582746 Modelling and Analysis in Bioinformatics

Lecture 2: Network motifs

2.11.2015

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Outline

Network motifs

Statistical significance of motifs



Network motifs

- Network motifs are a way to analyze the local structure of a network:
 - What kind of local substructures (motifs, graphlets) does the network have
 - Assessing the statistical significance of these substructures





(b)



Fig. 2. An example graph (a), a pattern (b) and all different matches of the pattern (c, M₁ − M₂). The vertices of the graph and of the matches are numbered consecutively for identification purposes.

What is a motif?

- A motif is a statistically overrepresented pattern of local interactions in the network
- Overrepresentation = occurring more frequently than expected by chance
- The rationale is that overrepresentation may denote possible function
 - The motif has emerged several times
 - and it has been conserved in the evolution of the network

What is a motif?

- A motif is a small connected subgraph G' = (V', E')
- Size of motif is measured either by the number of vertices or the number of edges

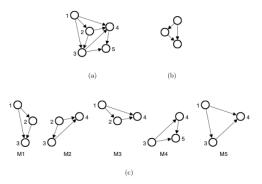


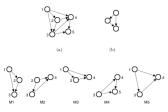
Fig. 2. An example graph (a), a pattern (b) and all different matches of the pattern (c, M₁ - M₅). The vertices of the graph and of the matches are numbered consecutively for identification purposes.

Types of motifs

- Motifs can be
 - Directed or undirected
 - Cyclic (loopy) or acyclic

matching the type of underlying network to be analyzed, e.g.

- Protein-protein interactions: undirected
- Gene regulatory interactions: directed, cyclic



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Matching motifs

- A match of a motif G' in the target graph G = (V, E) is a subgraph G'' = (V'', E'') which is *isomorphic* to motif G'
- ► Two graphs G' and G" are isomorphic if there is a bijective mapping between the edge and vertex identities
 - ▶ i.e. G' is transformed to G" by changing the vertex and edge identities

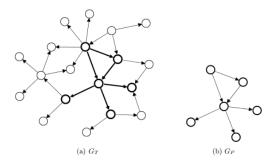
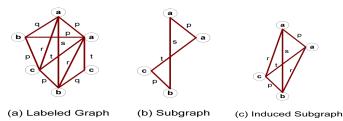


Fig.1. (a) A graph with a randomly selected subgraph (highlighted with bold lines). This subgraph is isomorphic to the graph G_P shown in (b). The highlighted subgraph in G_T is also a match of G_P in G_T . $\Rightarrow \quad \langle \Xi \rangle \Rightarrow \quad \Xi \quad \bigcirc \bigcirc \bigcirc \bigcirc$

Alternative definition: Induced subgraph

- Isomorphic induced subgraph (graphlet): a subgraph G" = (V", E") in G = (V, E) is accepted as a match only if it contains all edges of the original graph beteeen the nodes in V": mathematically we require that if e = (n_i, n_j) ∈ E and n_i, n_j ∈ V" then e ∈ E"
- Motivation: leaving out interactions from the motif may give false ideas of the biological function



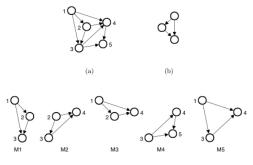
Hardness of isomorphism problems

- The complexity of graph isomorphism is in the 'grey area' of complexity:
 - It belongs to NP class of problems (problems where solution is easy to verify once found)
 - It is not known if graph isomorphism belongs to P class of problems (problems that can be solved efficiently)
 - It is not known if graph isomorphism is NP-complete (problems that are believed to be hard ot solve but easy to verify)

- Subgraph isomorphism, checking if a subgraph G" that is isomorphic to given graph G' exists in a larger graph G, is known to be NP-complete
- ► No hope for really fast algorithms for finding motifs.

