
p	q
p1: while flag = false	q1: flag ← true
p2: n ← 1 − n	q2:

- All scenarios should be fair
  - One requirement in correct solution

Machine Language Code

- What is atomic and what is not?
  - Assignment?
  - Increment?

X = Y;

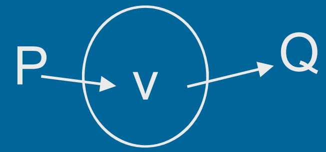
X = X+1;

Algorithm 2.6: Assignment statement for a register machine	
integer n ← 0	
p	q
p1: load R1,n	q1: load R1,n
p2: add R1,#1	q2: add R1,#1
p3: store R1,n	q3: store R1,n

Critical Reference

kriittinen viite

- Reference to variable v is critical reference, if …
  - **Assigned value in P and read in Q**
    - Read directly or in a statement



- Program satisfies limited-critical-reference (LCR)
  - Each statement has at most one critical reference
  - Easier to analyze than without this property
  - Each program is easy to transform into similar program with LCR

rajoitettu kriittinen viite

	P	Q	
Not LCR:	`n = n+1;`	`n = n+1`	Bad
Not LCR:	`n = m+1;`	`m = n+1`	Bad
LCR:	`tempP = n+1;` `n = tempP;`	`tempQ = n+1;` `n = tempQ;`	Good

LCR vs. atomicity?
(ouch)

Volatile and non-atomic variables

riskialtis

- Volatile variable
  - Can be modified by many processes (must be in shared memory)
  - Advice for compiler (pragma)
    - Keep something in memory, <u>not</u> in register
    - Pseudocode – does not generate code
- Non-atomic variables
  - Multiword data structures: long ints, arrays, records, …
  - Force access to be indivisible in given order

What if compiler/hw decides to keep value of n in a register/cache? When is it stored back to memory? What if local1 & local2 were volatile?

Algorithm 2.8: Volatile variables

integer n ← 0

p		q	
	integer local1, local2		integer local
p1:	n ← *some expression* store n?	q1:	local ← n + 6
p2:	*computation not using n* exec.	q2:	
p3:	local1 ← (n + 5) ∗ 7 order?	q3:	which n?
p4:	local2 ← n + 5 store n?	q4:	
p5:	n ← local1 ∗ local2	q5:	

Example Program with Volatile Variables

Algorithm 2.9: Concurrent counting algorithm	
integer n ← 0	
p	q
integer temp	integer temp
p1: do 10 times	q1: do 10 times
p2: temp ← n	q2: temp ← n
p3: n ← temp + 1	q3: n ← temp + 1

- Can implement it in any concurrent programming language
  - (Extended) Pascal and (Extended) C
  - BACI (Ben-Ari Concurrency Interpreter)
    - Code automatically compiled (from Extended Pascal or C)
  - Ada
  - Java

(Ben-Ari Concurrent Pascal)

possibly volatile

Concurrent Program in Pascal

```
program count;
var n: integer := 0;

procedure p;
var temp, i : integer;
begin
  for i := 1 to 10 do
    begin
    temp := n;
    n := temp + 1
    end
end;
```

```
procedure q;
var temp, i : integer;
begin
  for i := 1 to 10 do
    begin
    temp := n;
    n := temp + 1
    end
end;

begin { main program }
  cobegin p; q coend;
  writeln(' The value of n is ', n)
end.
```

n is volatile, because… it is assigned in one thread, and read in the other

Concurrent Program in C (Ben-Ari Concurrent C, C--)

```
int n = 0;

void p() {
  int temp, i;
  for (i = 0; i < 10; i++) {
    temp = n;
    n = temp + 1;
  }
}

void q() {
  int temp, i;
  for (i = 0; i < 10; i++) {
    temp = n;
    n = temp + 1;
  }
}

void main() {
  cobegin { p(); q(); }
  cout << "The value of n is " << n << "\n";
}
```

possibly volatile, use carefully

(volatile, if critically referenced)

