



Overlay (and P2P) Networks

Part II

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Models of Complex Networks to Model Overlay Networks

- Milgram's Experiment
- Duncan Watts Random Rewiring Model
- Scale Free Networks (Power-Law Networks)
 - Preferential attachment
 - Evolving Copying Model (Copying Generative Model)
- Navigation in Small World

**Complex
Networks**

**Overlay
Networks** 

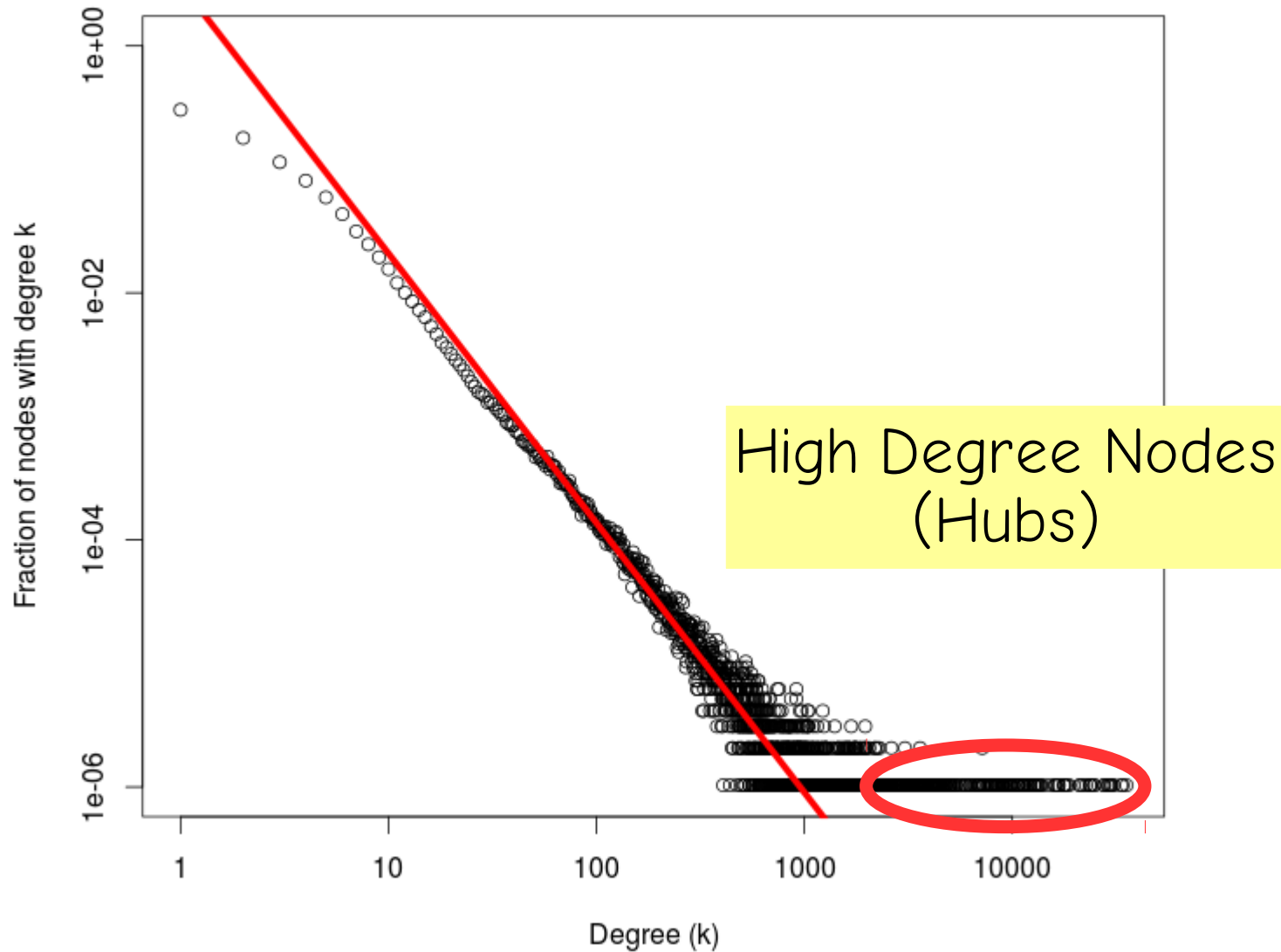


Outline for this lecture

- Error and Attack Tolerance of Complex Networks
- Navigation in Complex Networks
- Mathematics and the Internet: A Source of Enormous Confusion and Great Potential
- Summary on Modeling Overlay Networks



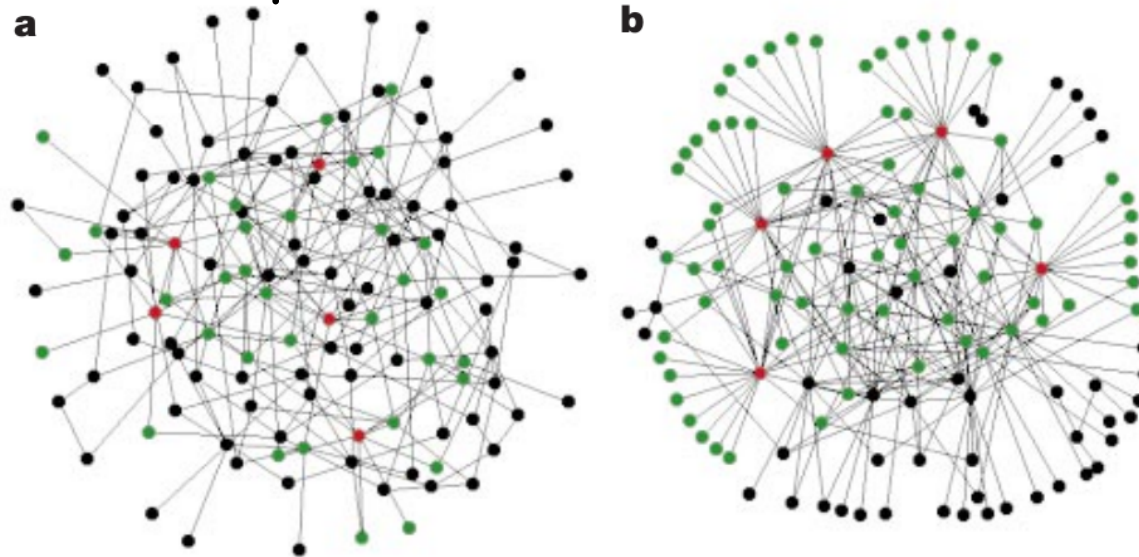
Scale-Free Model for AS-Graph



AS Topology of skitter dataset parsed by SNAP team
<http://snap.stanford.edu/data/as-skitter.html>



