

## Lecture 12 Parameter Estimation Summary

HW/SW Monitors  
Events/Sampling  
Overhead  
Example  
Summary

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## HW Monitors

- Collect raw data
- Wire monitors to physical units  
Fig. 2.7 & Tbl 2.3 [Ferrari 78]
- HW tool characteristics (1978)  
Tbl. 2.2 [Ferrari 78]
- Can be made to work with no (or very little) side effects or load  
– why?

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## Order Processing Example

- Create model, Fig 9.1
- How to get parameter values?
  - OS
  - TP monitor
  - DBMS
  - Application

Customers	disk A
Inventory	disk A?
Orders	disk B
Paging	disk C

Job classes:	Order entry	init R=3.1 sec
	Order query	

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## SW Monitors

- Modify SW to collect data  
Application? OS? DBMS?
- Use off-the-shelf monitoring system  
OS? DBMS? Other?
- Problem: runs in the same system that is being monitored  
– biased data. uses system resources
- Compare to HW monitors

Tbl 7.1 [Jain 91]

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## Measurement Process

- See Fig. 9.2
- Which data is needed?
- Where to place instrumentation?
- When to collect data?
- How long to collect data?
- How to account for measurement error?
  - load caused by the instrumentation?
  - does instrumentation have side effects?

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## Use Accounting Data for Parameter Estimation

- Good: it is there already (or is it?)
- Bad: it is built for accounting purposes
- Who used resources?
  - user id, program, project, account nr or class
- What resources were used?
  - CPU time, I/O op counts, NW use, mem use
- When were resources used?
  - user time only?
  - Granularity may be too coarse?

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## Program Analyzers

- Runs with some programs
- Designed for one program  
IBM DB2, IMS, CICS
- Generic analyzers
  - gprof
  - get nr. of transactions, CPU time, mpl

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## Events

- Predefined system events
  - “I/O request completed”
  - “process terminates”
- Create events
  - how much overhead?
- Log them into log files, or
- Send them to filtering and analysis

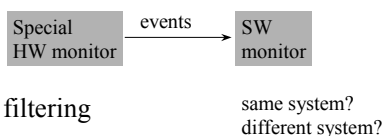
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## Hybrid Monitors

- HW + SW



- Data filtering
- Post run analysis
- Can be very large monitoring system
  - as large as actual monitored system

[Wybraniec &amp; Haban]

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## Sampling

- Good: smaller volume
- Bad: may miss something
- Sampling rate? too high? too low?
- Collect state information
  - “CPU busy”, “disk busy”
  - “amount of memory at use”
- Assign data to threads/processes/programs?
  - “CPU busy 5 ms due to process P” ???

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## Events or Sampling

- Event driven or time driven data gathering?
- Volume?
- Data collection method?
- Data storage? Where?
- Overhead?

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## Measurement Unit

- Module
- Subroutine
- Statement
  - language dependent?
- Machine instruction?

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## Time Unit

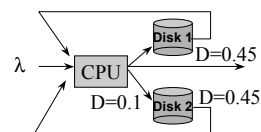
- Real time?
- System time?
- External time?
- CPU time?
- Elapsed time?

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## Parameter Values



- Job classes
  - clustering? factor analysis?
- $N_r$   $\lambda_r$   $Z_r$   $M_r$ 
  - external to system: plain measurements
- $D_{ir}$ 
  - internal: need internal data - e.g.,  $U_{ir}$
  - $D_{ir} = U_{ir} * T / C_{0r}$
  - sometimes get only  $U_i$

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## Instrumentation

- When?
  - source code
  - compiler generated code
  - before linking, or at load time
  - run time
- How?
  - source code? Obj. code?
  - run time env.? OS? HW?