What if compiler optimized and kept n in a register?
Lets hope not!
(in ExtPascal or C--
global (volatile) variables are seemingly kept in memory by default)

Concurrent Program in Ada

```
with Ada.Text_IO; use Ada.Text_IO;
procedure Count is
    N: Integer  := 0;
    pragma Volatile(N);

    task type Count_Task;
    task body Count_Task is
        Temp: Integer;
    begin
        for I in 1..10 loop
            Temp := N;
            N := Temp + 1;
        end loop;
    end Count_Task;

begin
    declare
        P, Q: Count_Task;
    begin
        null;
    end;
    Put_Line("The value of N is " & Integer'Image(N));
end Count;
```

advice compiler to keep N in memory

Concurrent Program in Java

```
class Count extends Thread {
    static volatile int n = 0;

    public void run() {
        int temp;
        for (int i = 0; i < 10; i++) {
            temp = n;
            n = temp + 1;
        }
    }
    public static void main(String[] args) {
        Count p = new Count();
        Count q = new Count();
        p.start ();
        q.start ();
        try {
            p.join ();
            q.join ();
        }
        catch (InterruptedException e) { }
        System.out.println ("The value of n is " + n);
    }
}
```

Thread.yield(); // force?

How many threads really in parallel?
• how to control it?

Execute on 8-processor vera.cs.helsinki.fi?

http://www.cs.helsinki.fi/u/kerola/rio/Java/examples/Adder8.java

http://www.cs.helsinki.fi/u/kerola/rio/Java/examples/Adder8b.java

> javac Adder8.java
> java Adder8

BACI

http://www.mines.edu/fs_home/tcamp/baci/

- Ben-Ari Concurrency Interpreter
  - Write concurrent programs with
    - *C--* or *Ben-Ari Concurrent Pascal* (.cm and .pm suffixes)
    - Compile and run in BACI
  - GUI for Unix/Linux
- jBACI http://stwww.weizmann.ac.il/g-cs/benari/jbaci/
  - Just like BACI
  - GUI for Windows
- Installation http://stwww.weizmann.ac.il/g-cs/benari/jbaci/jbaci1-4-5.zip
  - load version 1.4.5 jBACI executable files and example programs, unzip, edit config.cfg to have correct paths to bin/bacc.exe and bin/bapas.exe translators, click run.bat
- Use in class, homeworks and in project

BACI Overall Structure

add.cm

C- - (Concurrent C)

```
…
void main() {
  cobegin { add10();
            add10(); }
….
```

bacc.exe

C-- to PCODE Compiler

add.lst

```
…
17   24  void main() {
18   25      cobegin {add10(); add10(); }
….
```

add.pco

(many tables)

```
LOAD_ADDR, push sum
LOAD_VALUE, push local
PUSH_LIT 1
DO_ADD, pop(1), s[t]=(s[oldt-1]+s[oldt])
STORE, s[s[t-1]]=s[t], pop(2)
```