Importance of Hubs



(Random Graph) Exponential

Scale-free

Figure 1 Visual illustration of the difference between an exponential and a scale-free network. **a**, The exponential network is homogeneous: most nodes have approximately the same number of links. **b**, The scale-free network is inhomogeneous: the majority of the nodes have one or two links but a few nodes have a large number of links, guaranteeing that the system is fully connected. Red, the five nodes with the highest number of links; green, their first neighbours. Although in the exponential network only 27% of the nodes are reached by the five most connected nodes, in the scale-free network more than 60% are reached, demonstrating the importance of the connected nodes in the scale-free network Both networks contain 130 nodes and 215 links ($\langle k \rangle = 3.3$). The network visualization was done using the Pajek program for large network analysis: (<http://vlado.fmf.uni-lj.si/pub/networks/pajek/pajekman.htm>).

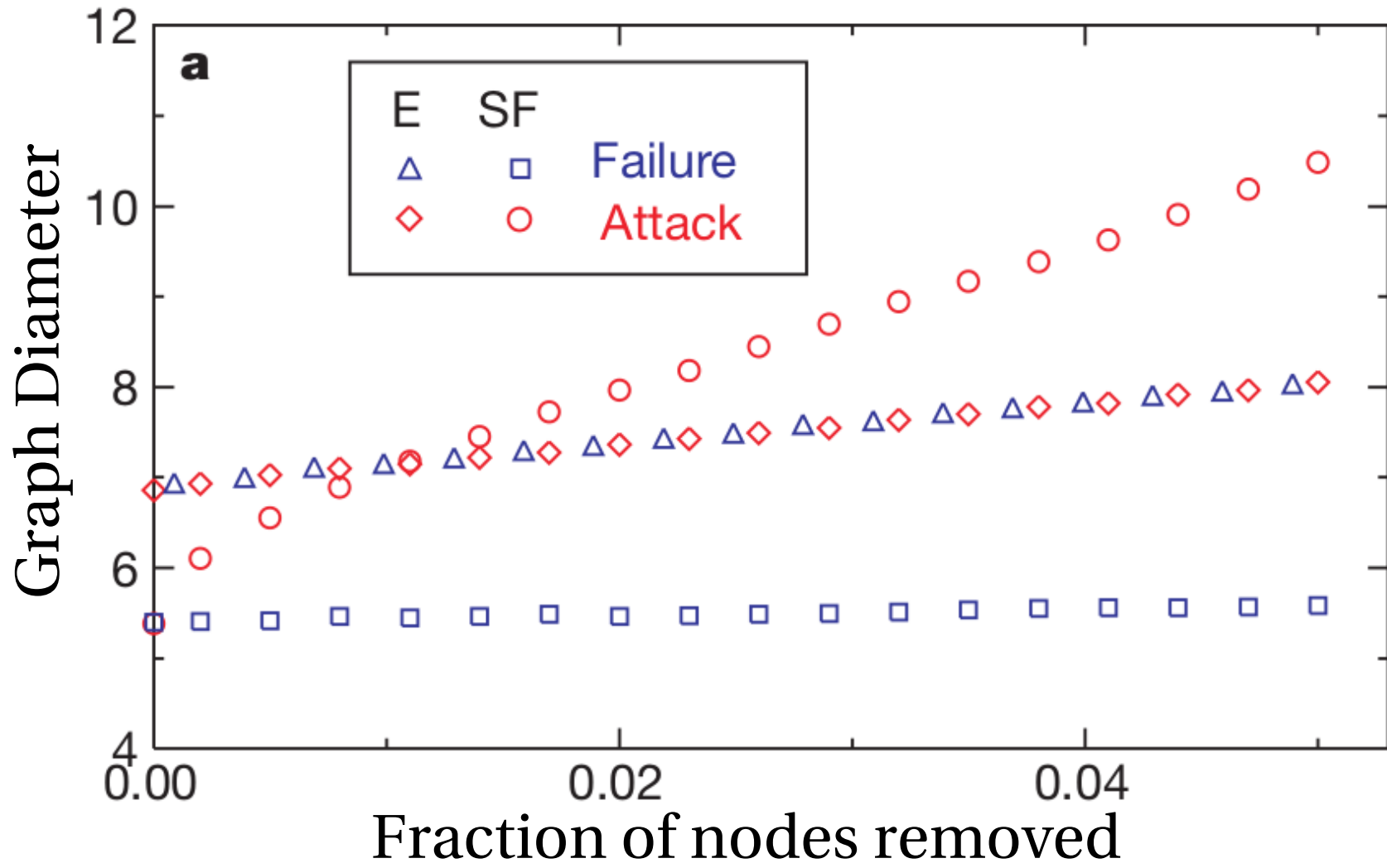


Error vs Attack

- Error (Node Failure)
 - random node fails (malfunction)
- Attack
 - Selected node with a given property is made to fail
 - Which nodes would you target if you knew the network is a scale-free network?
 - *Nodes with the highest degree*