Motif frequency

- How many times a motif occurs in the network to be analyzed?
- Depends on
 - Definition of a match (subgraph or induced subgraph)
 - Counting schemes for matches



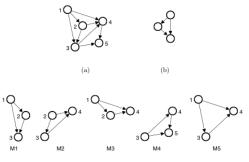
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Fig. 2. An example graph (a), a pattern (b) and all different matches of the pattern (c, $M_1 - M_5$). The vertices of the graph and of the matches are numbered consecutively for identification purposes.

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Counting schemes

- Possible schemes for counting matches
 - *F*₁: Two matches may overlap so that they share vertices and edges
 - *F*₂:Two matches may overlap so that they share vertices but not edges
 - ► *F*₃: Two matches may not overlap, they need to have disjoint sets of vertices



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Fig. 2. An example graph (a), a pattern (b) and all different matches of the pattern (c, M₁ − M₂). The vertices of the graph and of the matches are numbered consecutively for identification purposes.

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Counting schemes

	Graph elements shared by different matches		Values for the example in Fig. 2	
Concept	Vertices	Edges	Frequency	Selected matches
\mathcal{F}_1	yes	yes	5	$\{M_1, M_2, M_3, M_4, M_5\}$
\mathcal{F}_2	yes	no	2	$\{M_1, M_4\}$ or $\{M_3, M_4\}$
\mathcal{F}^*	no	yes	-	-
\mathcal{F}_3	no	no	1	one of $\{M_1, M_2, M_3, M_4, M_5\}$

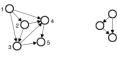








Fig. 2. An example graph (a), a pattern (b) and all different matches of the pattern (c, $M_1 - M_5$). The vertices of the graph and of the matches are numbered consecutively for identification purposes.

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Outline

Network motifs

Statistical significance of motifs

Statistical significance of motifs

- The frequency of a motif in some network does not directly tell us its importance
- Testing for statistical significance is more informative
 - How often we would expect to see this motif by chance in a similar random network

Need to formulate a *null hypothesis* and check the probability of the motif occurring as frequently under the null hypothesis

Testing fot statistical significance

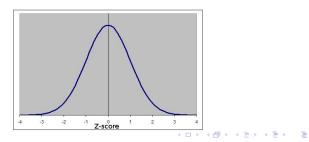
- For a null hypothesis:
- Estimate the probability distribution of the frequency of the motif in random networks
 - Analytically using a network model (e.g. ER networks)
 - By generating an ensemble of random networks
- Measure the statistical significance with Z-score or p-value

Measures of motif significance: Z-score

- Denote by *F*(*m*) the frequency of motif *m* and by *F_r(m*) and *σ_r(m*) the average and standard deviation of the motif frequency among the randomized networks.
- Z-score: "how far above the mean of the random networks"

$$Z(m) = \frac{\mathcal{F}(m) - \overline{\mathcal{F}_r(m)}}{\sigma_r(m)}$$

 Z-score above 2.0 is generally considered significant ("two standard deviations")

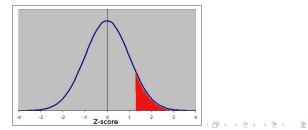


Measures of motif significance: P-value

P-value: "how often a random network has more motif ocurrences"

$$P(m) = \frac{1}{N} \sum_{r=1}^{N} \mathbf{1}_{\{\mathcal{F}_r(m) \geq \mathcal{F}(m)\}}$$

- ► 1_{A} denotes the indicator function, F_r(m) denotes the motif's frequency in r'th randomized network
- ▶ Requires a large number of randomized networks (≈ 1000) to be accurate
 - Estimating the tail of the distribution is harder than estimating its mean (as in Z-score)



Analytical approach using ER networks

- In ER networks an edge is present between two vertices with probability p
- Here we also allow self loops and also these edges are present with probability p
- The ER network should have a similar number of vertices and edges as the real network and so

$$p = \frac{E}{N^2}$$

where E is the number of edges in the real network and N is the number of vertices in the real network.

