

# Lecture 4

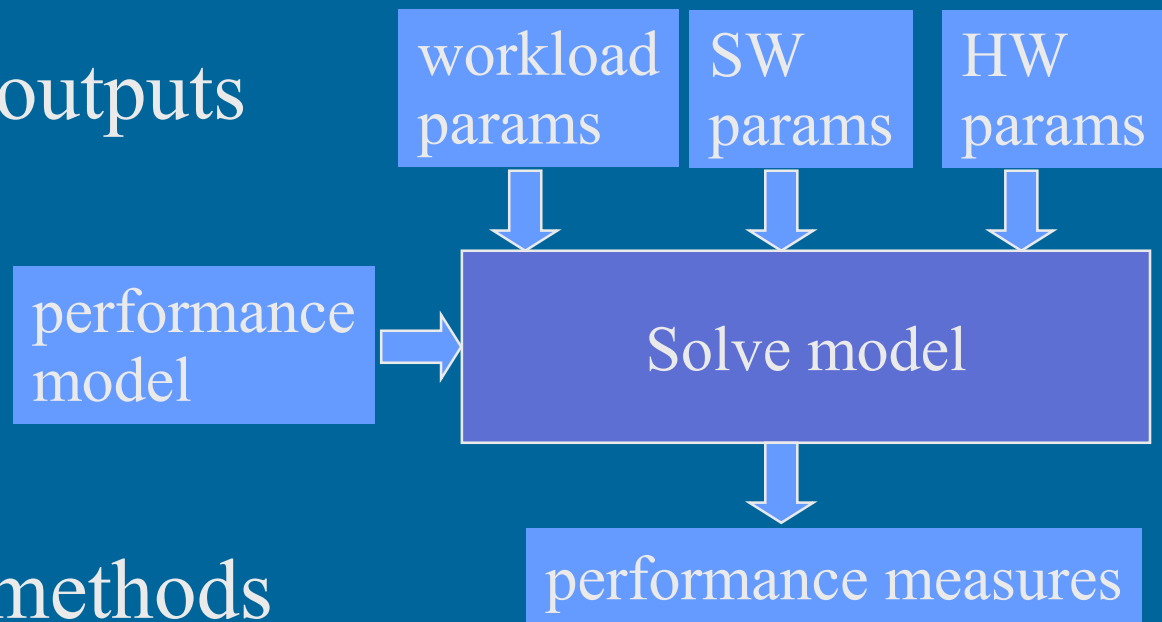
## Performance Evaluation Models



Building a Model  
Multiple Class Models  
Baseline Model  
Modification Analysis

# Performance Model

- Inputs & outputs



- Solution methods

- depend on model
- trivial: rules of thumb
- complex: analytical, simulation, benchmarks

# Example: Database Server

Query transactions



0.5 tps

Multiprogrammed  
OS

CPU

Disk 1

memory


Disk 2

5 partitions

(I.e., max 5 transactions  
processed at the same time)

- What if 1.0 tps?
  - need faster CPU? or more memory?
- Queues? Resources? Active? Passive?
- Use of resources? service time?

# Typical Transaction T

- Acquire memory partition
    - queue for memory?
  - Use CPU ....
    - queue for CPU?
  - .... until
    - I/O operation: use (and queue for) disk
    - timeslice expires: give CPU to next job
    - transaction completes: release mem & depart
- 

# Active resource

- Server, device
- Gives service
- Must have in possession during service
- Kept only during service
- Waiting queue or line
- Speed or rate of service, service time
  - parameter to model?
  - aver. value? distribution?
- Fig. 3.2 [Men 94]

E.g., aver 4.6 ms

CPU, Disks



# Passive Resource

- Allocated, reserved
- Deallocated, freed
- Waiting queue or line
- Must have in order to proceed
- Kept until deallocated
  - difficult for Markov Chain based analytical solutions
  - trouble: simultaneous resource possession
- Figs 3.3 & 3.4

Memory

# Queueing Network (QN)

- Network of Queues
- Open Queueing Network

Operational Analysis

- Database server: Fig. 3.5 [Men 94]
- system, arrivals
- transitions, transition probabilities
- queues, queue lengths
- subsystem, easy to solve!