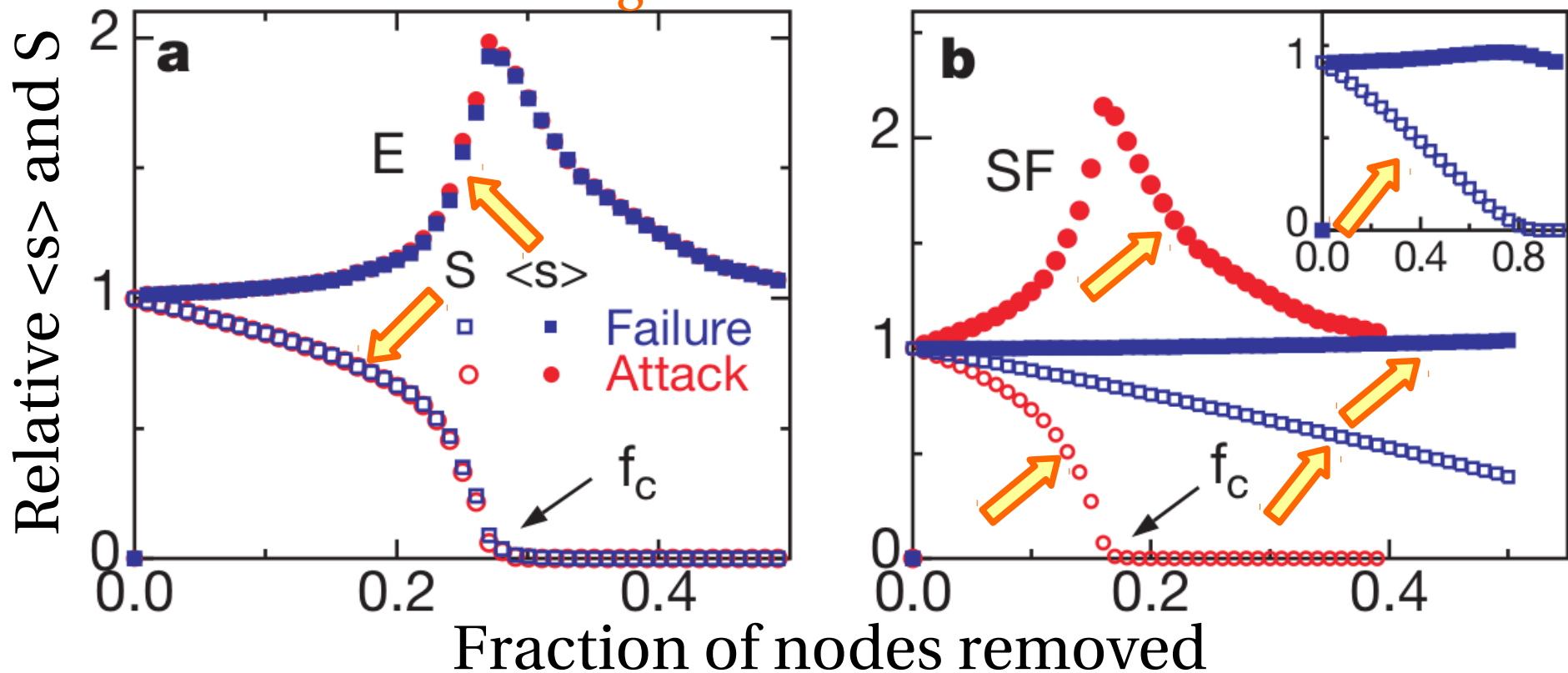
Impact of Errors and Attacks (Graph Diameter)





Impact of Errors and Attacks (Size of Largest Cluster)

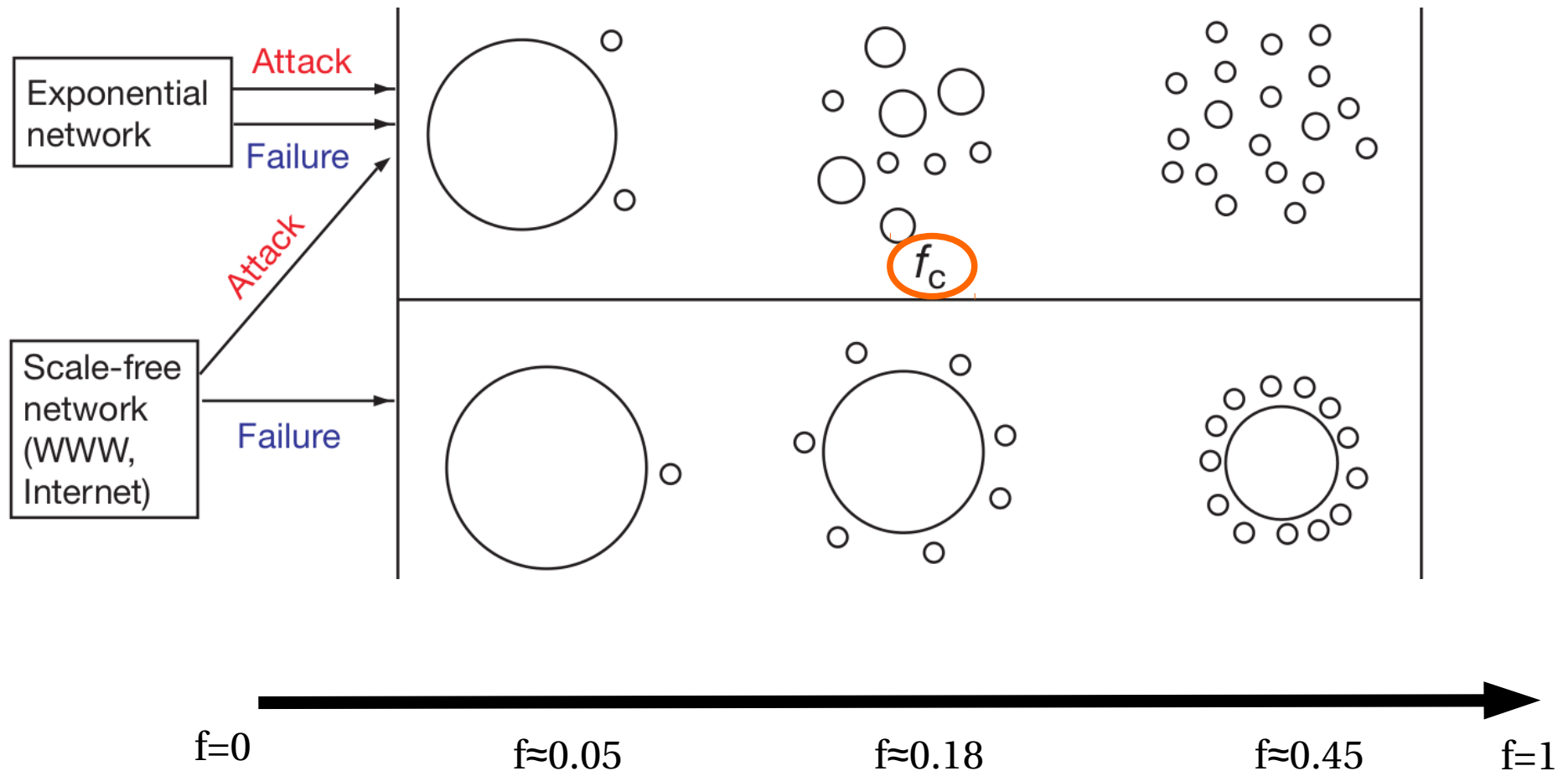
S : Fraction of nodes in largest cluster
 $\langle s \rangle$: average size of isolated clusters



Albert, Réka, et al. "Error and attack tolerance of complex networks."
nature 406, no. 6794 (2000): 378-382.



Network Response to Attacks and Failures



Albert, Réka, et al. "Error and attack tolerance of complex networks." nature 406, no. 6794 (2000): 378-382.