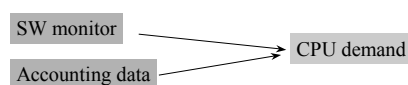
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## Parameter Example

Fig. 9.3 [Men 94]



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## Report

- Hierarchical
- Graphics
- What time is used?
  - inherited time
  - own time

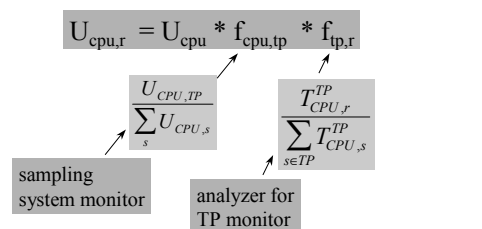
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## Transaction Processing System Data

- Have  $U_{total} = U_{cpu, OS} + U_{cpu, TP} + U_{cpu, prog}$
- Need  $U_{cpu, r}$  for each transaction class  $r$



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TP Example <sup>(1)</sup>

- System with 3 classes:
  - Batch (B), Interactive (I)
  - TP (queries & updates) → for our TP model we need demands for these!
- System monitor:  $T=1800$ ,  $U_{cpu}^{total} = 72\%$
- Accounting:  $U_{cpu,B}^{no OS} = 32\%$   $U_{cpu,I}^{no OS} = 10\%$   $U_{cpu,TP}^{no OS} = 28\%$
- TP analyzer:
  - 1200 queries in 120 sec
  - 400 updates in 140 sec
- Now,  $0.35 \text{ sec per update}$

$$U_{cpu,U} = 0.72 * \frac{28}{70} * \frac{140}{260} = 15.5\%$$

$$D_{cpu,U} = U_{cpu,U} * \frac{1800}{400} = 15.5\% * \frac{1800}{400} = 0.70 \text{ sec}$$

$$D_{cpu,Q} = \left[ 0.72 * \frac{28}{70} * \frac{140}{260} \right] * \frac{1800}{1200} = 0.20 \text{ sec}$$

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CPU Demand:  $D_{cpu,r}$ 

$$D_{cpu,r} = U_{cpu,r} * \frac{T}{C_{0,r}} = U_{cpu}^{Total} * f_{cpu,r} * \frac{T}{C_{0,r}}$$

where  $f_{cpu,r} = \frac{C_{0,r}}{\sum_s C_{0,s}} = f(\text{jobs})$

or  $f_{cpu,r} = \frac{U_{cpu,r}^{mon}}{\sum_s U_{cpu,s}^{mon}} = f(\text{time})$

or  $f_{cpu,r} = \frac{nrIO(r)}{\sum_s nrIO(s)} = f(nr \text{ IO's})$

or ...

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## Arrival Rate &amp; MPL

- Assume flow balance
  - arrival count  $\equiv$  completion count
  - $\lambda_r = C_{0r} / T$
- Average mpl:
 

$$\bar{N} = \lambda R = \frac{1}{T} * \frac{\sum e_i}{n} = \frac{\sum e_i}{T}$$

elapsed time for job i

nr of jobs

Fig 9.4 [Men 94]

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Capture Ratio <sup>(1)</sup>

- Capture ratio
  - class dependent!  $C_r = \frac{\text{measured } T_{cpu,r}}{\text{actual } T_{cpu,r}}$  measure estimate calibrate guess
- Example:

$$C_{batch} = 0.80 \quad C_{term} = 0.60$$

$$T = 7200 \text{ sec} \quad U_{cpu}^{Total} = 64\%$$

$$T_{cpu,batch}^{Meas} = T_{cpu,term}^{Meas} = 1500 \text{ sec}$$

$$T_{cpu,batch} = \frac{1500}{0.80} = 1875 \text{ sec (actual!)}$$

$$T_{cpu,term} = \frac{1500}{0.60} = 2500 \text{ sec} \quad T_{cpu} = 4375 \text{ sec}$$

$$U_{cpu,batch} = U_{cpu}^{Total} * f_{cpu,batch} = 0.64 * \frac{1875}{4375} = 0.27$$

actual, not measured

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M and Z <sup>(2)</sup>

- M - Number of terminals
    - session length (from accounting logs)
  - Z:
 

$$Z = \frac{M}{X} - R = \frac{M}{C_{0r}/T} - R$$

average measured elapsed time in system
- Example:  $T = 1 \text{ hour} = 3600 \text{ sec}$   
 $M = 40$  4900 interactive commands  
 2.5 sec aver resp time  
 $Z = 40 / (4900/3600) - 2.5 = 29.4 - 2.5 = 26.9 \text{ sec.}$