- Service demands ( $D_i$ )
  - Tbl 3.1

Utilization

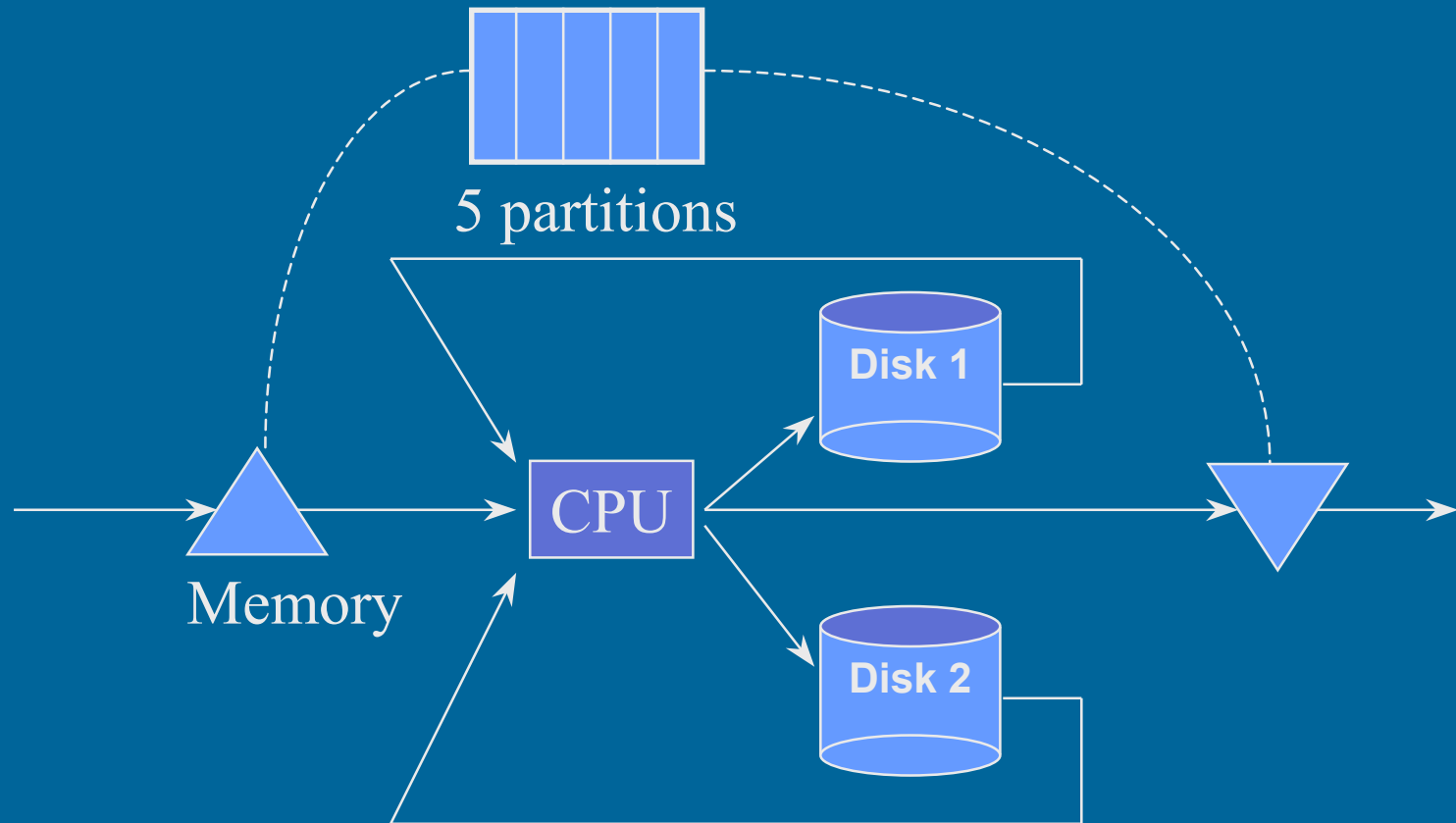
Total  
observation  
time

$$D_i = (U_i * T) / C_0$$

$$D_i = V_i S_i$$

Completions  
from system

# Open Queueing Network





# Service Demand $D_i$

Device Utilization

Total observation time

$$D_i = (U_i * T) / C_0$$

Completions  
from system

$$= U_i / (C_0 / T) = U_i / X_0$$

System throughput

$$= V_i S_i$$

Device Service Time

Device Visit Ratio = Nr of visits to device

# Example Open QN Model

- Fig. 3.5
- Minimum response time:
- Queueing time?

