582746 Modelling and Analysis in Bioinformatics

Lecture 3: Global Network Models

20.9.2016

Outline

Examples of biological networks

Global properties of networks

Distance measures

Degree measures

Local clusters

Network Models

Erdös-Renyi Model

Watts-Strogatz Model

Barabasi-Albert Model

More Network Properties

Statistical Testing of Network Properties

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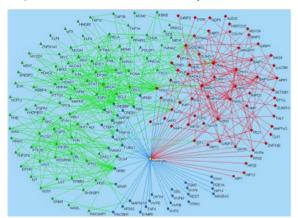
More Network Properties

Statistical Testing of Network Properties



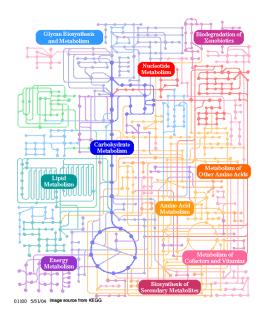
Protein-protein interaction network

- Vertices are proteins
- ▶ The proteins are connected if they interact with each other.



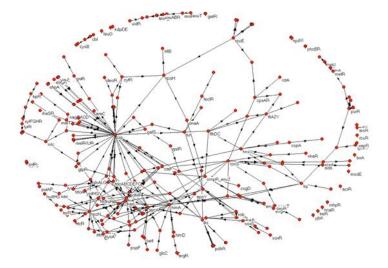
Metabolic network

- Vertices are metabolites, i.e. chemical compounds
- Edges describe how the cell can transform a metabolite into another



Gene regulatory network

- Vertices are genes
- ▶ Genes are linked if one regulates the other



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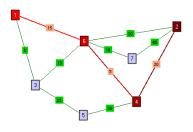
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(Shortest path) distance

- ▶ Distance d_{ij} is the length of the shortest path between vertices n_i and n_j, i.e. the minimal number of edges one needs to traverse to get from n_i to n_j
- The shortest path may not be unique, but the length of the shortest path is unique
- In directed network, we may have $d_{ij} \neq d_{ji}$
- ▶ If there is no path between n_i and n_j , we have $d_{ij} = \infty$

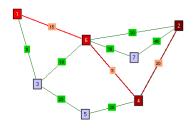


Ignoring weights, above we have $d_{12}=2$, $d_{13}=1,d_{14}=2,...$

Diameter and average path length

- ▶ The diameter $d_m = \max(d_{ij})$ is the maximal distance between any two nodes (= the longest shortest path)
- Average or characteristic path length

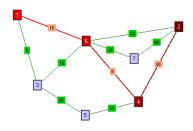
$$d = \langle d_{ij} \rangle = \frac{1}{N_V^2} \sum_{i=1}^{N_V} \sum_{j=1}^{N_V} d_{ij}$$



Ignoring weights, above $d_m = 3$, $d \approx 1.57$

Efficiency

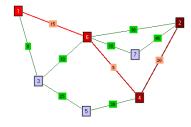
- ▶ Efficiency, or average inverse path length: $d_{eff} = \langle 1/d_{ii} \rangle$
- Useful when average path length is infinite (disconnected network)
- Fully connected network has efficiency $d_{\it eff}=1$, graph with no edges has $d_{\it eff}=0$



Ignoring weights, above $d_{eff} \approx 0.73$

Weighted graphs

- If the edges in the graph have associated weights w_{ij}, it is natural to define distances based on the weights:
- d_{ij} as the sum of weights in the minimum weight path between n_i and n_j
- Maximum and average path length as well as efficiency naturally generalize by changing the distance measure to the weighted version



With weights, $d_{12}=40$ (red path), $d\approx 24.2$, $d_m=50=d_{25}, d_{eff}=0.06$

Finding shortest paths

Finding shortest paths in graphs is part of classical algorithm theory, two efficient algorithms

- ▶ Dijkstra's algorithm: given a vertex find shortest paths to all other vertices, basic implementation runs in $O(N_V^2)$ time, can be implemented faster for sparse graphs
- ▶ Floyd-Warshall algorithm: find shortest paths for all pairs of vertices in the graph in $O(N_V^3)$ time; outputs a distance matrix $(d_{ij})_{i,i=1}^{N_V}$ in same time.
- ▶ Both work with weighted formulations

