



$$\frac{\Gamma \Rightarrow A \quad \Delta, B \Rightarrow C}{\Gamma, \Delta, A \supset B \Rightarrow C} L\supset_{ind} \qquad \frac{\Gamma \Rightarrow A \quad \Gamma, B \Rightarrow C}{\Gamma, A \supset B \Rightarrow C} L\supset_{sh}$$

$$\frac{\Gamma \Rightarrow A \quad \Delta, A \Rightarrow C}{\Gamma, \Delta \Rightarrow C} Cut_{ind} \qquad \frac{\Gamma \Rightarrow A \quad \Gamma, A \Rightarrow C}{\Gamma \Rightarrow C} Cut_{sh}$$

We will start with right conjunction: We will assume the premisses of  $R\&_{ind}$ , use weakening repeatedly to get the same context to both sides and derive the conclusion of  $R\&_{ind}$  using  $R\&_{sh}$  as follows ( $Wk^*$  means  $Wk$  applied  $0 \dots n$  times):

$$\frac{\frac{\Gamma \Rightarrow A}{\Gamma, \Delta \Rightarrow A} Wk^* \quad \frac{\Delta \Rightarrow B}{\Gamma, \Delta \Rightarrow B} Wk^*}{\Gamma, \Delta \Rightarrow A \& B} R\&_{sh}$$

The other direction can be proved by assuming the premisses of  $R\&_{sh}$ , by applying  $R\&_{ind}$  and by using contraction repeatedly in order to remove duplicated formulae from the context:

$$\frac{\frac{\Gamma \Rightarrow A \quad \Gamma \Rightarrow B}{\Gamma, \Gamma \Rightarrow A \& B} R\&_{ind}}{\Gamma \Rightarrow A \& B} Ctr^*$$

The other cases are similar:

$$\frac{\frac{\Gamma, A \Rightarrow C}{\Gamma, \Delta, A \Rightarrow C} Wk^* \quad \frac{\Delta, B \Rightarrow C}{\Gamma, \Delta, B \Rightarrow C} Wk^*}{\Gamma, \Delta, A \vee B \Rightarrow C} L\vee_{sh} \qquad \frac{\Gamma, A \Rightarrow C \quad \Gamma, B \Rightarrow C}{\Gamma, \Gamma, A \vee B \Rightarrow C} L\vee_{ind}$$

$$\frac{\Gamma, A \vee B \Rightarrow C}{\Gamma, A \vee B \Rightarrow C} Ctr^*$$

$$\frac{\frac{\Gamma \Rightarrow A}{\Gamma, \Delta \Rightarrow A} Wk^* \quad \frac{\Delta, B \Rightarrow C}{\Gamma, \Delta, B \Rightarrow C} Wk^*}{\Gamma, \Delta, A \supset B \Rightarrow C} L\supset_{sh} \qquad \frac{\Gamma \Rightarrow A \quad \Gamma, B \Rightarrow C}{\Gamma, \Gamma, A \supset B \Rightarrow C} L\supset_{ind}$$

$$\frac{\Gamma, A \supset B \Rightarrow C}{\Gamma, A \supset B \Rightarrow C} Ctr^*$$

$$\frac{\frac{\Gamma \Rightarrow A}{\Gamma, \Delta \Rightarrow A} Wk^* \quad \frac{\Delta, A \Rightarrow C}{\Gamma, \Delta, A \Rightarrow C} Wk^*}{\Gamma, \Delta \Rightarrow C} Cut_{sh} \qquad \frac{\Gamma \Rightarrow A \quad \Gamma, A \Rightarrow C}{\Gamma, \Gamma \Rightarrow C} Cut_{ind}$$

$$\frac{\Gamma, \Gamma \Rightarrow C}{\Gamma \Rightarrow C} Ctr^*$$

□

3. Show that reductio ad absurdum ( $Raa$ ) is derivable using the natural deduction rules for intuitionistic logic and the rule of excluded middle ( $Em$ ).