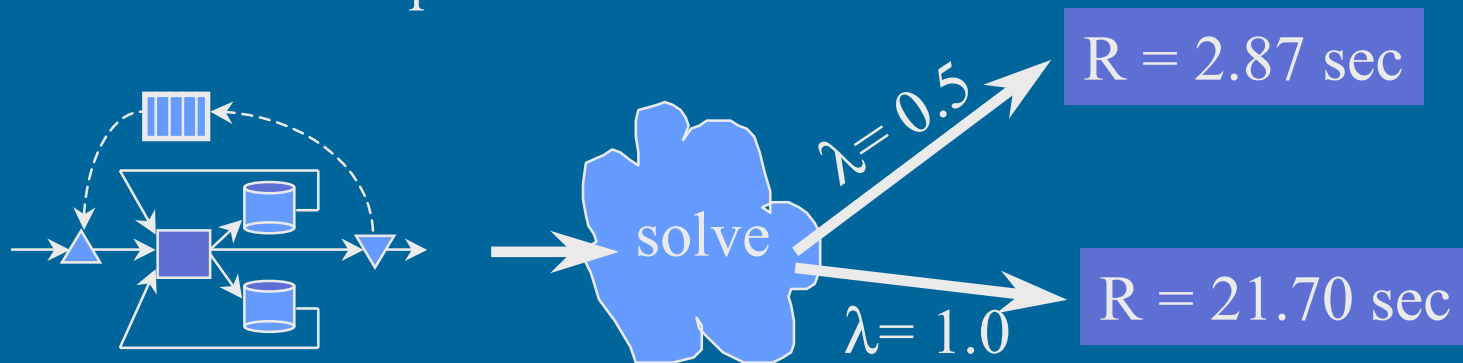
$$T_0 = \sum_{i=1}^K D_i$$

$$T = T_0 + W_{mem} + W_{cpu} + W_{disk1} + W_{disk2}$$

- Average arrival rate  $\lambda = 0.5$  tps
- Maximum degree of multiprogramming
  - how many jobs in subsystem?
  - $N^{\max} = 5$

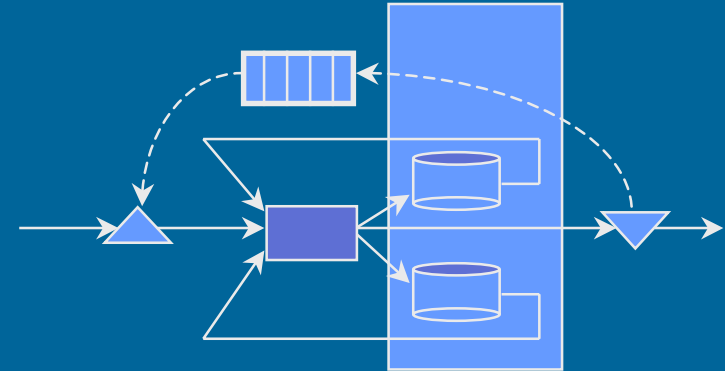
# How to Solve System Model?

- Depends on arrival rate!
  - Tbl 3.2
- Easy to solve with light load
- More complex to solve with heavy load
- What is  $\lambda_1$  ?



# Disk Subsystem

- I/O channels? SCSI?
- Heads of strings?
- Device controllers
- Disk cache
- File access protocol
- Rotation speed?



Average demand  
 $S_i = 0.75$  sec

OK or not OK?

# Interactive Systems

(osituskäyttö  
systeemit)

- People are part of model
- Terminals, work stations, ...
- Response time (R), Think time (Z)
- Fig. 3.6
- Tbl 3.7
- How to solve?
  - depends on number of terminals
  - Tbl 3.4



# Batch Systems

(eräajosteemit)

- No people
- Closed system
- Fixed number of multiprogramming level
- Tbl 3.5

# Multiple Job Classes

- More difficult than before
  - More parameters to estimate
  - More complex to solve
- Gives more usable information
- Open model: class arrival rates
- Closed model: class populations
- Tbl 3.6

# Aggregating Classes

(luokkien yhdistely)