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 $D_{ir}$  for Disks

IO count → Mean service time

$$U_i^{Total} = \frac{IOC_i * MST_i}{T} \quad (\text{or from HW monitor})$$

$$D_{ir} = \frac{U_i^{Total} * T}{C_{0,r}} * f_{ir} = \frac{IOC_i * MST_i}{C_{0,r}} * f_{ir}$$

where  $f_{ir} = \frac{IOC_{ir}}{IOC_{0,r}}$  (user data)

or  $f_{ir} = \frac{SWAPC_{ir}}{SWAPC_{0,r}}$  (swap data)

or  $f_{ir} = \frac{PageInC_{ir}}{PageInC_{0,r}}$  (paging data)

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## Delay Node

- High disk utilization: include all disks in model
- Low (< 5%?) disk utilization: may use one aggregate **delay** node in model
  - no queueing likely ....

$$D_{ir} = \sum_{\text{disks } j} \frac{U_j^{\text{Total}} * T}{C_{0,r}} * f_{jr} = \frac{T}{C_{0,r}} \sum_{\text{disks } j} U_j^{\text{Total}} * f_{jr}$$

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## Calibration

- Use parameters to build baseline model
- Use solver
  - e.g., PMVA fig.9.1.out
  - get  $R_{OE} = 2.5$  (not 3.25 as in Tbl 9.1)
- What now? Calibrate model!
  - scale (or add to) output to match data
  - adjust MPL or its distribution fig.9.1a.out
  - add new class fig.9.1b.out
  - modify demands (at bottleneck)
  - add ghost server for unaccounted for load

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## Order Processing Example

Tbl 9.1 [Men 94]

- Want: Arrival rates, demands
- CPU demands Tbl 9.3
- File ops per application class Tbl 9.4
- File ops per disk Tbl 9.5
- Disk util. & demands per class Tbl 9.6
- $Z_r$ 's when  $R_r$ 's are known: use Little's Law

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## Estimate $Z_r$ 's if $R_r$ unknown

- Use subsystem as open model
- Use  $\lambda_r = X_{0r}$
- Solve model, get  $R_r = (3.00, 0.98, 2.23)$
- Use those estimates to compute  $Z_r = M_r / X_{0r} - R_r$

$$\begin{aligned} Z_{OE} &= 10/1.43 - 3.0 = 3.99 \quad (\text{vs } 3.75) \\ Z_{OI} &= 5/0.83 - 0.98 = 5.04 \quad (\text{vs } 5.04) \\ Z_{Ot} &= 5/0.31 - 2.23 = 13.9 \quad (\text{vs } 13.9) \end{aligned}$$

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## Summary

Probability Theory  
Modelling  
Solution Methods  
Approximations  
Parameter Estimation

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## Probability Theory

- Distributions
- Sample
- Variance
- Confidence intervals

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## Approximations

- Flow equivalent server
- Aggregate model
- Iterative solutions
  - multiple class memory
- Load concealment
- Shadow server

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## Modeling

- Capacity planning process
- System model
- Workload model
- Capacity
- Performance

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## Parameter Estimation

- Measurement process
- HW/SW monitors
  - hybrid monitors
- Accounting data
- Program analyzers
- Parameter values
  - combine data from various sources
  - fractions
  - overhead

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## Solution Methods

- M/M/1, Markov Chain
- Operational analysis
- Closed models
  - Convolution
  - MVA
  - Approximate MVA
- Open models

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