Critical Threshold

(random node failures)

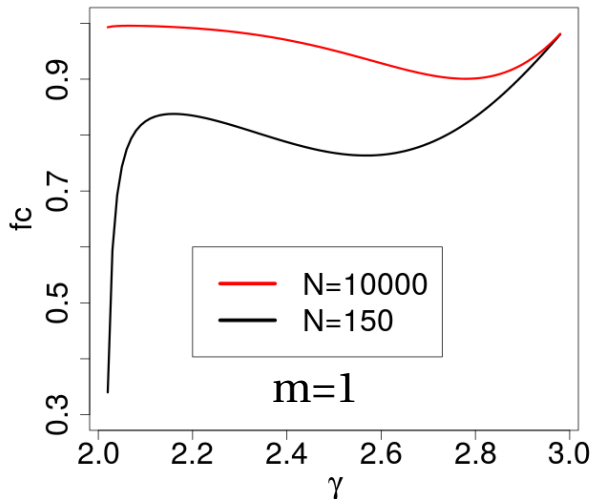
$$f_c = 1 - \frac{1}{\beta - 1} \begin{cases} \gamma : \text{exponent of power-law} \\ m : \text{smallest degree} \\ N : \text{number of nodes in the graph} \\ K : \text{largest degree, } K \approx mN^{\frac{1}{\gamma-1}} \end{cases}$$

where

$$\beta = \frac{|2 - \gamma|}{|3 - \gamma|} \times \begin{cases} m & \text{if } \gamma > 3 \\ m^{\gamma-2} K^{3-\gamma} & \text{if } 2 < \gamma < 3 \\ K & \text{if } 1 < \gamma < 2 \end{cases}$$

for $2 < \gamma < 3$

$$f_c = 1 + \left(1 - m^{(\gamma-2)} K^{(3-\gamma)} \frac{\gamma - 2}{3 - \gamma} \right)^{-1}$$



Cohen's technique can be extended to errors
(No closed form for f_c for errors)

Cohen, Reuven et al. "**Resilience of the Internet to random breakdowns.**" Physical review letters 85, no. 21 (2000): 4626.

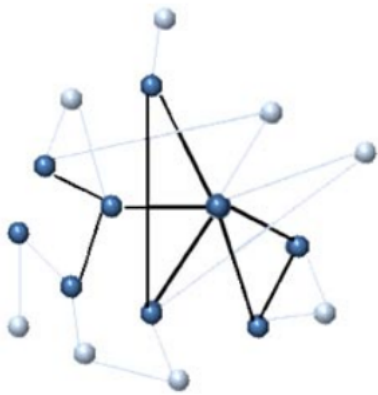


Summary on Attack and Error Tolerance of Complex Networks

Scale-free networks resilient to random failures but vulnerable to targetted attacks



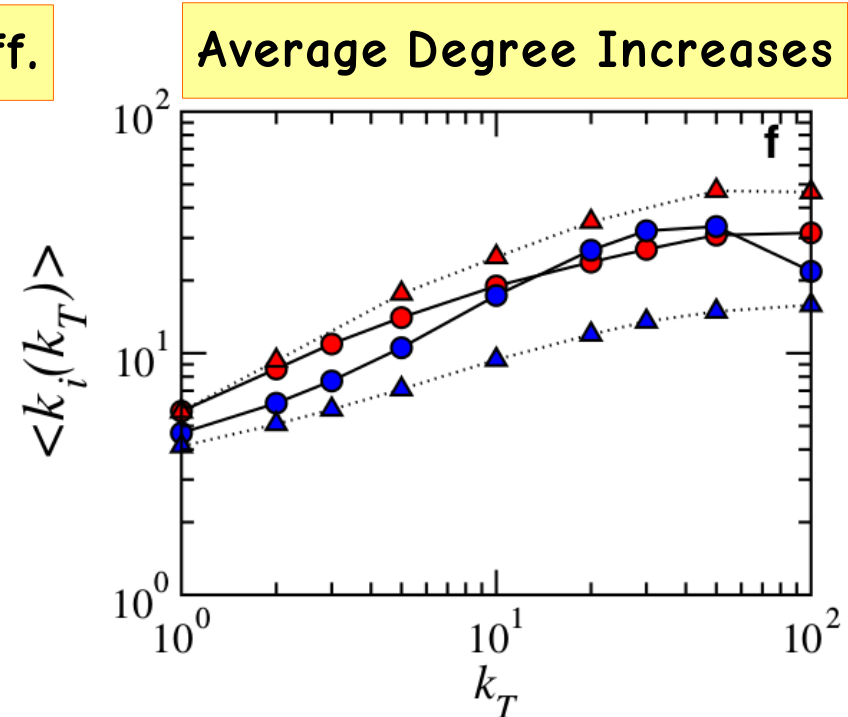
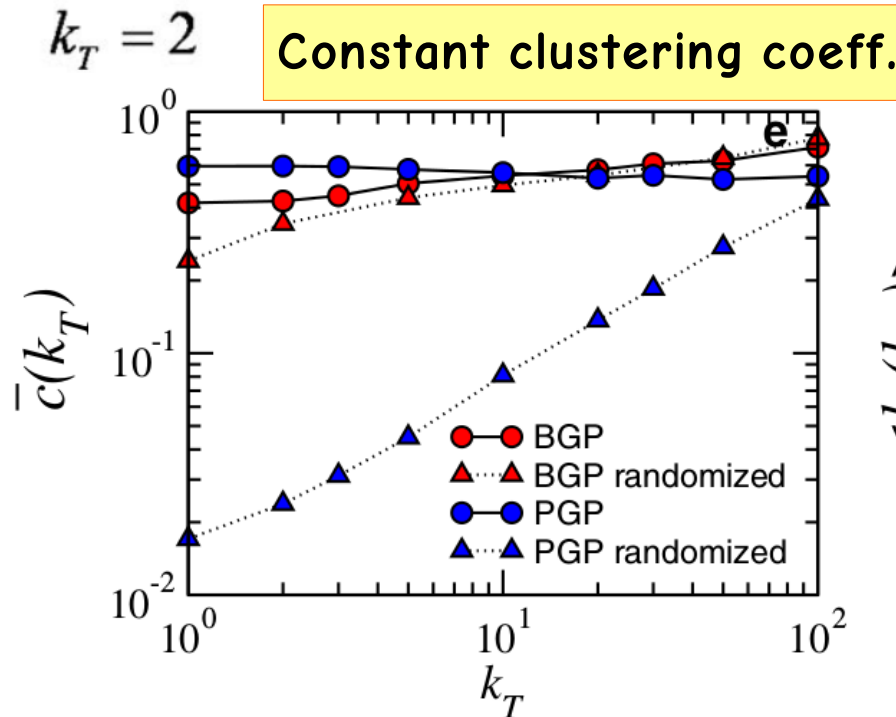
Clustering of Nodes



k_T : extract subgraph $G(k_T)$ with nodes having degree $k > k_T$

$\langle k_i(k_T) \rangle$: average degree of this subgraph $G(k_T)$

$\bar{c}(k_T)$: clustering coefficient of subgraph $G(k_T)$



How to Generate Scale-Free Graphs with Strong Clustering

Serrano, M. Angeles, Dmitri Krioukov, and Marián Boguná. "Self-similarity of complex networks and hidden metric spaces." Physical review letters 100, no. 7 (2008): 078701.



