

## 58093 String Processing Algorithms (Autumn 2016)

### Exercises 4 (Tuesday, November 22)

1. The Knuth–Morris–Pratt algorithm differs from the Morris–Pratt algorithm only in the failure function, which can be defined as

$fail_{\text{KMP}}[i] = k$ , where  $k$  is the length of the longest proper border of  $P[0..i)$  such that  $P[k] \neq P[i]$ , or  $-1$  if there is no such border.

- (a) Compute both failure functions for the pattern `ananassana`.
  - (b) Give an example of a text, where some text character is compared three times by the MP algorithm but only once by the KMP algorithm when searching for `ananassana`.
2. A pattern matching automaton for a string  $P$  is an automaton that recognizes the language  $\Sigma^*P$ , i.e., all strings that end with  $P$ . Draw the following types of pattern matching automata for  $P = \text{anas}$ :
- (a) MP-automaton (see slide 79 in lecture notes)
  - (b) KMP-automaton
  - (c) Non-deterministic finite automaton (see slide 84 in lecture notes)
  - (d) Deterministic finite automaton

*Hint: In (a) and (b), you can get the failure transitions from Exercise 4.1 above.*

3. Let us analyze the average case time complexity of the Horspool algorithm, where the average is taken over all possible patterns of length  $m$  and all possible texts of length  $n$  for the integer alphabet  $\Sigma = \{0, 1, \dots, \sigma - 1\}$  where  $\sigma > 1$ . This is the same as the expected time complexity when each pattern and text character is chosen independently and randomly from the uniform distribution over  $\Sigma$ .
- (a) Show that the average time spent in the loop on line 7 is  $\mathcal{O}(1)$ .
  - (b) Show that the probability that the shift is shorter than  $\min(m, \sigma/2)$  is at most  $1/2$ .
  - (c) Combine the above results to show that the average time complexity is  $\mathcal{O}(n/\min(m, \sigma))$ .
4. Simulate the execution of the BNDM algorithm for the pattern `anna` and the text `bananamanna`.
  5. Show how the following (single) exact string matching algorithms can be modified to solve the *multiple exact string matching problem*:
- (a) Shift-And
  - (b) Karp-Rabin

The solution should be more efficient than the trivial one of searching each pattern separately.