 Note that a directed network allowing self loops can have at most N² edges. Probability distribution of self loops in ER networks

The probability of having exactly k self loops is

$$P(k) = \left(egin{array}{c} N \ k \end{array}
ight) p^k (1-p)^{N-k}$$

The probability distribution is thus binomial with mean:

$$< N_{\rm self} >= Np = N \frac{E}{N^2} = \frac{E}{N}$$



and with variance (approximation via Poisson distribution)

$$\sigma_{\rm self} = \sqrt{\frac{E}{N}}$$

Z-score

The E. coli transcriptional network has 424 vertices and 519 edges (note that this is a different version of the network than what we use in the exercises):

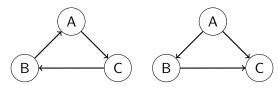
$$< N_{
m self} > = rac{E}{N} = rac{519}{424} = 1.2$$
 $\sigma_{
m self} = \sqrt{rac{E}{N}} = 1.1$

The real network has 40 self loops:

$$Z = rac{N_{
m self} - < N_{
m self} >}{\sigma_{
m self}} = rac{40 - 1.2}{1.1} = 32$$

Z-score is very high and thus the high number of self loops in the *E. coli* transcriptional network is statistically significant

- Consider a pattern graph G with n vertices and g edges
- How often would such a pattern occur in ER networks?
- ▶ We will use counting scheme *F*₁ (vertices and edges can overlap)



Motif 1

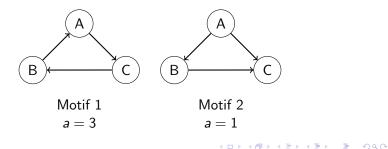
Motif 2

To generate an instance of this pattern in a random graph, we need to choose *n* vertices and place the *g* edges in appropriate places:

$$< N_G > = a^{-1} \cdot N \cdot (N-1) \cdot \ldots \cdot (N-n+1) \cdot p^g$$

 $\approx a^{-1} N^n p^g$

where a is the number of permutations of vertex labels of G that give the same graph.



The mean connectivity of a network is

$$\lambda = \frac{E}{N}$$

and then we get

$$< N_G > \approx a^{-1} N^n p^g$$

 $= a^{-1} N^n \left(\frac{E}{N^2} \right)^g$
 $= a^{-1} \lambda^g N^{n-g}$

 If we assume that the mean connectivity is constant regardless of the size of the network, then the number of subgraphs scales as

$$< N_G > \sim N^{n-g}$$

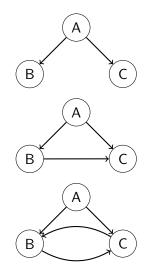
 V-shaped subgraphs (3 nodes, 2 edges) thus scale linearly with the size of the network:

 $< N_{\rm V-shaped} > \sim N$

Number of triangle shaped subgraphs (3 nodes, 3 edges) stays constant:

$$< N_{
m triangle} > \sim N^0$$

 Subgraphs with 3 nodes and more than 3 edges become rarer when the network gets larger



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A subgraph in ER networks and *E. coli* transcriptional regulation network

The E. coli transcriptional regulation network has 424 vertices and 519 edges:

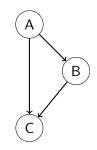
$$< N_{\rm G} > \approx a^{-1} \lambda^{g} N^{n-g} = \left(\frac{519}{424}\right)^3 424^{3-3} = 1.7$$

 The distribution of the motif in ER networks can be approximated by a Poisson distribution and thus the standard deviation is

$$\sigma_{\rm G} \approx \sqrt{} = 1.3$$

The E. coli transcriptional regulation network features 42 instances of the motif and so we get:

$$Z = \frac{42 - 1.7}{1.3} = 31$$



Null hypothesis from random networks

- In traditional hypothesis testing, one typically analytically formulates a probability distribution for the values of the random variable of interest (here frequency of a motif)
- In network analysis, analytically determining a suitable probability distribution may be difficult
- Instead, randomization tests are being used: a large set of random networks of appropriate structure are generated and the average frequency of the motif together with its variance is recorded.
- Computationally demanding process if the networks are large