Generating Scale-Free Graphs with Strong Clustering

Take all nodes and distribute them within an underlying circle

Assign each node an expected degree k where $P(k) \sim k^{-\gamma}$

Connect each pair of nodes with a connection probability $r(d; k, k')$

d is the distance between these two nodes in the circle

$d_c = kk'$ is also called the characteristic distance

$$r(d; k, k') = \left(1 + \frac{d}{d_c}\right)^{-\alpha}$$

Hubs will be connected with a high probability because of large d_c

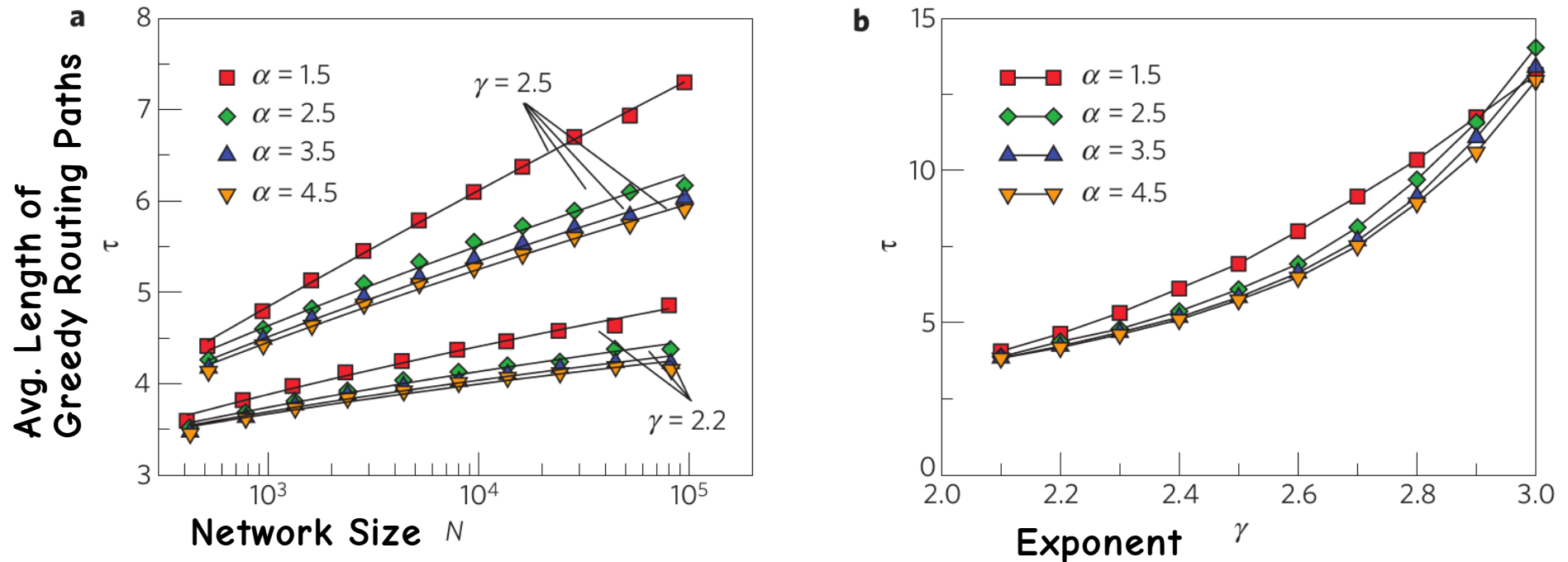
Low degree nodes connected only if (hidden distance) d is small

Hubs connected to low degree nodes at moderate hidden distance

α importance of hidden distance



Path Length (Greedy Routing)



Path length grows polylogarithmically with the network size

Paths shorter for smaller exponents and stronger clustering

Boguna, Marian et al. "**Navigability of complex networks.**"
Nature Physics 5, no. 1 (2009): 74-80.