The proof proceeds by assuming the premisses of the rule  $Raa$  and deriving its conclusion

without using the rule. The rule is  $\frac{[\sim A] \quad \dots \quad \perp}{\sim A}$  so we will assume that we have the derivation  $\dots \quad \perp$  and derive  $A$  as follows:

$$\frac{[\sim A]^1 \quad \dots \quad \perp}{A} \frac{\perp E}{Em, I}$$

□

4. Prove in natural deduction for classical logic (intuitionistic logic + *Em*):

(a)  $\sim(\sim A \ \& \ \sim B) \supset A \vee B$

$$\frac{\frac{\frac{[A]^2}{A \vee B} \vee I}{[B]^3 \quad A \vee B} \text{Em},3}{\frac{A \vee B}{\sim(\sim A \ \& \ \sim B) \supset A \vee B} \supset I,4} \quad \frac{\frac{[\sim(\sim A \ \& \ \sim B)] \quad \frac{\frac{[\sim A]^2 \quad [\sim B]^3}{\sim A \ \& \ \sim B} \&I}{\perp} \perp E}{[\perp]^1} \supset E,1}{\perp} \perp E,2}{\sim(\sim A \ \& \ \sim B) \supset A \vee B} \text{Em},2$$

(b)  $(A \supset B) \supset \supset (\sim A \vee B)$

$$\frac{\frac{\frac{[A \supset B]^2 \quad [A]^3 \quad [B]^1}{B} \supset E,1}{\sim A \vee B} \vee I}{(A \supset B) \supset (\sim A \vee B)} \supset I,2}{\frac{(A \supset B) \supset (\sim A \vee B)}{(A \supset B) \supset \supset (\sim A \vee B)} \text{Em},3} \quad \frac{\frac{[\sim A]^3}{\sim A \vee B} \vee I}{[\sim A \vee B]^7} \supset I}{\frac{[\sim A]^6 \quad [A]^5 \quad \frac{[\perp]^4}{B} \perp E}{B} \supset E,4}{\frac{A \supset B}{A \supset B} \supset I,5}{\frac{A \supset B}{(\sim A \vee B) \supset (A \supset B)} \supset I,4} \supset E,1}{\frac{(\sim A \vee B) \supset (A \supset B)}{(A \supset B) \supset \supset (\sim A \vee B)} \&I} \supset E,6$$

(c)  $((A \supset B) \supset A) \supset A$  (Peirce's law)

$$\frac{\frac{[A]^5}{((A \supset B) \supset A) \supset A} \supset I}{((A \supset B) \supset A) \supset A} \supset I,4}{\frac{[\sim A]^6 \quad [A]^2 \quad \frac{[\perp]^1}{B} \perp E}{B} \supset E,1}{\frac{[(A \supset B) \supset A]^4 \quad \frac{A \supset B}{A \supset B} \supset I,2}{A} \supset I,2}{\frac{A}{((A \supset B) \supset A) \supset A} \supset I,4} \supset E,3} \text{Em},5,6$$

(d)  $(A \supset B \vee C) \supset (A \supset B) \vee (A \supset C)$  (disjunction property under hypothesis)

$$\frac{\frac{[A \supset B \vee C]^6 \quad [A]^5 \quad [B \vee C]^1}{B \vee C} \supset E,1}{\frac{\frac{[B]^2}{A \supset B} \supset I}{(A \supset B) \vee (A \supset C)} \vee I}{\frac{[C]^2}{A \supset C} \supset I}{(A \supset B) \vee (A \supset C)} \vee I}{\frac{(A \supset B) \vee (A \supset C)}{(A \supset B) \vee (A \supset C)} \vee E,2} \supset E,3}{\frac{[A]^4 \quad [\sim A]^5 \quad \frac{[\perp]^3}{C} \perp E}{A \supset C} \supset I,4}{(A \supset B) \vee (A \supset C)} \vee I,5} \supset I,6$$