Shortest path distances in empirical networks

Path length analysis of many networks that occur in nature reveals the small-world property

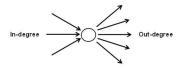
- Metabolite graphs: average path length $d \approx 3 \; (N_V \approx 10^3 10^4)$
- $ightharpoonup ext{WWW: links chains between two} \ ext{web documents } d pprox 16 \ (N_V > 10^9)$
- Frdös number: shortest co-author chain to Paul Erdös, d=4.65 ($N_V \approx 4 \times 10^5$)
- **.**..



Paul Erdös (1913-1996, a Hungarian mathematician, published over 1400 scientific papers over his lifetime with over 500

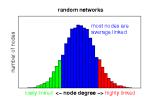
Node degree

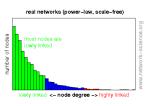
- Degree k_i of vertex n_i is the number of edges adjacent to a vertex
- In a network without self-loops and without multiple edges between any pair of edges: degree
 number of neighbours
- In directed networks: in-degree is the number of incoming edges and out-degree is the number outgoing edges



Degree distribution

- ▶ Given a fixed set of vertices, p(k) denotes the probability that a randomly chosen vertex has degree k.
- ▶ (Empirical) degree distribution is the list of probabilities (or relative frequencies) p(k), $k = 0 ... N_V$.
- Analysis of the degree distribution is an important means to characterize networks

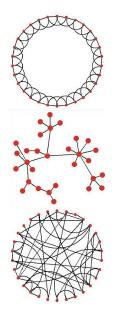




Degree distribution

Fitting the empirical degree distribution to a theoretical distribution given by a mathematical law is an important tool for network analysis

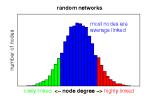
- ▶ Regular lattice: $p(k) \approx 1$, where k is a constant
- ▶ Scale free network: $p(k) \propto k^{-\gamma}$
- Random network: $p(k) \propto {N_V-1 \choose k} p^k (1-p)^{N_V-1-k}$

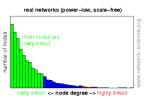


Degree distribution

The degree distributions of of scale-free network and random network look markedly different

- ▶ Scale free network: $p(k) \propto k^{-\gamma}$ (power law, heavy tail)
- ▶ Random network: $p(k) \propto {N_V-1 \choose k} p^k (1-p)^{N_V-1-k}$ (binomial, light tail)





Fitting degree distributions

- ▶ Typically the fitting of the empirical distribution is based on the histogram of observations for p(k)
- ► This is prone to errors in the region of high degree nodes due to low number of observations
- ▶ Binning can help: divide the range of *k* into intervals and put all observations in the interval into a common bin
- ▶ Cumulative degree distribution $p_c(k) = \sum_{l=k}^{\infty} p(l)$, the likelyhood that a given node has degree at least k, is more reliable and does not require binning

Degree correlations and assortative mixing

Degree correlation is a statistic that reveals additional information of the connection patterns of the nodes

- Assortative networks: high correlation between the degrees of adjacent nodes; highly connected nodes mostly connect to other highly connected nodes
- Disassortative networks: highly connected nodes mostly connect to low degree nodes
- Assortativity index $-1 \le r \le 1$: Pearson correlation coefficient of degrees of adjacent nodes, r > 0 assortative, r < 0 disassortative

Examples

Social networks are typically assortative, technological and biological networks tend to be disassortative

	Group	Network	Type	Size n	Assortativity r	Error σ ,
	a	Physics coauthorship	undirected	52 909	0.363	0.002
	a	Biology coauthorship	undirected	1 520 251	0.127	0.0004
	b	Mathematics coauthorship	undirected	253 339	0.120	0.002
Social	c	Film actor collaborations	undirected	449 913	0.208	0.0002
	d	Company directors	undirected	7 673	0.276	0.004
	e	Student relationships	undirected	573	-0.029	0.037
	f	Email address books	directed	16 881	0.092	0.004
	g	Power grid	undirected	4 941	-0.003	0.013
Technological	h	Internet	undirected	10 697	-0.189	0.002
	i	World Wide Web	directed	269 504	-0.067	0.0002
	j	Software dependencies	directed	3 162	-0.016	0.020
	k	Protein interactions	undirected	2 115	-0.156	0.010
	1	Metabolic network	undirected	765	-0.240	0.007
Biological	m	Neural network	directed	307	-0.226	0.016
	n	Marine food web	directed	134	-0.263	0.037
	o	Freshwater food web	directed	92	-0.326	0.031