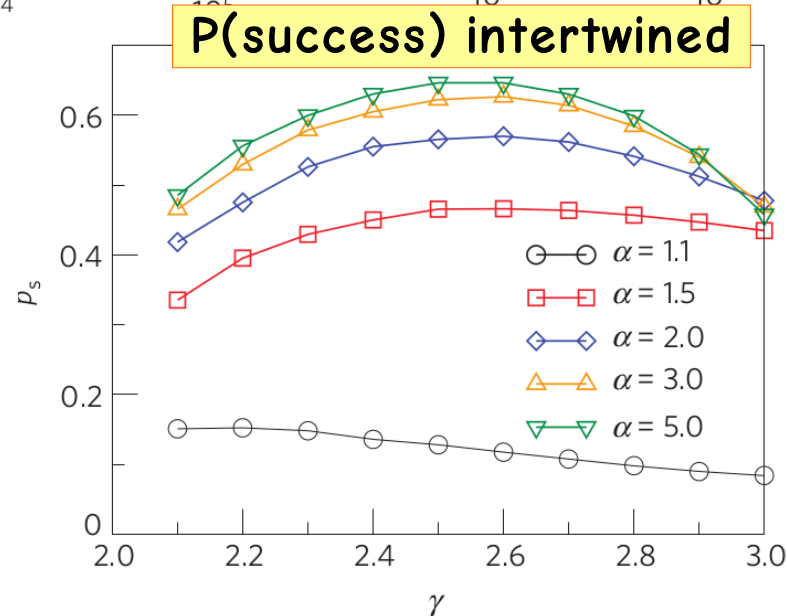
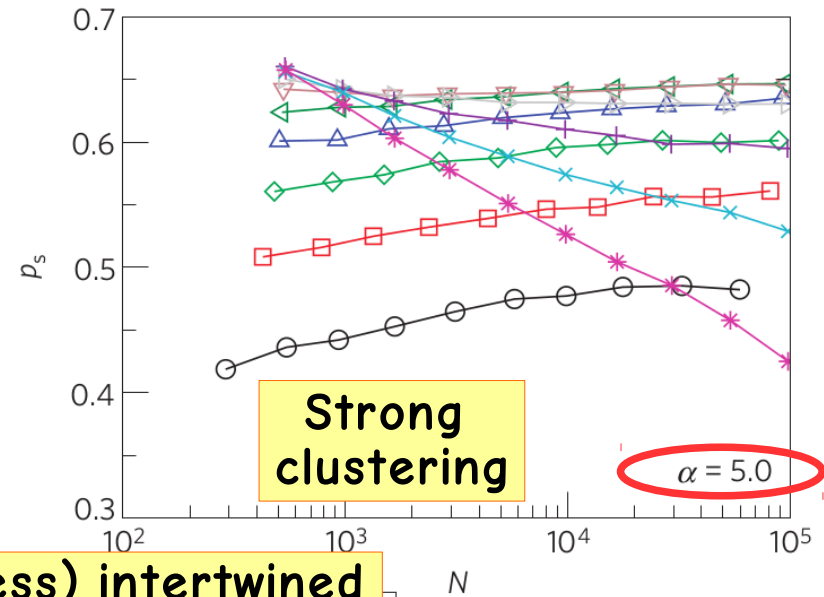
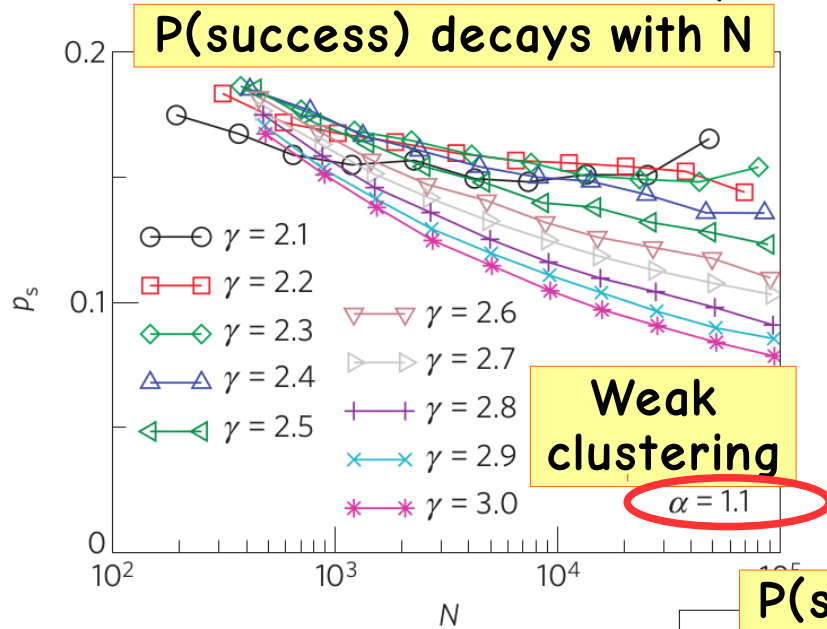


Greedy Routing

- Hidden Space as the coordinate space
 - *Hidden space is circle in this example*
- Greedy Routing: Send to neighbor who is closer to the destination (in hidden space)
- Unsuccessful Paths: None of your neighbors are closer to the destination in the hidden space



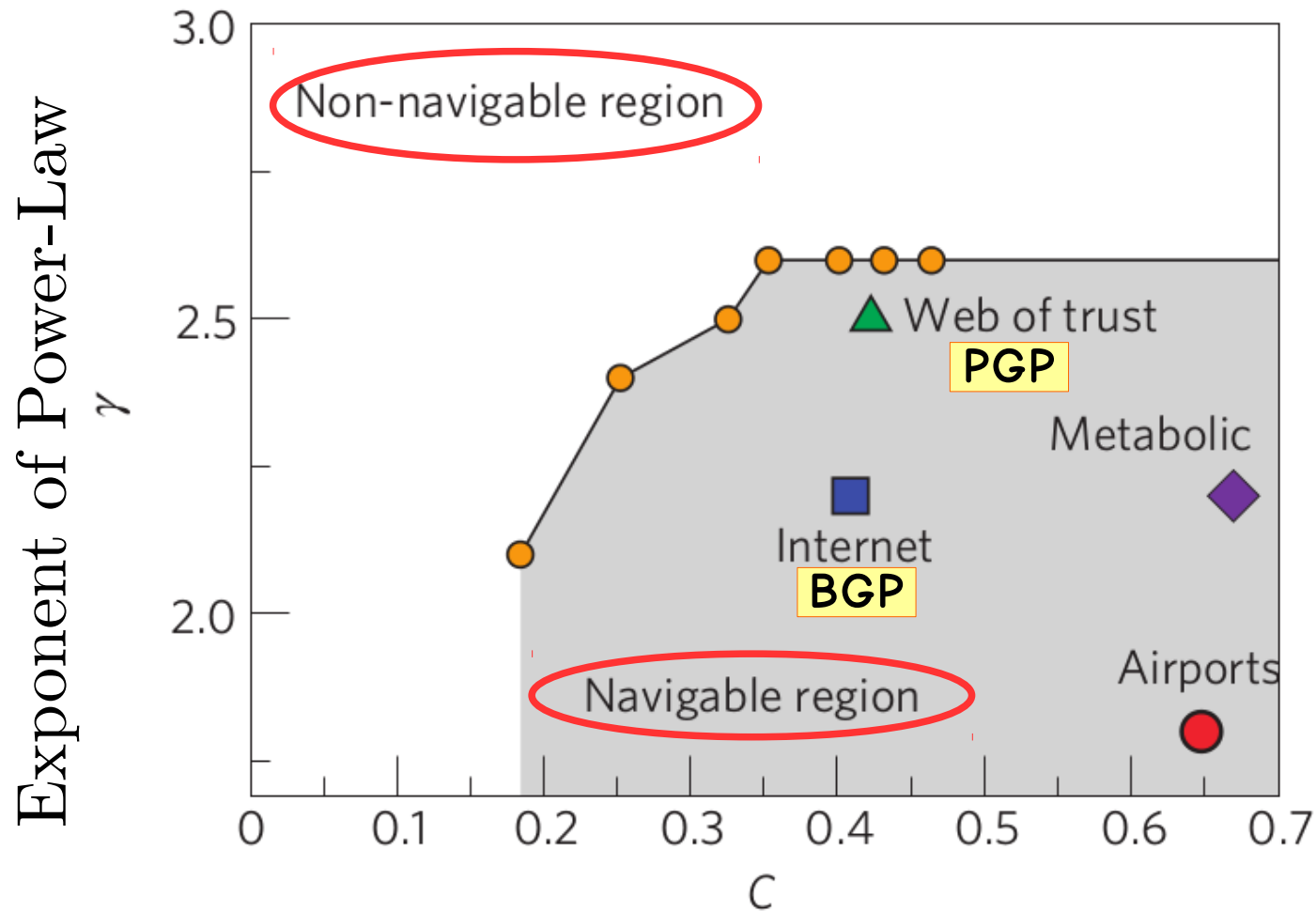
Success Probability (Greedy Routing)



Boguna, Marian et al.
"Navigability of complex networks."
 Nature Physics 5, no. 1 (2009): 74-80.