5. Give sequent calculus proof in **G3cp** of  $\Rightarrow D$ , where  $D$  is

(a)  $\sim(\sim A \ \& \ \sim B) \supset A \vee B$

$$\frac{\frac{\frac{A \Rightarrow A, B, \perp}{\Rightarrow A, B, \sim A} \text{Ax} \quad \frac{B \Rightarrow A, B, \perp}{\Rightarrow A, B, \sim B} \text{Ax}}{\Rightarrow A, B, \sim A \ \& \ \sim B} \text{R\&}}{\frac{\perp \Rightarrow A, B}{\sim(\sim A \ \& \ \sim B) \Rightarrow A, B} \text{L\perp}} \text{L\&}}{\frac{\sim(\sim A \ \& \ \sim B) \Rightarrow A, B}{\sim(\sim A \ \& \ \sim B) \Rightarrow A \vee B} \text{RV}} \text{L\&}}{\Rightarrow \sim(\sim A \ \& \ \sim B) \supset A \vee B} \text{R\supset}$$

(b)  $(A \supset B) \supset \supset (\sim A \vee B)$

$$\begin{array}{c}
 \frac{\overline{A \Rightarrow \perp, B, A}^{Ax}}{\Rightarrow \sim A, B, A}^{R\supset} \quad \frac{\overline{B \Rightarrow \sim A, B}^{Ax}}{}^{L\supset} \quad \frac{\overline{A \Rightarrow B, A}^{Ax} \quad \overline{B, A \Rightarrow B}^{Ax}}{A, \sim A \Rightarrow B}^{L\supset} \quad \frac{\overline{A, B \Rightarrow B}^{Ax}}{}^{L\vee} \\
 \frac{\overline{A \supset B \Rightarrow \sim A, B}}{\overline{A \supset B \Rightarrow \sim A \vee B}}^{R\vee} \quad \frac{\overline{A, \sim A \vee B \Rightarrow B}}{\overline{\sim A \vee B \Rightarrow A \supset B}}^{R\supset} \\
 \frac{\overline{\Rightarrow (A \supset B) \supset (\sim A \vee B)}}{\Rightarrow (A \supset B) \supset \supset (\sim A \vee B)}^{R\&}
 \end{array}$$

(c)  $((A \supset B) \supset A) \supset A$  (Peirce's law)

$$\begin{array}{c}
 \frac{\overline{A \Rightarrow A, B}^{Ax}}{\Rightarrow A, A \supset B}^{R\supset} \quad \frac{\overline{A \Rightarrow A}^{Ax}}{}^{L\supset} \\
 \frac{\overline{(A \supset B) \supset A \Rightarrow A}}{\Rightarrow (A \supset B) \supset A}^{R\supset}
 \end{array}$$

(d)  $(A \supset B \vee C) \supset (A \supset B) \vee (A \supset C)$  (disjunction property under hypothesis)

$$\begin{array}{c}
 \frac{\overline{A, A \Rightarrow A}^{Ax} \quad \frac{\overline{A, A, B \Rightarrow B, C}^{Ax} \quad \overline{A, A, C \Rightarrow B, C}^{Ax}}{A, A, B \vee C \Rightarrow B, C}^{L\supset}}{A, A, A \supset B \vee C \Rightarrow B, C}^{L\supset} \\
 \frac{\overline{A, A \supset B \vee C \Rightarrow B, A \supset C}^{R\supset}}{A \supset B \vee C \Rightarrow A \supset B, A \supset C}^{R\supset} \\
 \frac{\overline{(A \supset B \vee C) \Rightarrow (A \supset B) \vee (A \supset C)}^{R\vee}}{\Rightarrow (A \supset B \vee C) \supset (A \supset B) \vee (A \supset C)}^{R\supset}
 \end{array}$$