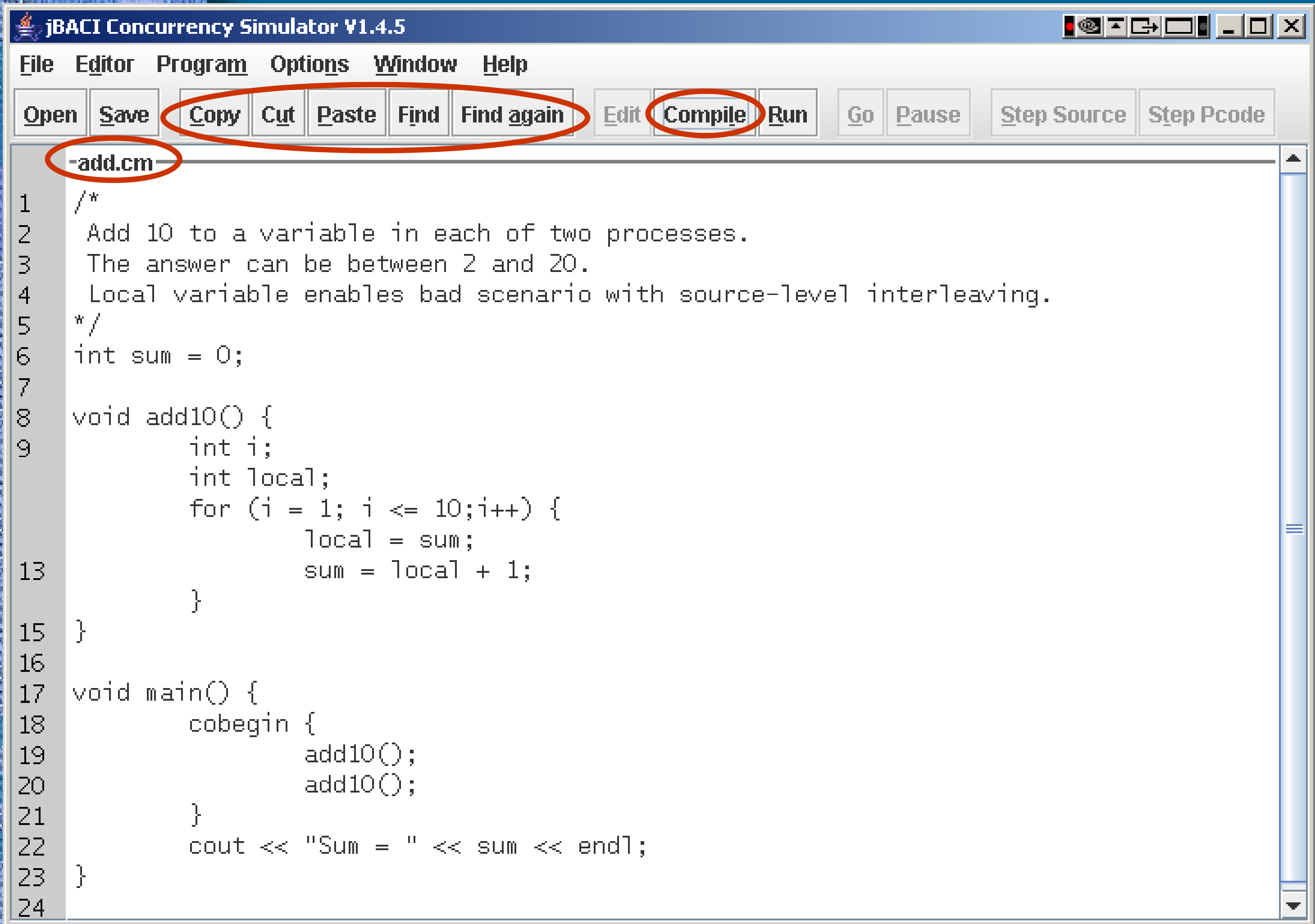
PCODE Interpreter

bainterp.exe

```
Executing PCODE ...
C  n =1  i =A  n =1 C2  i =
1 A
C  n =4  i =2 C
B  n =A  n =5  i =24 A
```