- Combine classes to make model simpler
  - aggregate all uninteresting job classes together?
- To make parameter estimation simpler
- Tbls 3.7 & 3.8
- Need to compute derived parameter values for aggregate class from those of component classes
  - method varies depending on network type
  - Figs 3.9, 3.10 & 3.11

# Priorities

- Priorities are used in real systems
  - CPU, disks, etc
- May be dynamically changing
- Difficult to model well
- Models with priorities are more complex to solve (than those with no priorities)
- Example with Tbl 3.9

# Shared Domains

(yhteisalueet)

- Class limits or passive resources shared with other classes
  - multiprogramming level
  - memory partitions
  - Fig. 3.12





# Multiple Class Model Parameters

- Tbl 3.10

# Baseline Model and Modification Analysis

- Baseline: Tbl 3.11
- A: Use DBMS: Tbl 3.12
- B: Use DBMS + Optimizing compiler
  - $D_{\text{CPU}}$  down 50%
  - only for applications, 40% of CPU path length
  - So,  $D_{\text{CPU}}$  really down only 20%





# Baseline Model and Modification Analysis (contd)

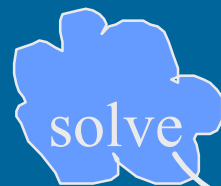
- C: DBMS + Larger DB record buffer pool (I.e., larger disk cache)
  - $D_{\text{DISK1}}$  and  $D_{\text{DISK2}}$  down 30%
  - Tbl 3.13

# Baseline Model and Modification Analysis (contd)

- D: Use also transaction logging for crash recovery?
  - disk update needs 2048 B record (38.7 msec)
  - logging only for complex transactions
  - assume each complex transaction causes one disk update, and so one log update (to DISK1)
  - $D_{\text{DISK1}}$  up 0.0387 sec
  - Tbl 3.14

HW specs

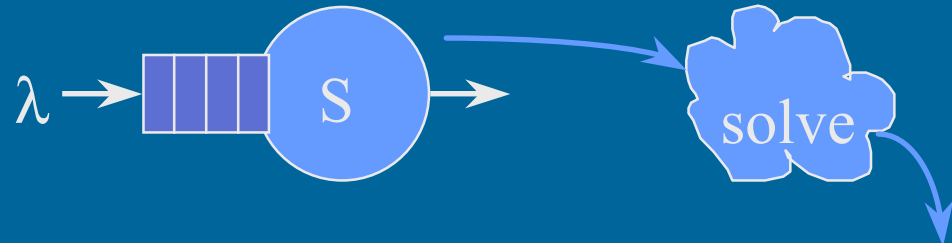
SW specs



	trivial	complex
R	1.17	2.26

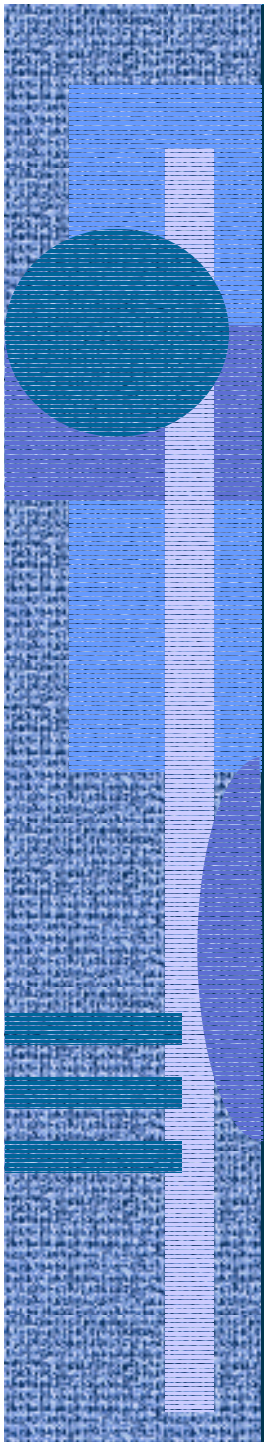
# Solution Methods

- Analytical
  - much more complex for complex models
- Approximate
  - approximate reality with simpler model
  - approximate exact solution for complex model
- Simulation
  - Monte Carlo
  - statistical analysis



$$R = \frac{S}{1 - \lambda S}$$





11.3.2002

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