Navigation in Scale Free Networks



Clustering Coefficient $C = f(\gamma, \alpha)$

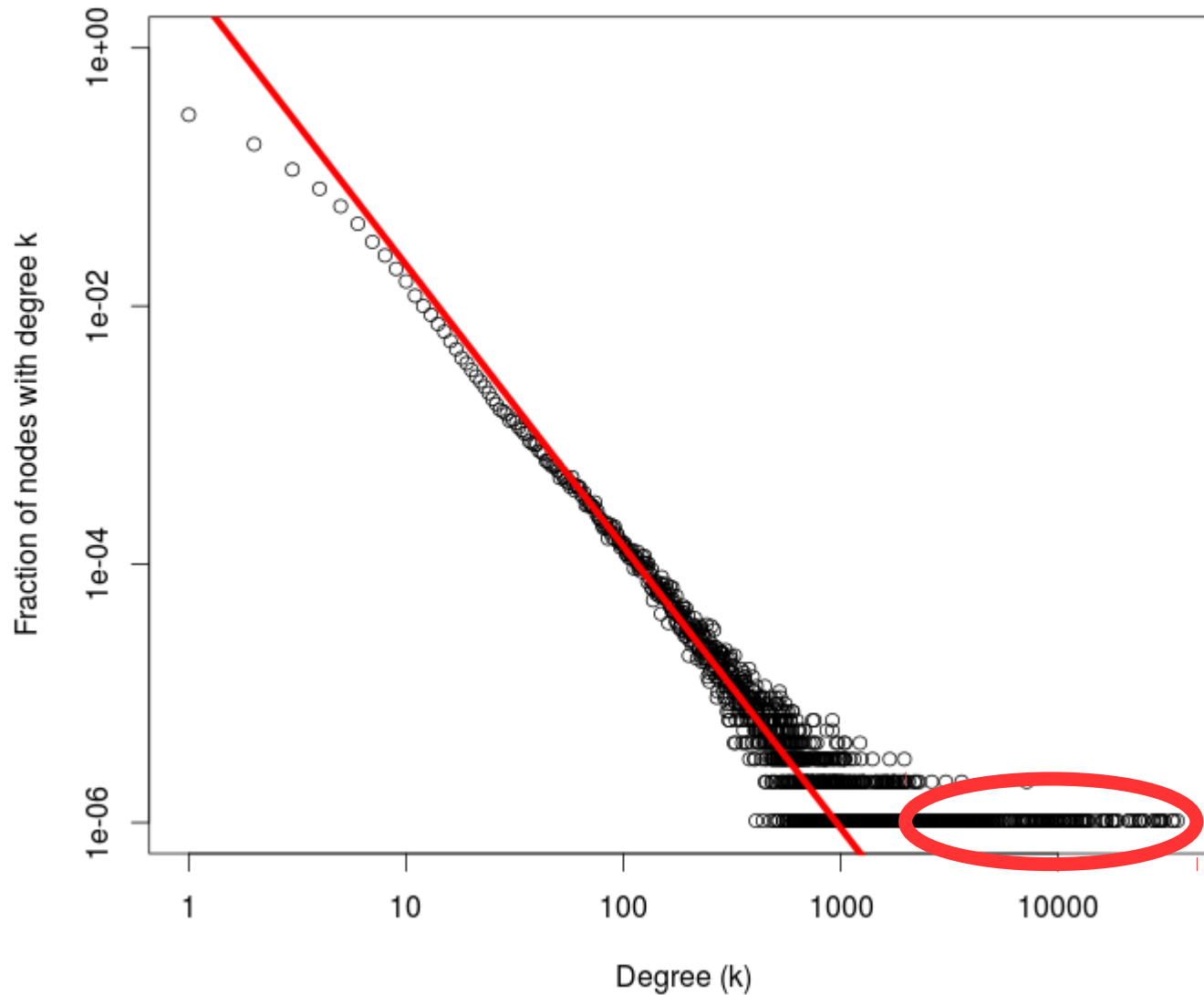


Implications of Result

- Internet Routing
 - Routers currently exchange signals to keep coherent view of network
 - Network size increasing with time
 - Hidden metric space eliminates the need for control signals exchanged to notify changes in network
- How to proceed to discover the hidden metric space
- Does Shortest Path imply Shortest Time to destination?
 - What happens in case of congestion at hubs?



Scale-Free Model for AS-Graph



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Is the Scale-Free Internet A Myth?

- What we have seen till now wrt to Preferential Attachment
 - Preferential attachment results in Hubs
 - Hubs vulnerable to coordinated attacks
 - Why is the Internet still up and running
- Is the Scale-Free modeling paradigm consistent with the engineered nature of the Internet and the design constraints imposed by existing technology?
 - Is the simplistic toy model too generic?
 - Do the available measurements, their analysis, and their modeling efforts support the claims made by "Error and Attack Tolerance" paper?