Randomization algorithm for Null model networks

- Typical method for null model generation is to take the original network being analyzed and make large number of randomized versions of it by modifying the network by a large number of random edit operations
- Commonly used edit operation is to rewire the network locally:
 - ► Take two edges (*A*, *B*) and (*C*, *D*) and replace them with edges (*A*, *D*) and (*C*, *B*)
 - Preserves degree distribution of nodes
 - If the nodes are chosen from a small neighborhood, also keeps average path length close to original

Randomization algorithm for Null model networks

Additional criteria to be preserved can be set, e.g.

- Preserve number of bidirectional edges
- ► Preserve number of motif of size n − 1 when searching for motifs of size n

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As a guideline, the null model should be as close to the original as possible, but randomize the property of interest.

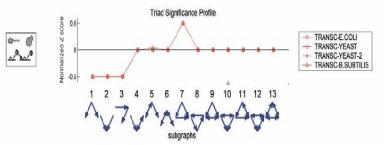
Motif significance profile

 Motif significance profile SP is a vector of normalized Z-scores for a particular set of motifs

$$SP = ((SP(m_1), \ldots, SP(m_2)),$$

where $SP(m) = Z(m_i) / \sqrt{\sum_j Z(m_j)^2}$.

- Motif significance profile allows comparing different size networks in terms of the motifs they contain
- Typically, the set of motifs contains all motifs of particular size



Motifs and antimotifs in PPI networks and internet router network

Pattern	Protein interactions	Internet routers
_	Not a motif	Not a motif
	C = 0.981	C = 0.977
	Motif $(Z = 48)$	Motif ($Z = 4600$)
$ \bigtriangleup $	C = 0.019	C = 0.023
	Motif $(Z = 15)$	Not a motif
$ \wedge $	C = 0.680	<i>C</i> = 0.931
	Anti-motif $(Z = -19)$	Motif $(Z = 18)$
	<i>C</i> = 0.024	<i>C</i> = 0.013
	Anti-motif $(Z = -18)$	Anti-motif $(Z = -7)$
	<i>C</i> = 0.292	<i>C</i> = 0.048
	Not a motif	Motif (Z = 356)
	C = 0.0013	C = 0.004
	Anti-motif $(Z = -4.5)$	Motif (Z=137)
	C = 0.0019	C = 0.002
	Not a motif	Motif (Z ND)
\mid	<i>C</i> = 0.0004	<i>C</i> = 0.0005

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Hardness of motif discovery

Several challenging subproblems:

- Graph isomorphism testing: required to check if two motifs are in fact the same. No polynomial time algorithm is known for this problem.
- Number of motifs: grows exponentially in the size of the motif. Especially with directed motifs grows very fast.
- ► Number of matches: theoretically the worst case number of potential matches is O(|E_t|^{|E_m|}) where E_t and E_m are the number of edges in the target and motif, respectively.
- Size of analyzed networks affects the above steps via the number of different patterns and matches that can be found.
- Calculation of statistical significance via randomization calls for generation and motif discovery from a large number of networks, multiplying the computation time of all the above points.

Study Group on Thursday

Group 1: Students whose first name starts with A-L

- F. Schreiber and H. Schwbbermeyer: Frequency concepts and pattern detection for the analysis of motifs in networks. Trans. on Comput. Syst. Biol: III, pp. 89–104, 2005.
- Concentrate on section 4.
- Group 2: Students whose first name starts with M-Z
 - N. Kashtan, S. Itzkovitz, R. Milo and U. Alon: Efficient sampling algorithm for estimating subgraph concentrations and detecting network motifs. Bioinformatics 20(11):1746–1758, 2004.

Concentrate on the Methods section.