(M. Newman. Phys. Rev. E 67, 026126 (2003))

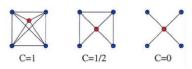


Clustering coefficient

- Clustering coefficient measures the probability that two vertices with a common neighbor are connected
- Let E_i denote the number of edges between the neighbors of v_i , and $E_{max} = k_i(k_i - 1)/2$ the theoretical maximum. Clustering coefficient for vertex n_i is now

$$C_i = \frac{E_i}{E_{max}} = \frac{2E_i}{k_i(k_i - 1)}$$

 Clustering coefficient for the whole graph is obtained by averaging over the vertices

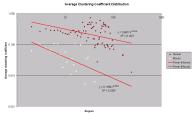


 $\label{lem:http://www.nd.edu/~swuchty/Download/Wuc.pdf, p. 24, Fig. 2} http://www.nd.edu/~swuchty/Download/Wuc.pdf, p. 24, Fig. 2$

Clustering coefficient in natural networks

- Natural networks often have relatively high clustering coefficient indicating local clustering within the network
- Negative correlation between the degree and the clustering coefficient has also been observed;
 - Low degree nodes lie in local clusters, while the neighbors of high degree nodes are less often connected
 - Indicates modular network structure

Example: PPIs in Mouse and Human:

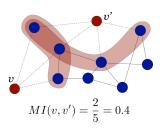


(http://bccs.bristol.ac.uk/toProgramme /_project/2008/Angela_Onslow_S08/)

Matching index

To be functionally related, two vertices do not need to be connected, examples:

- Two transcription factor proteins regulating the same gene
- Two metabolite molecules taking part in similar reactions



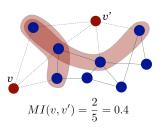
Zamora-Lopez. Frontiers in Neuroinformatics 4, 2010

Matching index

Matching index measures the amount of neighbors the two nodes share:

$$MI_{ij} = rac{Shared_{ij}}{k_i + k_j - Shared_{ij}}$$

 Similarity in terms of perceiving the neighborhood similarly



Zamora-Lopez. Frontiers in Neuroinformatics 4, 2010

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Models of complex networks

- Theoretical models of networks are needed as a basis for comparison to determine the significance of global properties or non-trivial substructures of natural networks.
- We will look at three specific models
 - Erdös-Renyi Model
 - Watts-Strogatz Model
 - Barabasi-Albert model



Erdös-Renyi Model

- ► ER network consists of N_V vertices
- Edge is drawn between a pair of nodes randomly with probability p
- ▶ Degree distribution of the ER model is binomial: $p(k) \propto {N_V-1 \choose k} p^k (1-p)^{N_V-1-k}$
- Degree distribution can be approximated by Poisson distribution for large graphs



Erdös-Renyi Model

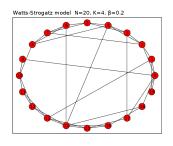
Significant body of theoretical research exists for the ER model, e.g.

- ▶ For $N_V p < 1$ the network almost surely has no large connected components
- ▶ For $N_V p \approx 1$ the network will almost surely have one large connected component
- For N_Vp > log N_V the network will almost surely be connected
- ▶ ER network has the small-world property when $p>1/N_V$ with average path length scaling as $I\sim \log N_V$
- No local clustering, expected clustering coefficient $C = p = \langle k \rangle / N_V$ for all nodes

Watts-Strogatz model

- 1. Arrange vertices in a ring structure
- 2. Connect each vertex to *K* closest neighbours
- 3. With probability p_{rew} , rewire each each edge by detaching from one end and attaching to a randomly chosen vertex.