Importance of Measurements

- Tool for measurement study for AS-measurements
 - Traceroute
- Biases of traceroute
 - Uses IPv4 Protocol
 - What about non-IPv4 protocols like MPLS?
 - Entry points to non-IPv4 regions can aggregate to Hubs
 - Only reports the interfaces traversed by the packet
 - Routers can have multiple interfaces and appear on different routes with different IP addresses



Leverage Domain Knowledge

- Device Constraints
 - Finite number of interfaces on routers
 - Finite capacity of routers
- Placement of High Degree Nodes
 - Edge vs Core

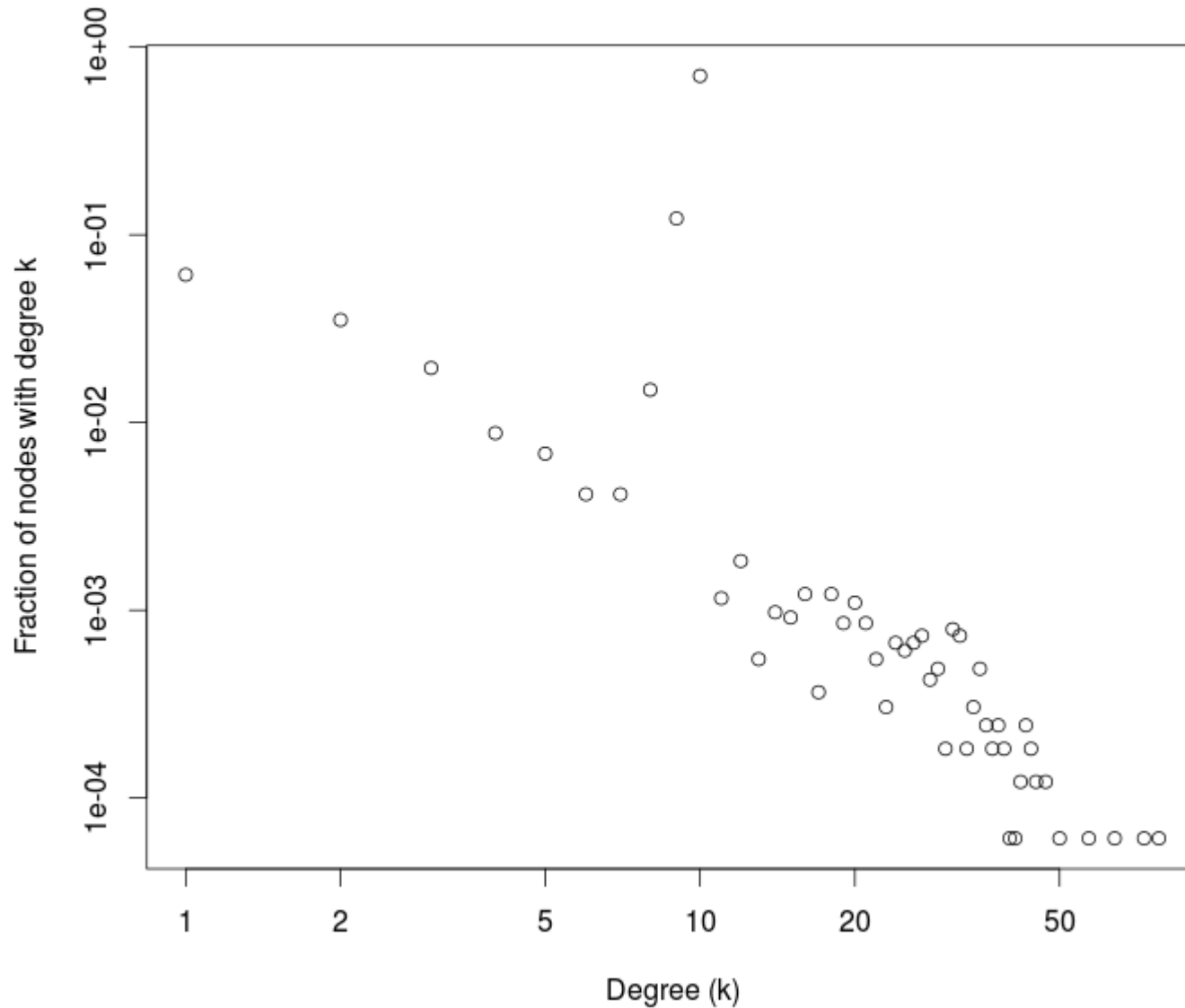
**What about
Overlay Networks?**

How would you deploy the network if you are a network engineer?

- Leverage domain knowledge to identify driving forces behind the design of high engineered systems such as the Internet



Scale Free for Gnutella?





Recap of Modeling Overlay Networks

- Milgram's Experiment
- Duncan Watts Random Rewiring Model
- Scale-Free Networks
 - Preferential attachment
 - Evolving Copying Model (Copying Generative Model)
 - Scale-Free with Strong Clustering
- Error and Fault Tolerance of Complex Networks
- Navigation (Greedy Routing)
 - In Small World (Kleinberg's Small World)
 - In Complex Networks (Scale-Free with Strong Clustering)
- Mathematics and the Internet: A Source of Enormous Confusion and Great Potential



Commonly used metrics

- Clustering Coefficient
- Diameter
- Degree Distribution



Methodology

- 1) Make observations (conduct measurement studies)
- 2) Build model to explain observations
 - Choose the right level of granularity (zoom level)
 - Strip the problem to a simple form
 - Attempt to formulate the problem and model the system
- 3) Validate model
 - Reproduce observations/measurements
 - Explain observations
- 4) Revisit step 2 (and 1) to improve understanding



Important Articles

- Milgram, Stanley. "**The small world problem.**" *Psychology today* 2.1 (1967): 60-67
- Watts, Duncan and Strogatz, Steven. "**Collective dynamics of 'small-world' networks.**" *Nature* 393.6684 (1998): 440-442.
- Barabási, Albert-László, and Albert, Réka. "**Emergence of scaling in random networks.**" *Science* 286, no. 5439 (1999): 509-512.
- Kleinberg, Jon. "**The small-world phenomenon: An algorithmic perspective.**" In *ACM Symposium on Theory of computing*, pp. 163-170. 2000.
- Ravi Kumar et al. "**Stochastic models for the web graph.**" In *Annual Symposium on Foundations of Computer Science*, 2000.
- Albert, Réka, and Barabási, Albert-László. "**Statistical mechanics of complex networks.**" *Reviews of modern physics* 74.1 (2002): 47.
- Newman, Mark. "**The structure and function of complex networks.**" *SIAM review* 45, no. 2 (2003): 167-256.
- Mitzenmacher, M. (2004). "**A brief history of generative models for power law and lognormal distributions.**" *Internet mathematics*, 1(2), 226-251.
- Mark Newman. "**Power laws, Pareto distributions and Zipf's law.**" *Contemporary physics* 46, no. 5 (2005): 323-351
- Jure Leskovec et al. "**Graphs over time: densification laws, shrinking diameters and possible explanations.**" In *ACM SIGKDD*, pp. 177-187. 2005.
- Boguna, Marian et al. "**Navigability of complex networks.**" *Nature Physics* 5, no. 1 (2009): 74-80.
- W Willinger et al. "**Mathematics and the internet: A source of enormous confusion and great potential.**" In *Notices of the AMS*. 2009.



Sources for these slides

- Sasu Tarkoma "**Overlay and P2P Networks**", 2015
- Datasets from Stanford Network Analysis Project (SNAP)