http://www.cs.helsinki.fi/u/kerola/rio/BACI/baci-c.pdf

jBACI

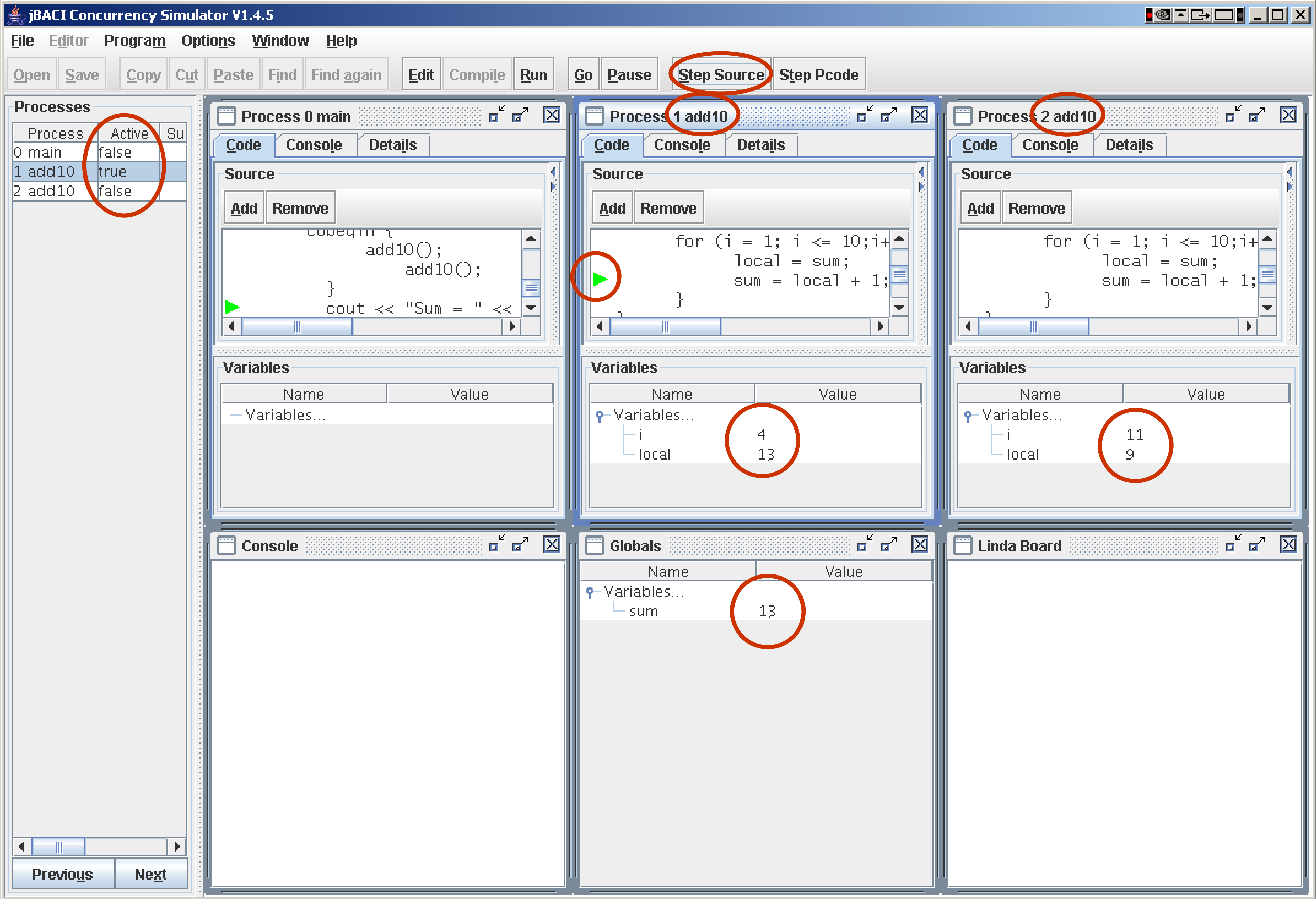


Just like BACI, but with Java

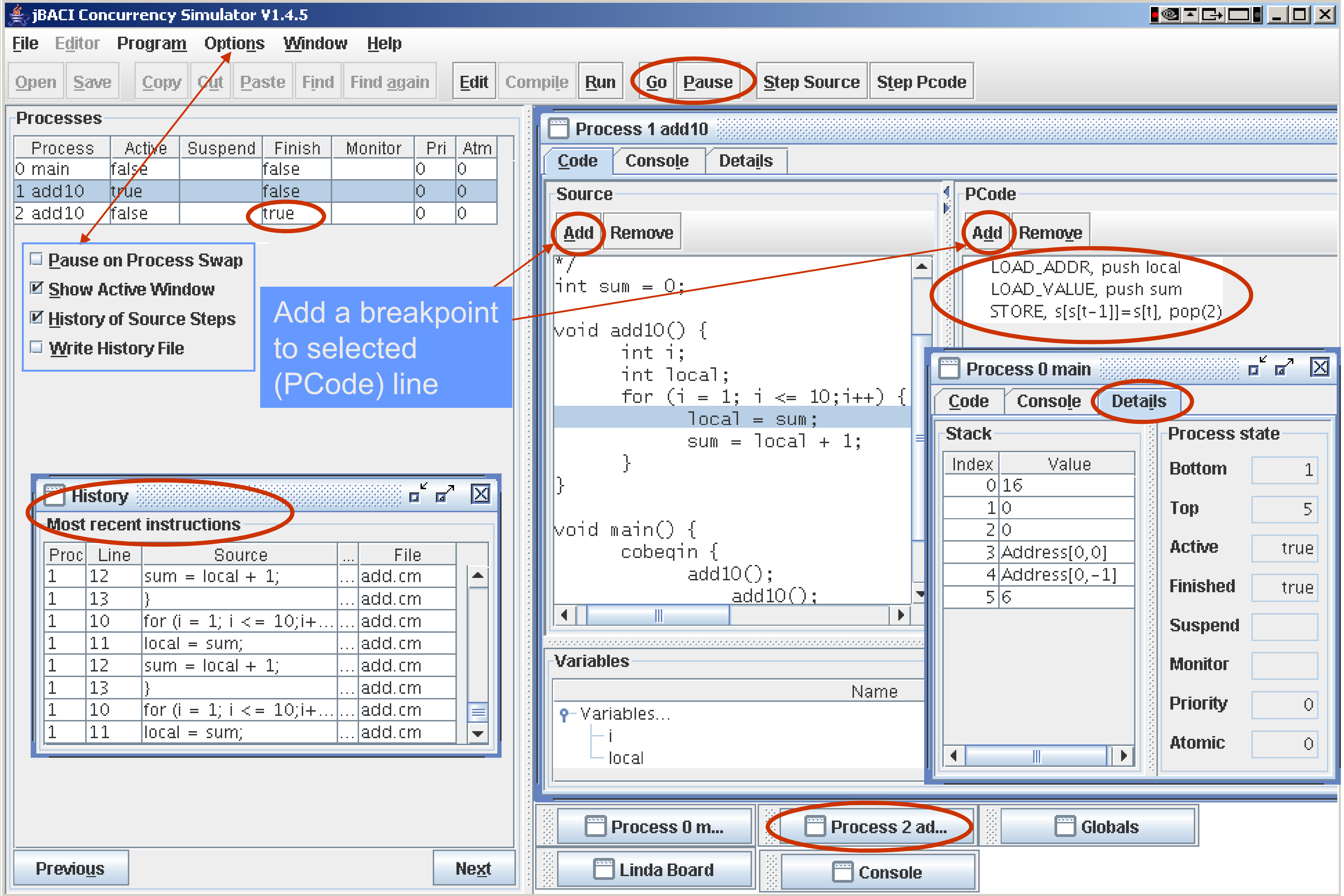
- requires Java v. 1.4 (SDK or JRE)
- Built-in compiler and interpreter
- edit state
- run state

http://www.cs.helsinki.fi/u/kerola/rio/BACI/jbaci.pdf

jBACI IDE (integrated development environment)



jBACI IDE (integrated development environment)



Summary

- Abstraction, atomicity
- Concurrent program, program state
- Pseudo-language algorithms
- High level language algorithms
- BACI