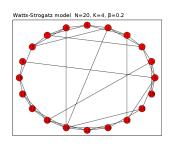
After steps (1-2) there is local clustering, step (3) lowers average path length by creating shortcuts



 $http://en.wikipedia.org/wiki/File:Watts_strogatz.svg$

Watts-Strogatz model

- Even for low rewiring probability $(p_{rew} << 1)$ the average path length goes down rapidly
- Small average path length and local clustering is retained for intermediate p_{rew}
- ▶ When $p_{rew} \rightarrow 1$, we get ER model, i.e. local clustering is destroyed
- Degree distribution is similar to ER graph: homogeneuos and peaked around k = K

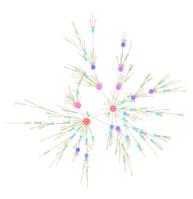


http://en.wikipedia.org/wiki/File:Watts_strogatz.svg

Barabasi-Albert model

- Start with an initial small connected network of N₀ vertices
- Iteratively add new vertices and connect the new vertex to m ≤ N₀ vertices
- Draw the nodes that will be connected to the new vertex with probability proportional to their degree (preferential attachment):

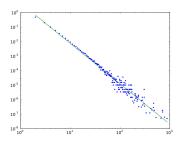
$$\rho(n_i) = k_i / \sum_i k_j$$



 $(\mathsf{BA}\ \mathsf{graph}\ \mathsf{from}\ \mathsf{http://melihsozdinler.blogspot.com/})$

Barabasi-Albert model

- Unlike ER or WS model, Barabasi-Albert model explain the inhomogeneuos degree distribution observed in natural graphs
- ▶ With enough iterations, the degree distribution of the BA model is scale-free, with $p(k) \sim k^{-3}$
- Average path length in BA networks has been found to be smaller than in ER and WS models



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Robustness and Attack tolerance

- Robustness against pertubations (mutations, environment changes) is a preferable property for biological networks
- Networks analysis is interested in preservation of network topology under perturbations (usually: removals of vertices or edges)





(a) Random network

(b) Scale-free network

Robustness and Attack tolerance

- Both ER networks and scale-free networks (such as BA model) are robust towards random deletions of nodes and connections
 - A random mutation is likely to hit a low degree node in BA model
- Scale-free networks are not robust towards intentional attacks
 - Removal of a set of highly connected nodes may collapse the global structure
 - "Robust, yet fragile"



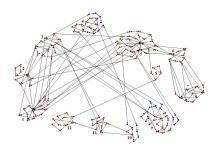


(a) Random network

(b) Scale-free network

Modularity and hierarchical organization

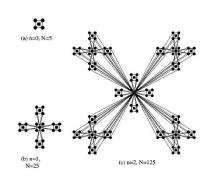
- Many natural networks are observed to posses modular structure with densely connected functional clusters of nodes that are sparsely connected to other nodes.
- Also, hierarchical organization of network structure can be observed
- The random network models discussed above, do not directly explain these phenomena



(Zhao et al. BMC Bioinformatics 2006, 7:386)

Modularity and hierarchical organization

- Barabasi and Albert model has been later extended to that direction
 - Based on replicating basic modules and wiring them to the central module of rest of the network
 - Recursive application leads to hierarchical organization
 - Deterministic rather than random procedure



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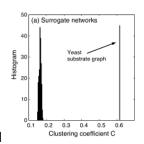


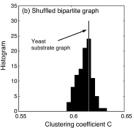
Statistical testing of network properties

- ► How to determine if an observed property of the network is significant or if it occured just by chance?
- Set up a null hypothesis
- Test if the observed property is consistent with the null hypothesis

Statistical testing of network properties: Example

- Suppose that we have observed a clustering coefficient C for a given network. Is the network highly clustered?
- Null hypothesis: The clustering coefficient is consistent with a network of the same size and degree distribution.
- Create an ensemble of random networks with same size and degree distribution and compute the clustering coefficient of each network
- Reject the null hypothesis if the probability of a network with clustering coefficient of at least C is low enough
- ▶ If the null hypothesis can be rejected, we can conclude that the network is highly clustered as compared to the null model.







What next?

- ► Thursday 10-12: Study group on global network models
- ► Thursday 12-14: Exercise session